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**Reference Manual**

**IBM 7070 Series Programming Systems**

**Autocoder**

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# Introduction

IBM 7070/7074 Autocoder is a symbolic programming system designed to simplify the preparation, correction, and interpretation of programs for the IBM 7070 and 7074 Data Processing Systems. This manual is a reference text and contains information to enable the programmer to use the Autocoder system.

Autocoder is a component of the 7070/7074 Compiler Systems Tape and forms an interlocking system with FORTRAN and the Report Program Generator. The Compiler Systems Tape may be obtained by sending a full reel of magnetic tape to the IBM Program Librarian. The compilers will be written on the tape which will be returned. It is suggested that a duplicate of the tape be made as soon as it is received so that one copy of the program can be kept in reserve. The extra copy should be used only in case the working copy becomes unusable. The reel of magnetic tape for the 7070/7074 Compiler Systems Tape should be sent to:

IBM 7070/7074 Program Librarian  
International Business Machines Corporation  
590 Madison Avenue  
New York 22, New York

This manual assumes that the programmer is familiar with the methods of data handling and the functions of instructions in the 7070/7074 Data Processing System. This information is included in the IBM Reference Manual "7070 Data Processing System," form A22-7003-2.

The 7070/7074 Autocoder system is designed for use in installations which have a minimum of six 729 Model IV (or Model II) tape units and a machine with 5,000 words of core storage. This minimum configuration permits compilation of programs whose input is in the form of a card image tape and whose output is to be written on tapes for printing and/or punching off-line. The following additional equipment may be added to perform the indicated operations:

1. IBM 7500 *Card Reader* is necessary, in addition to the six tape units, if *any* input is in the form of punched cards.
2. IBM 7550 *Card Punch* is necessary if on-line punching of output is desired.
3. IBM 7400 *Printer* is necessary if on-line printing of output is desired.
4. *As many as four additional tape units* may be used if the input is to be on more than one tape unit and/or if more than one program is to be compiled during a single machine run.

The specific machine requirements for each type of run which may be made using the Compiler Systems Tape are included in the 7070/7074 Data Processing System Bulletin "IBM 7070/7074 Compiler Systems: Operating Procedure," form J28-6105. Detailed operating instructions for each run, as well as the control cards necessary, are also included in this bulletin.

7070/7074 Autocoder is one of the powerful programming languages of today. Such languages have steadily evolved from languages requiring highly codified instructions closely related to the arithmetic capacities of the machine.

The first step in the evolution was the introduction of symbolic programming systems, such as 7070/7074 Basic Autocoder, in which a symbolic instruction is

written in place of each machine-language instruction. The programmer is thus able to code more easily and with greater meaning and the chance of errors is materially reduced.

The introduction of macro-instructions was a further step towards simplifying programming and reducing the time required to write a program. A macro-instruction is a symbolically-coded instruction which will produce a group of machine-language instructions. Two types of macro-instructions exist, substitution-type and generator-type.

The macro-instructions handled by 7070/7074 Four-Tape Autocoder are the substitution-type. The processor completes a "skeleton" routine by inserting parameters from the macro statement operand into this routine. The macro statement is then replaced by the completed routine. Four-Tape Autocoder also accepts symbolic machine instructions.

Autocoder, in addition to accepting symbolic machine instructions, handles generator-type macro-instructions. The task of compiling the proper sequence of instructions for the given macro statement is performed by the appropriate macro generator in the Library portion of the Compiler Systems Tape. In general, the operand of each macro-instruction must conform to a basic format. However, Autocoder also accepts certain macro-instructions with operands whose formats have not been pre-established "symbol by symbol." For example, the operand of ARITH or of LOGIC contains an "expression," the value of which is to be computed or the truth or falsity of which is to be determined. The macro generators scan the macro statements and compile the corresponding sequence of symbolic machine instructions; the operand need only conform to the rules of format which have been established and must not exceed certain specified lengths. Numerous illustrations of source program macro-instructions and their corresponding series of generated machine instructions are included with the macro-instructions under "Imperative Statements."

Macro-instructions to handle input/output operations are included in both the Autocoder and Four-Tape Autocoder systems. With certain restrictions, the source-program input/output statements are written the same for both Autocoders. However, the statements are handled differently by the processors of the two systems. In Four-Tape Autocoder, the processor substitutes the completed "skeleton" routine for the macro statement; in Autocoder, the macro-instruction is processed by means of a macro generator. A description of the input/output macro-instructions is included in this manual; a full discussion can be found in the 7070 Data Processing System Bulletin "IBM 7070 Input/Output Control System," form J28-6033-1. Input/output operations are estimated to constitute an average of 40% of a program; the ability to handle these operations by means of macro-instructions represents a substantial gain in programming simplicity and efficiency.

The Autocoder language may be extended by adding new macro-instructions. Appropriate macro generators may be added to the system to process the newly-created macro-instructions. The necessary generator is written in the Autocoder language according to the instructions presented in the 7070 Data Processing System Bulletin "Additions to the IBM 7070 Autocoder; Writing Macro Generators for the IBM 7070 Autocoder," form J28-6053. The generator is then compiled and entered on a new Compiler Systems Tape and the corresponding macro-instruction may then be included in any program.

While macro-instructions provide the programmer with a set of powerful tools to solve problems without becoming enmeshed in the tedious details of analysis

and of storage assignment, the programmer may still exercise direct control over the minor details of his program, should it be necessary. Autocoder accepts and processes all symbolic machine instructions as well as macro-instructions; in fact, any program written in Basic Autocoder can be assembled without change by Autocoder. In addition, storage allocation can be specified by a set of control statements which have been provided.

In summary, the Autocoder system provides a number of advantages by introducing a powerful macro language. The programming and processing advantages are as follows:

1. Macro language eliminates the need for breaking down many frequently encountered tasks into a number of small steps by turning these tasks over to macro generators.
2. Full use is made of the symbolic programming devices of the system. The need for attention to the details of data flow, actual storage allocation, and decimal-point positioning is eliminated.
3. The programmer is allowed to write program steps that are meaningful in terms of the problem to be solved, rather than in terms of machine capacities.
4. The program may be easily broken into meaningful segments, allowing for greater flexibility and accuracy in reprogramming or recombining of program segments.
5. The program is made much more readable.
6. Programming speed is materially increased.
7. Programming is much easier to learn because fewer instructions must be written by the programmer and because the need for concern with many machine details is eliminated.
8. Macro language reduces programming errors by making use of tested macro generators rather than many individual machine instructions.
9. Errors in the input statement are detected by macro generators themselves; a message is issued indicating the location of the error and, generally, the type of error committed.
10. Programs are largely independent of individual machine characteristics and therefore are easier to transfer from one system to another.





# Coding Sheet

A programmer, coding a program for the 7070 or 7074 in Autocoder, writes all Autocoder statements on the 7070 Autocoder Coding Sheet, form X28-6417-2 (see Figure 1). The coding sheet indicates by column numbers the input card format for both the Basic Autocoder and the Autocoder Systems. Each line of the coding sheet is punched into the indicated columns of a corresponding IBM 7070 Autocoder Input Card, electro A18265 (see Figure 2). An explanation of the purpose of each heading on the coding sheet is given below.

## Heading Line

The heading line consists of the spaces labeled "Program," "Programmed By," and "Date." The information entered in these spaces is for identification of the program and is not to be punched into input cards.

## Page Number (Columns 1-2)

A two-character page number sequences the coding sheets. Any alphanumeric characters may be used, providing they can be read into or can be translated on output by the input/output equipment of the 7070 or 7074 system used to process the Autocoder source program. (This applies in general to the usage of alphanumeric characters in all Autocoder statements.) The standard collating sequence should be followed in sequencing the pages. Alphanumeric characters which are not acceptable to various input/output equipment and the collating sequence may be found on page 9 of the IBM Reference Manual "7070 Data Processing System," form A22-7003-2.

FIGURE 2

### **Line (Columns 3-5)**

The first twenty-five lines on each sheet are prenumbered 01 through 25. Also provided are five unnumbered lines at the bottom of the sheet that may be used for additional lines or for inserts. Since provision is made for a three-character line number for the sequencing of the coding entries on the sheet, inserts may be readily made through the use of the optional third character. For example, inserts between lines 10 and 11 may be made by writing 101, 102 and 103 or 10A, 10B, and 10C, etc. in columns 3-5 of the unnumbered lines and placing the resulting cards after the card for line 10. Any alphameric characters may be used for all three characters of the line number so long as the standard collating sequence is followed in sequencing the lines.

The sequence of the cards entered into the processor will be checked by the page and line numbers punched in the source program deck. Any variation from the collating sequence will be noted in the warning and error message area of the program listing produced during a compilation. However, source-language input will be compiled in the order encountered.

### **Label (Columns 6-15)**

The label column is used to represent the location of data or instructions in the machine. Only instructions or data which will be referred to elsewhere in the program need have a label. In all other cases, the label column is blank. A label may be a symbolic location or an actual address. Each symbolic label must be unique, i.e., it may *not* appear more than once in the label column.

### **Operation (Columns 16-20)**

The operation column contains either a macro statement or the mnemonic representation of the machine operation to be performed. In certain cases, the column may be left blank. Actual machine operation codes are never used. The Autocoder symbolic machine instructions and macro-instructions are composed of from one to five alphameric characters, and are written left-justified in the operation column. Operation codes are categorized as "declarative," "control," and "imperative" statements. A description of these three types of statements is contained on pages 20, 87 and 103, respectively. If an invalid operation is used, a NOP will be generated.

### **Operand (Columns 21-75)**

The operand contains the actual or symbolic address of the information which is to be acted on by a particular command, or other parameters to be utilized by a macro generator. The operand may contain 55 columns of information as input to the Autocoder processor as opposed to the 40 columns for the Basic Autocoder processor. When field definition, address adjustment, or indexing is used in conjunction with the address, it is included in the operand. The operand may contain the actual data to be operated on by an instruction, referred to as a "literal." It may also be used to specify index words, electronic switches, channels, units, channel and unit, arm and file, inquiry and unit record synchronizers, latch numbers, and alteration switches.

## Identification (Columns 76-80)

Program identification is punched into columns 76-80 of all cards in the source program deck. The identification which appears on the *first* card of the source program deck will be punched in the identification field of each card of the object program deck and printed on each page of the output program listing by the Autocoder processor. A means is thus provided for relating all output to the proper symbolic source-language input.

Any alphameric characters acceptable to the 7070/7074 input/output equipment may be used in the identification entry. Alphabetic and special characters will print as such on the program listing. However, only the second digit of each double-digit representation of these characters will be punched in the identification field of the condensed cards of the object program decks since these cards must be numerical. For example, "A" will be punched as "1," "8" will be punched as "6," etc. If the first card of the source program deck has a blank identification field, the identification on the program listing will be blank and zeros will be punched in the cards of the object program deck, the double-digit representation of a blank being 00.

## Remarks and Comments

Remarks and comments may be included for description. They will appear in the symbolic output but will produce no entry in the object program deck and consequently will not affect the operation of the program.

Remarks may be included anywhere in the operand, provided they are separated by at least two blank spaces from the operand of the instruction. As noted on page 60, an @ symbol may not appear in remarks which are on the same line as an alphameric constant. Otherwise, remarks may include any acceptable characters.

Comments cards allow the programmer to insert complete lines of descriptive information in the program. A comments card is identified by placing an asterisk in column 6 of the label column. Any part of the label, operation, or operand columns may be used for the description. Comments cards are useful as descriptive headings for various sections of a program, such as operating instructions, or where the operand column of an instruction does not allow enough room for necessary remarks.

Remarks and comments may be used in a program as follows

Line	Label	Operation	OPERAND	Basic
3	56	1516	2021 25 30 35 40 45	50
01	*NET, PAY, CALCULATION,			
02		ZAL	GROSS, PUT, GROSS, PAY, IN, ACC, 1	
03		S1,	TAX, DEDUCT, INCOME, TAX,	
04		S1,	FICA, DEDUCT, FICA	
05				

# Parameters

## Address Types

The following types of addresses may appear in the label and/or operand fields of Autocoder statements: blank, actual, symbolic and literal. A description of these address types and the rules governing their usage follow.

### Blank

The label column may be blank if the corresponding entry is not referred to elsewhere in the program. A blank operand is valid for certain machine operations (see Appendix D) and in the following control and declarative statements:

- |                      |                        |
|----------------------|------------------------|
| 1. ORIGIN Control    | 6. XRELEASE Control    |
| 2. LITORIGIN Control | 7. SRELEASE Control    |
| 3. END Control       | 8. DC (Header Line)    |
| 4. XRESERVE Control  | 9. DLINE (Header Line) |
| 5. SRESERVE Control  |                        |

The effect of a blank operand varies for each of the control and declarative statements in which it may appear. These are explained fully under the discussion of the respective statements. If a Priority Release command has a blank operand, the processor will insert 0097. In all other machine operations where a blank operand is valid, the processor will insert 0000. If a blank operand is invalid, an error message will be produced.

### Actual

An actual address may be from one to four digits, written left-justified on the coding sheet. Leading zeros may be omitted. An actual address is valid only in the label and operand of symbolic machine instructions and in the operand of certain macro-instructions and control statements.

In certain instructions, an actual operand may only assume values within a limited range, e.g., a value of 1 through 4 in an alteration switch instruction and a value of 1 to 30 in an electronic switch instruction.

An actual address in the label column of a symbolic machine instruction will cause the instruction to be assigned that actual location. The contents of the location assignment counter being used will not be changed and the actual location, with the exception of index word addresses, will not be reserved. Hence, the programmer should be extremely cautious when using an actual label.

The following example illustrates the use of actual addresses in symbolic machine instructions:

Line	Label	Operation	OPERAND
3	56	15 16	20 21 25 30 35 40 45
0 1		N O P	3 2
0 2		.	
0 3	0	B	0 3 0 8
0 4		.	
0 5		H P	1
0 6			

### Symbolic

A symbolic address is valid in the operand of most statements and in the label column of all but the control statements. A symbolic address may contain from

one to ten characters with the following restrictions: the first, or leftmost, position must be a letter; the remaining characters may be letters or numbers (no special characters); blanks may not appear within the symbolic address. A symbolic address placed in the label column of a declarative or imperative statement is automatically associated with an actual storage location assigned by the processor's location counter. Further reference to a certain instruction may be made or operations on a particular data field may be performed by writing the symbolic name assigned to the instruction or data field in the operand of an imperative, declarative, or control statement.

Consider the following example:

Line	Label	Operation	OPERAND				
3	5 6	15 16 20 21	25	30	35	40	45
0 1	R, O U T I N E	Z A 1	F I E L D				
0 2							
0 3							
0 4		B H	R O U T I N E				
0 5							

If FIELD is a one-word field assigned to location 3000, the first entry above will result in the assembled instruction +1300093000. If the assignment counter had been at 5000 when the first entry was encountered, the label ROUTINE would have been associated with location 5000. Thus the entry in line 04 above would result in the assembled instruction -4000095000.

The asterisk, \*, is a special symbol which is valid in the operand only. If an actual address has not been written in the label column, the processor will assign the location of the instruction being processed, contained in the location assignment counter, to \*. For example, if the instruction

Line	Label	Operation	OPERAND				
3	5 6	15 16 20 21	25	30	35	40	45
0 1		Z A 1	*				
0 2							

has been assigned to location 4440, then the \* in the operand of that instruction will also be assigned location 4440; i.e., the assembled machine instruction will be +1300094440.

If an actual label *has* been entered in the label column and an \* in the operand, the processor will assign the actual address to \*. Thus, the instruction

Line	Label	Operation	OPERAND				
3	5 6	15 16 20 21	25	30	35	40	45
0 1	3 2 4	Z A 1	*				
0 2							

would result in the assembled machine instruction +1300090324.

Use of the asterisk address will reduce the number of symbols required in the label column. Unless there is a note to the contrary, the special symbol, \*, may be used as an operand address wherever this manual indicates that a symbolic address is valid.

## Literal

A literal is the actual data to be operated on by an instruction. The literal is valid only in the operand of an imperative statement, and its appearance causes the processor to assign a storage location to the literal. In order to conserve storage, literals are packed into words. The processor assigns field definers to the literal and incorporates the storage address and field definers into the instruction being processed. Once a literal has been stored, it will be re-used each time it is referenced again, except when the Litorigin Control statement is used (see page 88). Address adjustment and indexing (both discussed in later sections) should *not* be appended to literals. Literals should *not* be used for temporary storage, i.e., the operands of Store instructions should not be literals.

Certain imperative instructions (e.g., Priority Control, Index Word Load, Tape Write) operate on full words and do not permit field control. If a literal which is less than ten digits, or an adcon, is used as the operand of any of these instructions, it will be converted to ten-digit form by being right-justified in a word. Thus, if -50 is entered as the address portion of an Index Word Load instruction (XL), the actual constant stored will be -0000000050. Also, assume that ATABLE is the label of an instruction or data occupying location 2000. If the adcon -ATABLE is entered as the address portion of an Index Word Load and Interchange command (XLIN), the representation will be -0000002000, and, at object time, the contents of the specified index word will be replaced by -0020000000. In general, however, a literal should be written in the exact form in which it is to be used with leading zeros included where necessary.

Four principle classes of literals are permissible: automatic-decimal numbers, floating-decimal numbers, address constants (adcons), and alphameric constants.

### Automatic-Decimal Numbers

A literal having either of the following sets of characteristics is included in the automatic decimal class.

1. A signed number, one to *twenty* digits in length, which is referred to by a *macro-instruction*. A decimal point *may be* included in the literal.
2. A signed number, one to *ten* digits in length, which is referred to by a *symbolic machine instruction*. A decimal point *may not be* included in the literal.

An automatic-decimal number, referred to by a macro-instruction and described by the first set of characteristics above, will be examined by the Autocoder processor for decimal-point inclusion. A decimal point indicates the magnitude of the number according to ordinary usage and the desired decimal-point placement; it is neither stored in the constant nor saved with it. Using Autocoder macro generators, the processor generates instructions for shifting, decimal alignment, bridging words, and handling field definers. If the decimal point falls to the right of the rightmost digit of the number, it may be omitted and the number will be considered an integer. The following are examples of automatic-decimal literals which might appear in the operand of a macro-instruction:

```
+1234567890123 . 4567890
- . 12345
-1 . 23
+1
```

An automatic-decimal number described by the second set of characteristics above is actually a subset of one described by the first set of characteristics. Because the number is referred to by a symbolic machine instruction, the literal is restricted to ten digits in length and all instructions for shifting and decimal alignment must be written by the programmer.

An automatic-decimal number may also be defined by a DC subsequent entry (see page 55).

### Floating-Decimal Numbers

Floating-decimal numbers are permitted as a literal entry in macro-instructions only. A number may be expressed in the form

$$(\pm n) (10^{\pm m})$$

where  $n$  is an integer or decimal number of not more than eight digits and  $m$  is a one-digit or two-digit exponent. Floating-decimal numbers are related to this form and are entered according to the format  $\pm nF^{\pm m}$ . Thus, the number  $-.12345678 \times 10^3$  would be represented by  $-.12345678F+3$ . If the sign preceding  $m$  is omitted,  $m$  is considered to be positive. The exponent  $m$  may be omitted if equal to 0, provided  $\pm nF$  is not followed by another literal entry. The Auto-coder processor will consider the signs, the value of  $n$ , and the value of  $m$ , and generate a standard 7070 floating-decimal word. The following are examples of floating-decimal numbers which might appear in the operand of a macro-instruction:

-3.4F  
+34.567893F-2  
-31.92F+7  
-29567.1F-3  
+12546F+15

Additional examples of floating-decimal numbers may be found in the IBM Reference Manual "7070 Data Processing System," form A22-7003-2.

A floating-decimal number may also be defined by a DC subsequent entry as described on page 57.

### Address Constants (Adcons)

An adcon is a special-purpose numerical literal used to produce a four-digit constant whose value is the actual address assigned to a symbolic address. The address constant is treated in the same manner as a numerical literal. An adcon is entered in the form  $\pm$ SYMBOL. The following example illustrates the use of an adcon:

Line	Label	Operation	OPERAND								
3	5	6	15	16	20	21	25	30	35	40	45
01	ANY LABEL	S	FIELD								
02		.									
03		.									
04		ZA 1	+ ANY LABEL								
05											

If the symbolic address, ANYLABEL, is assigned the actual location, 2000, the use of the adcon, +ANYLABEL, in the operand of line 04 will cause the processor to produce a constant of +2000. If the address constant, +2000, is assigned to location 4576 with field definers (6, 9), the entry on line 04 would result in the assembled instruction +1300694576. The execution of this instruction at object program time would place the number, +0000002000, the actual address assigned to ANYLABEL, in accumulator 1.

An adcon may also be defined by a DC subsequent entry (see page 58).

### Alphameric Literal

An alphameric literal consists of alphameric characters preceded and followed by the @ character. A literal in the operand of a symbolic machine instruc-



tion may not be more than five characters in length and, in the operand of a macro-instruction, may not be more than 120 characters in length. All characters between the initial @ character and the second @ will be converted to double-digit form and assigned to core storage locations. The sign of each word used to contain the characters will be alpha. If the constant produced by a literal is more than one word in length, but *not* a multiple of five alphameric characters, the double-digit representation of alphameric blanks will be generated in the unused low-order positions of the last word of the constant, i.e., the constant will *not* be packed.

Alphameric literals may *not* include the @ character. Alphameric literals may, however, contain any other character (except for the record mark) which may be read by the input device used.

The following are examples of alphameric literals which might appear in the operand of macro-instructions or symbolic machine instructions:

```
@ABCD@
@A100X@
@12345@
@.□+*$@
@-/,%#@
```

The following literal would be valid in the operand of a macro-instruction only:

```
@THIS LITERAL IS LONGER THAN 5 CHARACTERS@
```

Alphameric constants, including those containing the @ character, may be defined by a DC subsequent entry (see page 60).

## Field Definition

Field definition may be written immediately following the operand address of symbolic machine instructions. The field definers are the digit positions of a ten-digit word numbered 0 to 9. The format for entering field definers is as follows:

Line	Label	Operation	OPERAND	Basic Autocod
3	56	1516	2021 25 30 35 40 45	50 55
01			ADDRESS( STARTPOSITION, ENDPOSITION )	
02				

The starting and ending positions are the field definers. When operating on a single digit, the comma and ending digit position may be omitted. Field definers may be omitted entirely when operating on a whole word or a field defined by a declarative statement (see pages 22 through 81).

Field definition may not be used in the operand of a macro-instruction. However, both macro-instructions and symbolic machine instructions may refer to the label of a Define Area (DA), Define Constant (DC), or Define Line (DLINE) subsequent entry, each of which may specify field definition.

Autocoder symbolic machine operation codes which permit field definition to be associated with the address are indicated in Appendix D. Field definition

may be used with both actual and symbolic addresses as illustrated in the following examples:

Line 3	Label 56	Operation 1516 2021	OPERAND				
			25	30	35	40	45
0 1		S 1	*( 8 , 9 )				
0 2		A 4	MANNO( 0 , 2 )				
0 3		S T 2	1 0 2 4 ( 2 , 8 )				
0 4		Z A 1	ALPHA( 5 )				
0 5		C 1	MANNO( 0 , 9 )				
0 6		A S 1	1 0 2 4 ( 0 , 9 )				
0 7		Z A 1	-SYMBOL( 2 , 3 )				
0 8							

In lines 05 and 06 above, field definers (0, 9) could have been omitted since the instructions will be operating on the entire word.

When used with literals, field definition will be relative to the literal itself. Thus, the adcon `-SYMBOL(2, 3)` refers to the third and fourth digits of the location assigned to `SYMBOL`. If location 9876 has been assigned to `SYMBOL`, `-SYMBOL(2, 3)` would be equivalent to `-76`. (See "Relative Field Definition" on page 41.)

## Address Adjustment

Address adjustment allows the programmer to refer to an entry which is a given number of locations preceding or following a symbolic address. Address adjustment is permitted with adcons and with all symbolic addresses, except the single-address operand of a `DRDW` statement (see page 78). It should not be used with actual addresses or other literals.

With *symbolic machine instructions*, address adjustment is indicated by writing a plus or minus sign followed by one to four digits immediately after the symbolic address and following any field definers.

Address adjustment with symbolic machine instructions is written as follows:

Line 3	Label 56	Operation 1516 2021	OPERAND				
			25	30	35	40	45
0 1			MANNO+5				
0 2			MANNO-2				
0 3			MANNO( 0 , 2 )+5				
0 4			*-1				
0 5			*( 0 , 4 )+10				
0 6							

If location 2150 has been assigned to the symbolic address `MANNO`, then 2155 will be assigned to `MANNO+5`, and 2148 will be assigned to `MANNO-2`. The entry `MANNO(0, 2)+5` refers to the first three digit positions of location 2155. Similarly, if 2196 is the location of an instruction containing `*-1` as an operand, then location 2195 will be assigned to this operand.

With *macro-instructions*, address adjustment is indicated by writing a plus or minus sign followed by one to four digits after the symbolic address. The sign

and digits must be enclosed by parentheses. The left parenthesis must be in the column immediately following the last character of the symbolic address being modified, except when the address adjustment goes to a continuation card (see "Continuation Cards").

Address adjustment with macro-instructions is written as follows:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1			LIST(+3)
0 2			LIST(-10)
0 3			

If location 1560 has been assigned to the symbolic address LIST, then 1563 will be assigned to LIST (+3), and 1550 will be assigned to LIST (-10).

Address adjustment of an *adcon* will cause a special function. The value of the address will be modified before the constant is created. Address adjustment of an *adcon* in a symbolic machine instruction is written as follows:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1	ANYLABEL	ZA1	+ANYLABEL+1
0 2			

If ANYLABEL is assigned to location 2000, the above entry will cause the creation of a literal (or *adcon*) value of 2001 to be stored for reference. In a macro-instruction, the format in the operand would be +ANYLABEL(+1).

The programmer should be careful when using address adjustment since it may become a source of error when a program is modified. For example, inserts and deletions of program entries could change addresses in such a way that  $*+10$  should now be  $*+9$ . It should also be noted that since it is not known how many machine instructions will be generated in place of a macro-instruction, address adjustment must not be used on a symbolic label in amounts that would carry the address into or across a macro-instruction.

## Index Words

### Indexing

The use of an index word in an instruction for the purpose of indexing will cause the indexing portion of the index word to be added algebraically to the address portion of the instruction and this new address is used for the operation. Indexing of symbolic and actual addresses may be specified in the operand of all imperative statements and Branch Control and End Control statements. Literals should not be indexed.

With the *Branch Control and End Control statements* and with *symbolic machine instructions*, the address of an index word follows the operand address, after field definers and address adjustment, if any, and is always preceded by a plus sign. An index word may be written symbolically or as the actual one- or two-digit number (1-99). When a symbolic name is used, the processor will

automatically assign an actual index word address. When the actual number is used, the format is Xn. The X indicates that an index word address rather than address adjustment follows the plus sign. The n is the actual one- or two-digit number of the word. When used for other than indexing purposes, the form Xn will be considered a symbol. Indexing with Branch Control and End Control statements and symbolic machine instructions is written as follows:

Line 3	Label 56	Operation 1516	OPERAND				
			2021	25	30	35	40
01			MANN0,+X2				
02			MANN0-15+X2				
03			MANN0(0,4)+15+LOOP				
04			2344+X2				
05			2344+LOOP				
06			2344(0,5)+LOOP				
07							

Autocoder will interpret +X2 as index word 2 and +LOOP as the symbolic designation of an index word. Note that blanks are not permitted within the address modification.

With *macro-instructions*, the address of an index word follows the operand address and is enclosed by parentheses. The left parenthesis must be in the column immediately following the last character of the address being modified, except when the indexing goes to a continuation card (see "Continuation Cards"). Actual index word numbers are not preceded by an X and the actual or symbolic index word is not signed. In macro-instructions, the absence of a plus or minus sign preceding an index word distinguishes the index word from a signed, one- or two-digit address adjustment. If both address adjustment and indexing are used, one set of parentheses must enclose them both and the indexing must *precede* the address adjustment. Indexing with macro-instructions is written as follows:

Line 3	Label 56	Operation 1516	OPERAND				
			2021	25	30	35	40
01			TABLE,(ROW)				
02			LIST,(34)				
03			TABLE,(ROW-10)				
04			LIST,(34+17)				
05							

Autocoder will interpret 34 as index word 34 and row as the symbolic designation of an index word. Note again that blanks are not permitted within the address modification and, also, that field definition may not be used with macro-instructions.

### Uses Other Than Indexing

Index words may be specified in symbolic machine instructions as the first entry in the operand of index word commands such as Index Word Load (XL) and Index Word Load and Interchange (XLIN). They may also be used in commands such as Record Gather (RG) and Record Scatter (RS). The specification of an *index word* in a symbolic machine instruction does not prevent the use of *indexing*. The index word may be written in actual or symbolic form. When written in actual form, however, only the one- or two-digit address of the index word should be written; Xn would be interpreted as the symbolic designation of an index word.

The following are examples of the use of index words for other than indexing purposes:

Line	Label	Operation	OPERAND
3	5,6	15,16	20,21 25 30 35 40 45
0.1		XL	2, CONSTANT
0.2		XL	X3, +0,000,010,100
0.3		RG	BASE, SCATAREA
0.4			

BASE and X3 will be interpreted as the symbolic designation of an index word and 2 as index word 2.

## Electronic Switches

Electronic switches may be referred to by a symbolic name or by their one- or two-digit actual number (1-30). Symbolic names may be assigned to an actual switch number by use of the declarative statement `equ` (see page 82). As explained later, symbolic references to electronic switches will be assigned to an actual address during compilation.

- Unlike index words (see page 18), electronic switches will *not* be reserved if the location assigned to an imperative statement or if any location defined by a declarative statement falls within the range 0101-0103.

Instructions referring to electronic switches are written as follows:

Line	Label	Operation	OPERAND
3	5,6	15,16	20,21 25 30 35 40 45
0.1		BES	19, COMPUTE
0.2		BSF	END, LOOP
0.3		ESN	28
0.4			

END will be interpreted as the symbolic designation of an electronic switch and 19 and 28 as electronic switches 19 and 28.

## Input/Output Unit and Alteration Switch Designations

The following items may be specified in actual or symbolic form in the operands of those instructions which refer to the particular items: channel, unit, combined channel and unit, combined arm and file, unit record synchronizers, inquiry synchronizers, and alteration switches. The declarative operation `equ` is used to equate symbolic names to item numbers (see page 85).

## **Continuation Cards**

Certain Autocoder statements make provision for more parameters than may be contained in the operand (columns 21-75) of a single line on the Autocoder coding sheet. When this is the case, the appropriate section of this manual will indicate that "Continuation Cards" may be used. Thus, when specifically permitted, the operand of a given line on the Autocoder coding sheet may be continued in the operand of from one to four additional lines which immediately follow.

The label and operation columns must be blank and the continuation of the operand must begin in column 21; i.e., it must be left-justified in the operand column of the coding sheet. The operand need not extend across the entire operand column of either the header card or continuation cards but may end with the comma following any parameter. Remarks may appear to the right of the last parameter on each card provided they are separated from the operand by at least two blank spaces.

Illustration of the use of continuation cards are included throughout the examples illustrating the various statements.

If a continuation card follows a statement that does not permit continuation cards, the compiler will generate a NOP and issue an error message. Additional restrictions regarding the use of continuation cards with macro-instructions appear on page 106.

## Reservation of Index Words and Electronic Switches

The assignment of actual addresses to symbolic index word and electronic switch names occurs in Phase III of the Autocoder processor. The initial availability of index words and electronic switches is determined by a table which is included in the Compiler Systems Tape. This table originally indicates that index words 1 through 96 and electronic switches 1 through 30 are available for assignment to symbolic references; index words 97 through 99 are not available. The initial setting of this table may be altered, however, as described in the 7070/7074 Data Processing System Bulletin "IBM 7070/7074 Compiler System: Operating Procedure," form J28-6105.

During the first pass of Phase III, references to the *actual* addresses of index words and electronic switches are collected and the availability table is updated. At the end of this pass, the table indicates which index words and electronic switches are not available for assignment to symbolic references.

Both index words and electronic switches may have been made unavailable before the start of assignment in one of the following ways:

1. The initial setting of the availability table indicated that the index word or electronic switch was not available for assignment.
2. The one- two-digit number of the index word or electronic switch was used in the operand of a symbolic machine instruction to specify indexing or as a parameter which is *always* an index word or electronic switch, e.g.,

Line	Label	Operation	OPERAND															
			3	5	6	15	16	20	21	25	30	35	40	45				
0.1		B,L,X								5								
0.2		E,S,N								1	6							
0.3																		

3. The one- or two-digit number of the index word or electronic switch was used in the operand of an EQU statement, e.g.,

Line	Label	Operation	OPERAND															
			3	5	6	15	16	20	21	25	30	35	40	45				
0.1	N,A,M,E	E,Q,U								3								
0.2																		

When the index words or electronic switches are reserved because of actual usage in the statements described above, the position or order of the statements within the program is not important; any such reference will make the index word or electronic switch unavailable at the end of this pass.

During the assignment pass of Phase III, index words and electronic switches are reserved as they are encountered during assignment. Index words and electronic switches may be reserved in the following ways. The first two methods apply to both index words and electronic switches; the third applies only to index words.

1. During the assignment pass, each instruction is examined for reference to the symbolic name of an index word or electronic switch. When such a reference is found, an actual address is assigned and the availability table is changed so that the assigned index word or switch is no longer available for later assignment.
2. If the one- or two-digit address of an index word or electronic switch is used or is included in the operand of an `XRESERVE` or `SRESERVE` statement (see page 99), the corresponding index word or electronic switch is reserved.
3. If a statement has been assigned an address in the index word area
  - a. by means of an actual label or
  - b. by means of an `ORIGIN` statement which refers to an *actual address*

the corresponding index word will be reserved. These entries should normally appear at the beginning of the program or immediately following each `LITORIGIN` statement. Otherwise, symbolic names may have previously been assigned to these same index words. (This method does not apply to electronic switches.)

The preceding methods allow efficient use of index words and electronic switches during a sectionalized or multi-phase program, particularly when used in conjunction with the `LITORIGIN` statement. Extreme caution should be used, however, to avoid the conflicting usage of an index word or electronic switch which may result from the assignment of more than one name or function to the same address.

If the symbolic name or actual address of an index word or electronic switch appears or is included in the operand of an `XRELEASE` or `SRELEASE` statement (see page 101), the specified index word or electronic switch will again be made available, regardless of the method by which it was reserved. It will not, however, be used for symbolic assignment until all other index words or electronic switches have been assigned for the first time.

If, at any time during the assignment pass, the compiler finds that there are no more index words available for assignment, the warning message “NO MORE INDEX WORDS AVAILABLE” will be placed in the object program listing, the table will be altered to show that index words 1 through 96 are available, and the assignment will continue as before. If the compiler finds that there are no more electronic switches available for assignment, the warning message “NO MORE ELECTRONIC SWITCHES AVAILABLE” will be placed in the object program listing, the table will be altered to show that electronic switches 1 through 30 are available, and assignment will continue as before. The resultant conflicting usage of index words or electronic switches may be avoided by reducing the number of symbolic names used, e.g., through the proper use of the `EQU`, `XRELEASE`, or `SRELEASE` statements.

As noted in Appendix C, index words 97 through 99 are *never* available for assignment to symbolic names by the compiler; also, index words 93 through 96 may have been made unavailable for assignment.



## Declarative Statements

Autocoder declarative statements provide the processor with the necessary information to complete the imperative operations properly. Declarative statements are never executed in the object program and should be separated from the program instruction area, placed preferably at its beginning or end. Otherwise, special care must be taken to branch around them so that the program will not attempt to execute something in a data area as an instruction. If the compiler does encounter such statements, a warning message will be issued. 7070/7074 Autocoder includes the following declarative statements: `DA` (Define Area), `DC` (Define Constant), `DRDW` (Define Record Definition Word), `DSW` (Define Switch), `DLINE` (Define Line), `EQU` (Equate), `CODE`, `DTF` (Define Tape File), `DIOCS` (Define Input/Output Control System), and `DUF` (Descriptive Entry for Unit Records). `DA`, `DC`, `DTF`, and `DLINE` require more than one entry.

The `DA` statement is used to name and define the positions and length of fields within an area. The `DC` statement is used to name and enter constants into the object program. Since the 7070 and 7074 make use of record definition words (`RDWS`) to read, write, move, and otherwise examine blocks of storage, the `DA` and `DC` statements provide the option of generating `RDWS` automatically. When so instructed, Autocoder will generate one or more `RDWS` and assign them successive locations immediately preceding the area(s) with which they are to be associated. An `RDW` will be of the form  $\pm 00xxxxyyyy$ , where `xxxx` is the starting location of the area and `yyyy` is its ending location. These addresses are calculated automatically by the processor.

In some cases, it may be more advantageous to assign locations to `RDWS` associated with `DA` and `DC` areas in some other part of storage, i.e., not immediately preceding the `DA` or `DC` areas. The `DRDW` statement may be used for this purpose. The `DRDW` statement may also be used to generate an `RDW` defining *any* area specified by the programmer.

As many as ten digital switches may be named and provided by the `DSW` statement for consideration by the `SETSW` and `LOGIC` macro-instructions. Each switch occupies one digit position in a word, can be set `ON` or `OFF`, and is considered as logically equivalent to an electronic switch. It cannot, however, be referred to by electronic switch commands, e.g., `ESN`, `BSN`, etc. An individual switch or the entire set of switches in a word may be tested or altered as desired.

Through use of the `DLINE` statement, a means is provided for specifying both the editing of fields to be inserted in a print line area and the layout of the area itself. The area may include constant information supplied by the programmer. The area may also be provided with additional data during the running of the object program by means of `EDMOV` or `MOVE` macro-instructions.

The declarative statement `EQU` permits the programmer to equate symbolic names to actual index words, electronic switches, arm and file numbers, tape channel and unit numbers, alteration switches, etc., and to equate a symbol to another symbol or to an actual address.

The `DIOCS`, `DTF`, and `DUF` statements are used when required by the Input/Output Control System. `DIOCS` is used to select the major methods of processing to be used, and to name the index words used by `IOCS`. Each tape file must be described

by Tape File Specifications, produced by `DTFS`. In addition to information related to the file and its records, the File Specifications contain subroutine locations and the location of tape label information. A `DUF` entry must be supplied for every unit record file describing the type of file and the unit record equipment to be used. The `DUF` also supplies the locations of subroutines written by the user that are unique to the file.

A full description of the `DIOCS`, `DTF`, and `DUF` statements is contained in the 7070 Data Processing System Bulletin "IBM 7070 Input/Output Control System," form J28-6033-1. Brief descriptions of these three declarative statements and detailed descriptions of the formats and functions of each of the other 7070/7074 Auto-coder declarative statements follow below.

# DIOCS—Define Input/Output Control System

When the Input/Output Control System is to be used in a program, a DIOCS statement must be used to select the major methods of processing to be used. This statement also allows the naming of the index words used by IOCS.

## Source Program Format

The basic format of the DIOCS statement is as follows:

Line	Label	Operation	OPERAND					Basic Autocoder	Autocoder								
3	56	15	16	20	21	25	30	35	40	45	50	55	60	65	70		
0 1	ANY LABEL	D I O C S	I O C S I X F	,	I O C S I X G	,	I O C S I X H	,	C H A N n	,	O P E N n	,	E O R n	,	C H P T	,	I G E N n
0 2																	

ANYLABEL is any symbolic label; it may be omitted. The entry DIOCS must be written exactly as shown.

The first item in the operand, IOCSIXF, is used to specify the first IOCS index word for programs using tape files. This item may be a symbolic name or an actual one-digit or two-digit index word address in the range 3-94. If the first item in the operand is omitted, the symbolic name IOCSIXF will be assigned. When an actual index word or a symbolic address is specified, Autocoder will equate the name IOCSIXF to it.

The second item in the operand, IOCSIXG, is used to specify the second IOCS index word for programs using tape files. This item may be a symbolic name or an actual one-digit or two-digit index word address in the range 3-94. If the second item in the operand is omitted, the symbolic name IOCSIXG will be assigned. When an actual index word or a symbolic address is specified, Autocoder will equate IOCSIXG to it.

The third item in the operand, IOCSIXH, is used to specify an IOCS index word for programs using unit record files. This item may be a symbolic name or an actual one-digit or two-digit index word address in the range 3-94. If the third item in the operand is omitted, the symbolic name IOCSIXH will be assigned. When an actual index word or a symbolic address is specified, Autocoder will equate IOCSIXH to it.

The fourth item in the operand, CHANN, is used to specify the number of the highest tape channel to be used by the program. Thus, the programmer would write CHAN1, CHAN2, CHAN3, or CHAN4 to show that the program was to use channel 1, channels 1 and 2, channels 1, 2, and 3, or channels 1, 2, 3, and 4, respectively.

The fifth item in the operand, OPENN, is used to specify the method of handling the OPEN macro-instruction. The value of n may be 1-6. If 1 or 5 is used, the special procedure discussed under "Use of OPEN1" in the "IBM 7070 Input/Output Control System" bulletin should be followed.

1. If OPEN1 is entered in the operand, the OPEN subroutine will not be preserved in storage after it is used; other subroutines will be loaded into the locations

used by the OPEN subroutine. Thus, all tape files must be opened at the same time. The DTFs and File Schedulers must have been loaded into storage before this routine is loaded and executed.

2. If OPEN2 is used, the OPEN subroutine will be retained in storage for use whenever needed.
3. If OPEN3 is entered in the operand, the OPEN subroutine will be written on the tape provided for checkpoint records and read into storage whenever needed. The storage locations required for the OPEN subroutine will be used for other subroutines during the time the OPEN subroutine is on tape. When OPEN3 is specified, EOR1 and CHPT must also be specified.
4. If OPEN4 is used, the OPEN subroutine will be retained in storage for use whenever needed as for OPEN2, except that Form 3 and Form 4 records cannot be processed and three input/output areas cannot be used for one file. The OPEN4 subroutine will occupy fewer storage locations than the OPEN2 subroutine.
5. If the three-area rotating method is used, either OPEN5 or OPEN6 must be specified. OPEN5 and OPEN6 contain provisions for the three-area rotating system; otherwise they are the same as OPEN1 and OPEN2, respectively. The DTFs and File Schedulers must have been loaded into storage before OPEN5 is loaded and executed.

The sixth item in the operand, EORN, is used to specify whether tape labels are to be processed in the End-of-Reel subroutines. The value of n in EORN may be either 1 or 2.

1. The use of EOR1 in the operand specifies that the reading or writing of tape labels is to be determined by the LABELINF entry in the appropriate DTF for each input and output file. EOR1 is required when OPEN3 and/or CHPT is specified.
2. If EOR2 is used in the operand, none of the input tapes may have labels nor will any labels be written on output tapes.

The seventh item in the operand, CHPT, is used to specify whether checkpoint records are to be written. If CHPT is entered in the operand, checkpoint records will be written under the control of the DCHPT statement; EOR1 must also be specified. If CHPT is omitted from the operand, no checkpoint records may be written.

The eighth item in the operand, IGENN, is used to specify the use of SPOOL programs and illegal double-digit character checking in the IOCS tape error subroutine. The value of the final n in IGENN may be 1, 2, 3, or 4.

1. Entering IGEN1 in the operand indicates that a SPOOL program(s) may operate with this main program and that the tape error subroutine is to check for illegal double-digit characters.
2. Entering IGEN2 in the operand indicates that a SPOOL program(s) may operate with this main program but that the tape error subroutine will not check for illegal double-digit characters.
3. Entering IGEN3 in the operand indicates that a SPOOL program will never be run with this main program but that the tape error subroutine is to check for illegal double-digit characters.
4. Entering IGEN4 in the operand indicates that no SPOOL program can be run with this main program nor will the tape error routine check for illegal double-digit characters.

## Processing Techniques

The DIOCS statement may appear only once in a program. As noted previously, the first three items in the operand may be used to specify names for certain index words used by the Input/Output Control System. Each of the remaining entries is used by Autocoder to determine the version of the corresponding iocs subroutine that will be produced. Since the generated material is located at the point where the DIOCS statement is encountered, the programmer should not include the DIOCS statement within a series of imperative statements.

As noted previously, IOCSIXF, IOCSIXG, IOCSIXH, and/or CHPT may be omitted from the operand. In each case, separating commas must be written. Thus, if all four of the above items are omitted, the operand would appear as follows:

Line	Label	Operation	OPERAND							
3	56	1516	2021	25	30	35	40	45		
0 1	ANY LABEL	D I O C S	,	, C H A N n	,	, O P E N n	,	, E O R n	,	, I G E N n
0 2										

If a DIOCS operand does not contain eight items (including omitted items indicated by a separating comma) the iocs subroutines will be generated as though the following DIOCS statement had been entered:

Line	Label	Operation	OPERAND					Base				
3	56	1516	2021	25	30	35	40	45	50			
0 1	ANY LABEL	D I O C S	,	, C H A N 2	,	, O P E N 2	,	, E O R 1	,	, C H P T	,	, I G E N 1
0 2												

If any of the individual items in a DIOCS statement are invalid, the processor will generate one of the following subroutines corresponding to the invalid item: CHAN2, OPEN2, EOR1, CHPT, or IGEN1. Thus, an incorrect entry as the fourth item in the operand would cause the processor to assume CHAN2, etc.

## Error and Warning Messages

CHAN ENTRY INVALID. TWO CHANNELS ASSUMED

If the CHANN entry does not specify CHAN1 or CHAN2, two channel schedulers will be generated.

CHPT ENTRY INVALID. CHPT INCLUDED

If the CHPT entry is neither blank, nor "CHPT," this message will be produced. The checkpoint routine will be generated.

CHPT REQUIRES EOR1. EOR1 GENERATED

If the CHPT entry is not blank and if the EORN entry specifies EOR2, this message will be produced. EOR1 will be generated.

EOR ENTRY INVALID. EOR1 ASSUMED

If the EORN entry does not specify EOR1 or EOR2, EOR1 will be generated.

IGEN ENTRY INVALID. IGEN1 ASSUMED

If the IGENn entry is not in the range IGEN1 to IGEN4, IGEN1 will be generated.

IMPROPER OPERAND. CHAN2, OPEN2, EOR1, CHPT, IGEN1 ASSUMED

If the operand of the DIOCS entry does not contain 8 parameters, this message will be produced. An IOCS package consisting of the items named in the message will be generated.

IOCSIXF(G, H) ENTRY OUT OF RANGE, IGNORED

If any of the index words in the DIOCS entry is specified as actual, and if it does not lie in the range 3-94, this message will be produced. The specified index word will be ignored.

OPEN ENTRY INVALID. OPEN2 ASSUMED

If the OPENn entry is not in the range OPEN1 to OPEN6, OPEN2 will be generated.

OPEN3 REQUIRES CHPT. OPEN2 GENERATED

If the OPENn entry specifies OPEN3 and if the CHPT entry is blank, this message will be produced. OPEN2 will be generated.

## DTF—Define Tape File

Each input and output tape file required by a program must be described by a set of File Specifications in a DTF statement. The DTF statement consists of a header line and 35 subsequent entries which describe the file, subroutine locations supplied by the user, and the location of tape label information. These lines are entered on the 7070 File Specifications Coding Sheet shown in Figure 3.

Each DTF statement causes the generation of a nine-word File Specifications Table and a Tape File Scheduler which are used by the Input/Output Control System in conjunction with input and output tape file operations. In addition, the Autocoder compiler used the information furnished by the DTF in order to generate the proper instructions when the name of the file is used in a macro-instruction.

### Source Program Format

#### DTF Header Line

The basic format for the DTF header line is as follows:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
01	ANYLABEL	DTF	FILENAME					
02								

ANYLABEL may be any symbolic label; if it is omitted, the processor will generate a label. The entry DTF must be written exactly as shown. The operand must contain the name of the tape file defined. This name will be used in the operand of macro-instructions which refer to this tape file.

#### DTF Subsequent Entries

Although the line numbers for each DTF subsequent entry need *not* be the same as those which appear in the File Specifications Coding Sheet, a card must appear for each line, in the order shown. If a DTF does not consist of exactly 36 entries, the processor will discontinue compilation. A label for each line is not necessary unless the entry will be referred to by the program; the labels of other entries may be blank. The operand for each entry must be as specified in the bulletin, "IBM 7070 Input/Output Control System."

### Processing Techniques

Since the Input/Output Control System requires that all File Specifications Tables appear in consecutive locations, the programmer must insure that all DTFs are entered together at the desired point in the program. The processor will determine the location of the first DTF encountered and make this location available to the appropriate input/output routines. A nine-word constant will be assembled at the point where each DTF is encountered.

In addition, the processor will generate a corresponding Tape File Scheduler. This scheduler will be located with literals, adcons and other material generated out-of-line; placement may be determined by use of a LITORIGIN statement.





When the processor finds that all consecutive DTFs have been assembled, it will generate a NOP to signal the end of the series of Tape File Specifications Tables. If a separate DTF is encountered at some later point in the program, it will not be included in the series of File Specifications Tables and a warning message will be issued. The DTF will, in general, be assembled properly, but a corresponding File Scheduler will *not* be generated. Therefore, the ninth word of the assembled DTF will not contain the customary address of a File Scheduler. In addition, the assembled DTF will not be available to the Input/Output Control System during the object program unless it is moved to a point within the previously mentioned series of File Specifications Tables and the address of an appropriate File Scheduler inserted in the ninth word.

The product of the blocking factor (BLOCKING, line 11) and the tape input/output areas (TIOAREAS, line 17) in the DTF for a particular tape file determines the number of areas to be specified in the header line of the DA entry which defines the area for the tape file records. (See "N (Area Number)," page 32.)

When the name of an input or output file is to be used in Autocoder macro-instructions which are not a part of the Input/Output Control System, the following requirements must be considered:

1. The name of the file used in the operand of the DTF entry must be identical to the name of the DA entry which defines the areas for the input or output records.
2. The operand of the DA entry which defines the areas must specify an implicit index word and that index word must be the same as the operand of the INDXRDA entry in the appropriate DTF.

## DUF—Descriptive Entry for Unit Records

When unit record files are to be handled by the Input/Output Control System, a DUF entry for each file must be supplied which describes the type of file and the unit record equipment to be used. The DUF entry also supplies the locations of subroutines written by the user that are unique to the file.

### Source Program Format

The basic format of the DUF statement is as follows:

Line	Label	Operation	OPERAND	Basic
3	56	1516	2021 25 30 35 40 45	50
0 1	ANY LABEL	DUF	FILENAME, FILETYPE, CARDSYNC,	
0 2			LISTADDR, INDXWORD, EOFADDRS,	
0 3			ERRADDRS	
0 4				

ANYLABEL is any symbolic label; it may be omitted. The entry DUF must be written exactly as shown.

The first item in the operand, FILENAME, is the name of the unit record file to be described by the DUF entry. This name will be used in the operand of macro-instructions which refer to this unit record file.

The second item in the operand, FILETYPE, is a one-digit number from 1 through 4 to specify the type of unit record file and the operating conditions.

1. A 1 indicates that the unit record file is an input file and that the input unit used by the file will never be shared with a SPOOL program.
2. A 2 indicates that the unit record file is an output file and that the output unit used by the file will never be shared with a SPOOL program.
3. A 3 indicates that the unit record file is an input file and that the input unit may be shared with a SPOOL program.
4. A 4 indicates that the unit record file is an output file and that the output unit may be shared with a SPOOL program.

The third item in the operand, CARDSYNC, is a one-digit number to specify the synchronizer to be used for the unit record file. A 4 must be used when an input file is to be read through an IBM 7501 Console Card Reader.

The fourth item in the operand, LISTADDR, is the address, either actual or symbolic, of the RDW(s) for the unit record area generated by a DA or DRDW statement.

The fifth item in the operand, INDXWORD, is either a two-digit number from 03 through 94 or a symbolic name which specifies the index word to be associated with the unit record file. The indexing portion (positions 2 through 5) of the index word will contain the location of the first word of the current unit record.

The sixth item in the operand, EOFADDRS, is an optional address that may be either actual or symbolic. The address specifies the location of a card reader end-of-file

routine or a printer carriage tape channel 9 routine to be entered if either condition occurs. When this item is omitted and a card reader end-of-file condition occurs, the program will come to a programmed halt if no `SPOOL` program is to be run in conjunction with it. If a `SPOOL` program is to operate in conjunction with the program, the program will enter a loop to permit the `SPOOL` program to continue. When this item is omitted, a channel 9 condition occurring during the printing of a unit record file will have no effect and the program will continue normally.

The seventh item in the operand, `ERRADDRS`, is an optional address that may be either actual or symbolic. The address specifies the location of an error routine which will be entered when an error occurs during the execution of a macroinstruction which refers to the file named in the first item of the `DUF` entry. If this item is omitted and an error occurs, the error record will first be typed on the console typewriter. When the file is an output file, the error record will also be punched or printed and processing then resumes automatically. When the file is an input file, the operator may correct the error card immediately and read it in again or he may depress the Start key and resume processing if immediate correction is not required.

## Processing Techniques

The `DUF` entries are entered with the source program when the program is assembled. They should be positioned in the same manner as declarative entries.

As noted earlier, the sixth and seventh items in the operand are optional. Whenever the sixth item is omitted and the seventh item is included, the omission of the sixth item must be indicated by a comma; i.e., two commas will appear between the fifth and seventh items. When both items are omitted, commas *are not* required; only the first five items need appear in the operand.

## Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified.

### ACTUAL ADDRESS NOT ALLOWED

This message is issued if `FILENAME` is not a symbolic name. The entire `DUF` statement is ignored.

### AN ELEMENT OF THE OPERAND STARTS ILLEGALLY

This message is issued in connection with the `DUF` if either of the following conditions occur:

1. The character following a comma is not alphabetic, numerical, blank or comma.
2. `FILENAME` begins with a blank.

Processing will occur on all parameters (if any) which appear before the parameter in error; the remainder will be ignored.

### LABEL SHOULD BE BLANK

This message is issued if any continuation cards for the `DUF` operand have an entry in their label column. The entry is processed properly but the label is ignored; if referenced by the program, it will be undefined.

OPERAND OUTSIDE OF ALLOWABLE RANGE

This message is issued if:

1. FILETYPE is not 1, 2, 3, or 4.
2. CARDSYNC is not 1, 2, 3, or 4.
3. INDXWORD, if actual, is greater than 99.

The processor assumes that the value is 1 and normal processing continues.

STMT SHOULD OR SEEMS TO BE ENDED BUT CARDS REMAIN

This message is issued if there are continuation cards remaining but the current card includes the seventh parameter or ends with a blank. The remaining cards are ignored.

SYMBOLIC ADDRESS NOT ALLOWED

Issued if FILETYPE or CARDSYNC begins with an alphabetic character. Preceding parameter(s) will be processed; the remaining parameters are ignored.

## DA—Define Area

The declarative statement, DA, may be used to define and reserve any portion of storage. An area may be reserved for use as an input, output, or work area, or contiguous areas may be reserved to contain a number of records, all of which are identical in format. The DA statement instructs the processor concerning the positions, lengths, and names of fields which make up the record area(s) being defined, as well as the characteristics of the data which is to occupy each field. Such characteristics include format, implied decimal-point position, and double- or single-digit representation. Locations, field definers, and, if specified, implicit indexing are assigned by the processor to enable the programmer to refer to the fields by name in imperative statement operands. Thus, the programmer need not be concerned with the actual locations of the fields within storage. It should be remembered that the DA statement does not provide the data for the field it defines, but only reserves the space which the data is to occupy. The data itself must be brought into the defined area in core storage from some external source such as cards or magnetic tape, or from other locations in storage.

A DA statement consists of a header line and one or more subsequent entries. The DA header line is used to *initiate* the reservation of a portion of storage. The header line specifies the number of identical record areas to be reserved. An indication to generate rows corresponding to the areas may be specified. Relative addressing and implicit indexing, which facilitate referencing a field within a record area, may also be specified in the DA header line. The subsequent entries define the fields within the record area, specify the amount of storage to be reserved, and describe the data which will appear in each field.

### Source Program Format

#### DA Header Line

The basic format of the DA header line is as follows:

Line	Label	Operation	OPERAND					Basic Ad
			25	30	35	40	45	
0 1	ANY LABEL	DA	N	, ± RDW	, ADDRESS ± ADDR	, DJ	+ INDEX	WORD
0 2								

ANYLABEL is any symbolic label; it may be omitted. The operation code, DA, must be written exactly as shown. With the exception of the first entry, N, the items in the operand are optional. An explanation of the entries in the operand follows.

*N (Area Number).* N, the first entry in the operand of the DA header line, may *not* be omitted. N is replaced by the number of identical record areas to be reserved by the DA operation, the format of which is defined by the programmer. For example, suppose that records of identical format are to be read into storage in blocks of ten. The programmer might enter a DA header line with N equal to 10, followed by subsequent entries which specify the starting and ending posi-

tions, characteristics, and names of the fields composing one record. If all other items in the operand were omitted, the DA header line would appear as follows:

Line	Label	Operation	OPERAND						
3	5/6	15/16	20/21	25	30	35	40	45	
0 1	A N Y L A B E L	D A	1 0						
0 2									

$N$  may be any unsigned number from 1 to 9999; it is limited only by the size of storage. (Frequently,  $N$  will equal 1, as in a work area or a unit record input or output area.) The number of storage words to be reserved for the entire DA area (excluding RDWs, if any) will be  $N$  times the number of words reserved for the one record defined by the subsequent entries. The maximum number of words which may be reserved for an entire DA area is also limited only by the size of storage.

If the DA is to reserve contiguous areas of storage that are to contain a number of input or output records having identical format, then  $N$  must equal the product of the blocking factor and the tape input/output areas indicated in the DTF for the file (see page 28). If the blocking factor for a tape file is 10 and the programmer only wishes to use one area of storage for input or output, then  $N$  should equal 10. However, if the programmer wishes to use two consecutive areas of storage for input or output, then  $N$  must equal 20 and if he wishes to use three consecutive areas, then  $N$  must equal 30.

*RDW (Record Definition Words).* Record definition words are required by the 7070 and the 7074 for reading in and writing out data and for moving blocks of data within storage. If "RDW" is written in the operand of the DA header line, the processor will automatically generate  $N$  RDWs associated with the  $N$  defined areas. These RDWs will be assigned  $N$  locations immediately preceding the first word of the first area defined.

If RDW is not preceded by a sign, all generated RDWs will be plus except the last, which will be minus. If a + or - sign precedes RDW, all generated RDWs will be given the indicated sign.

If RDW is not written in the operand, RDWs will not be generated; the first word reserved by the DA statement will be the first word of the first record area.

For example, suppose that an area is to contain a group of four 10-word records and that the following DA header line is entered, the fields within a record being defined by subsequent entries:

Line	Label	Operation	OPERAND						
3	5/6	15/16	20/21	25	30	35	40	45	
0 1	R E C O R D	D A	4	R D W					
0 2									

Assume that Autocoder's location assignment counter contains 1000 when the above DA statement is encountered. The processor will generate the following

RDWs associated with the four record areas and assign them to locations immediately preceding the first record area:

<i>Symbol</i>	<i>Location</i>	<i>RDW</i>
RECORD	1000	+0010041013
	1001	+0010141023
	1002	+0010241033
	1003	-0010341043

Note that the first RDW, located at 1000, defines the first 10-word record area beginning at 1004; the second RDW, located at 1001, defines the second 10-word record area beginning at 1014; etc. If RDW had been omitted from the DA header line, the first word of the first 10-word record would have been located at 1000; no RDWs would have been generated.

RDWs may also be defined elsewhere by using the DRDW statement.

*ADDRESS (Relative Addressing)*. The fields defined in subsequent entries may be assigned addresses relative to (i.e., beginning with) an ADDRESS specified in the operand of the DA header line. ADDRESS may be written in actual, \*, or symbolic form. When a relative address is specified, it will usually be 0.

Assume that a relative address of 0 is specified, as in the following example:

Line	Label	Operation	OPERAND
3	5/6	15/16	20/21 25 30 35 40 45
0 1	ANY LABEL	DA	2 0 , R D W , 0
0 2			

In this case, any field which occupies the first, second, or third word, etc., of the area beyond the generated RDWs will be assigned relative addresses of 0000, 0001, 0002, etc., respectively. Thus, if FIELD C is the name of a field occupying the third word of each record, the instruction

Line	Label	Operation	OPERAND
3	5/6	15/16	20/21 25 30 35 40 45
0 1		Z A, 1	F I E L D C ( 0 , 4 ) + X, 1, 1
0 2			

would be translated in the object program as +1311040002. Since twenty identical areas have been reserved and since the third word of each of the twenty areas is addressed FIELD C and assigned the relative address 0002, one of the twenty areas referred to by an instruction must be specified. The desired area may be indicated by indexing the individual instruction (as illustrated) or by implicit indexing (as explained in the next section). Referring to the illustration, the indexing portion of index word 11 must have been programmed to contain the starting location of the area to be processed. Then, when indexing takes place, the indexing portion of index word 11 is added algebraically to the relative address of FIELD C, 0002, resulting in the actual address of FIELD C.

The presence of a relative address in the DA header line does not affect the reservation of storage. The generated RDWs define the actual record areas reserved; they are also unaffected by the absence or presence of relative addressing. For example, assume that the location assignment counter is at 1000 when the following DA header line is encountered:

Line	Label	Operation	OPERAND						
3	56	1516	20	21	25	30	35	40	45
0 1	ANY LABEL	DA	2	0	,	RDW	,	0	
0 2									

Twenty RDWs associated with the twenty record areas will be assigned locations 1000 through 1019. If the subsequent entries define a record that is ten words in length, the first RDW generated at 1000 would be +0010201029 and the first word of the first record would be assigned to 1020, the second to 1021, etc.

If relative addressing is specified, but RDWs are not to be generated, commas must appear in two consecutive columns, as in the following example:

Line	Label	Operation	OPERAND						
3	56	1516	20	21	25	30	35	40	45
0 1	ANY LABEL	DA	2	0	,	,	ADDRESS	+	INDEX
0 2									

*ADDRADJ (Address Adjustment)*. Address adjustment is permitted with a relative address in symbolic or \* form. However, when a relative address is specified, it will usually be 0 and address adjustment will not appear.

Address adjustment, when used with an \* or symbolic address, will appear as in the following examples:

Line	Label	Operation	OPERAND						
3	56	1516	20	21	25	30	35	40	45
0 1	AREA NAME	DA	3	,	RDW	,	ADDRESS	-	1
0 2									
0 3									
0 4									
0 5	RECORD	DA	1	0	,	,	*	+	1
0 6									



*INDEXWORD* (*Implicit Indexing*). An index word, either symbolic or actual, may be specified in a DA header line as in the following examples:

Line	Label	Operation	OPERAND
3	56	15 16	20 21 25 30 35 40 45
0 1	A R E A N A M E	D A	1 5 , R D W , A D D R E S S + I N D E X W O R D
0 2			
0 3			
0 4			
0 5	I N P U T F I L E	D A	1 3 , , A D D R E S S - 1 + I N D E X W O R D
0 6			
0 7			
0 8			
0 9	O U T P U T F I L E	D A	2 0 , , A D D R E S S + X 1 5
1 0			
1 1			
1 2			
1 3	C H A N G E F I L E	D A	8 , R D W , A D D R E S S + 1 0 + X 7
1 4			

In these examples, *INDEXWORD* is the symbolic name of an index word; X15 and X7 are the actual addresses of index words. Note that a plus sign must always precede the name of the index word and that actual numbers of index words must be preceded by "+X." The specification of indexing follows address adjustment, if present, or the relative address itself, if address adjustment is not used.

The naming of an index word in the header line has no effect upon the storage reservation and location assignments produced by the DA. It does facilitate writing instructions which reference fields in one of the records within the defined area. When the compiler encounters an instruction referencing such fields, the instruction will be examined for the presence of indexing. If indexing already appears in the instruction, it will be left unchanged. If indexing does not appear in the instruction, however, the address of the index word specified in the DA header line will be inserted in the index word positions of the instruction. Thus, if it is known when coding a DA that the record defined by it is to be processed by indexing, the indexing may be caused by making the single notation in the DA header line, rather than by supplying the indexing in every instruction which acts upon the record.

Consider an example in which the compiler encounters the following DA header line:

Line	Label	Operation	OPERAND
3	56	15 16	20 21 25 30 35 40 45
0 1	A R E A N A M E	D A	1 0 0 , , 0 + X 2 5
0 2			

Each field in a record will be assigned a location relative to 0. If FIELD2 is the label of a field occupying the second word of each record in the area, the instruction

Line	Label	Operation	OPERAND						
3	5/6	15/16	20/21	25	30	35	40	45	
0.1		Z A 1	F I E L D 2						
0.2									

will result in an assembled machine instruction of +1325090001. Notice that the 6-9 portion of the instruction is 0001, but that this address is to be indexed by index word 25. The indexing portion of index word 25 must have been programmed to contain the starting location of the area to be processed. Then, by indexing the instruction by index word 25, the indexing portion of the index word is added algebraically to the relative address of FIELD2, 0001, resulting in the actual address of FIELD2.

If, however, the field had been referred to by an instruction which already included indexing, such as

Line	Label	Operation	OPERAND						
3	5/6	15/16	20/21	25	30	35	40	45	
0.1		Z A 1	F I E L D 2 + X 2 9						
0.2									

the assembled instruction would have been +1329090001, thus ignoring the specification of index word 25 in the DA header line.

As demonstrated by the previous examples, the use of a DA header line to specify a relative address and implicit indexing is most valuable in processing blocked tape records. If the fields of a record are defined as relative to 0 and the rows associated with the records are successively loaded into the index word specified in the DA header line, instructions to act upon the fields of the record may be written as if no indexing was required. Autocoder will insert the address of the index word into the index word positions of all such instructions.

For another example, assume that an area is to contain a group of four 10-word records. Assume also that it is desired to perform the following operation on three fields named FIELD A, FIELD B and FIELD C, which occupy the first, second and third words, respectively, of each record: subtract the contents of FIELD B from the contents of FIELD A and store the result in FIELD C. This could be accomplished by the following coding:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
01	I N A R E A	D A	4 , + R D W , 0 + X 2
02		.	
03			(Subsequent entries define
04			the fields within a record.)
05		.	
06		X L	3 , + 0 0 0 0 0 0 0 0 0 3
07	L D R D W	X L	2 , I N A R E A + X 3
08		Z A 1	F I E L D A
09		S 1	F I E L D B
10		S T 1	F I E L D C
11		B I X	3 , L D R D W
12			

Assume that Autocoder's assignment counter contains 1000 when the DA statement is encountered, 2000 when the first instruction of the routine is encountered, and the literal +0000000003 has been assigned location 4432. The corresponding assembled machine-language instructions would be as shown below:

<i>Location</i>	<i>Assembled Instruction</i>
2000	+4500034432
2001	+4503021000
2002	+1302090000
2003	-1402090001
2004	+1202090002
2005	+4900032001

### DA Subsequent Entries

Subsequent entries under a DA header line are used to name fields, to indicate their starting and ending positions within the record area, and to specify their format. The subsequent entries, taken collectively, indicate a logical assemblage of words, the number of which is multiplied by the number *N* in the DA header line to determine the amount of storage (excluding *rdws*, if any) to be reserved by the DA.

The label column of a DA subsequent entry may include any symbolic name or may be left blank (except as noted for the special operation, *CODE*, which is explained on page 50). The operation column is left blank, except when used for the *CODE* entry.

*Field Position and Length.* The position and length of a field within the record area is indicated in the operand by writing the starting digit position, a comma, and the ending digit position. If the area is considered to be a consecutive string of digits, the first being 0, the position of any digit is identical to its placement in the string. Thus, a field occupying the entire first word of a record area would be defined as follows:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	A,R,E,A,N,A,M,E	D,A	1					
0 2	F,I,E,L,D,A		0	,	9			
0 3								

In the above example, 0 is the digit position of the first character and 9 is the digit position of the tenth and final character in FIELD A.

A field occupying the entire second word of a record area would be defined as follows:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	A,R,E,A,N,A,M,E	D,A	1					
0 2	F,I,E,L,D,B		1	0	,	1	9	
0 3								

A field occupying the last two digits of the third word of a record area would be defined as follows:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	A,R,E,A,N,A,M,E	D,A	1					
0 2	F,I,E,L,D,C		2	8	,	2	9	
0 3								

A digit position, therefore, is a combination of a word number (consisting of one, two, three, or four numerical characters) followed immediately by a digit number (always one numerical character). The first word of the area is referred to as word 0, the second as word 1, the third as word 2, etc.; digits within a word are numbered 0 through 9, left to right, in the normal fashion. Thus, digit position 49 would refer to the low-order digit in the fifth word of a record area; digit position 9990 would refer to the high-order digit in the thousandth word of a record area.

It is not necessary to enter leading zeros; the following operands would have exactly the same effect:

12,25  
00012,00025

A one-digit field may be entered by writing the single digit position only. Thus, the following operands would have exactly the same effect:

19,19  
19

In either case, the processor would reserve the low-order digit position of the second word of the record area.

The following example provides additional illustrations of entries which define the length and position of fields within an area.

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	R E C O R D	D A	1	R D W				
0 2	N A M E O N E		0 0	1 6				
0 3	N A M E T W O		1 7	1 9				
0 4	N A M E T H R E E		2 0	2 9				
0 5	N A M E F O U R		5 3	5 7				
0 6	N A M E F I V E		5 8					
0 7			7 9					

An eight-word record is defined (digits 0-79). The first field, NAMEONE, occupies a total of 17 digit positions beginning with the high-order position of the first word of the record. NAME TWO and NAME THREE are fields of lengths 3 and 10 digits, respectively.

A gap of 23 digits then occurs, followed by NAMEFOUR, a five-digit field. This is followed by NAMEFIVE, a one-digit field. (Since the starting and ending digit positions are identical, the one number, 58, suffices.)

A subsequent entry is always required when a field is to be named for reference by the object program. Other fields in a record need not be defined by a subsequent entry except for an entry to terminate the record. The field *must* appear but need not have a label. Thus, in the example above, the one-digit field in digit position 79 is not named, but it is necessary in order to establish that eight words are occupied by the record. Any digit in the eighth word (73, for example) would cause the DA area to be comprised of eight full words, since a record is never terminated in the middle of a word.

Subsequent entries may be entered in any order. However, fields are normally entered in ascending storage-position order for convenience and accuracy.

The net effect of the following coding would be identical to that of the preceding example, despite the order of the subsequent entries:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	R E C O R D	D A	1	R D W				
0 2			7 9					
0 3	N A M E O N E		0 0	1 6				
0 4	N A M E F I V E		5 8					
0 5	N A M E T W O		1 7	1 9				
0 6	N A M E F O U R		5 3	5 7				
0 7	N A M E T H R E E		2 0	2 9				
0 8								

Subfields are entered in the same DA which defines the fields. Any desired breakdown of the fields may be indicated by writing the starting and ending digit positions of subfields which overlap, fall within, or bridge other fields. Again, the order is irrelevant. The following example illustrates the coding of subfields:

Line	Label	Operation	OPERAND
3 56	1516	2021	25 30 35 40 45
0 1	CALENDAR	DA	1, RDW
0 2	DATE		7 4, 7 9
0 3	DAY		7 6, 7 7
0 4	YEAR		7 8, 7 9
0 5	MONTH		7 4, 7 5
0 6	DECADE		7 8
0 7			

In this example, DATE will reference the entire date field; the subsections of the date field may be referred to as MONTH, DAY, or YEAR. In addition, DECADE references a one-digit field within YEAR and also, therefore, within DATE.

*Relative Field Definition.* The primary function of the DA declarative operation is to instruct the processor as to the positions and lengths of fields within an area, thereby allowing the processor to assign storage locations and field definers automatically when a field is referred to by name in the operand of an instruction. Sometimes it may be desirable to refer to a portion of a field which has been defined by a declarative operation. For the convenience of the programmer, this is done relative to the field itself, so that it is not necessary to remember the actual digit positions of the field.

To illustrate, assume that the name CUSTNO is a field of seven digits, defined under a DA as follows:

Line	Label	Operation	OPERAND
3 56	1516	2021	25 30 35 40 45
0 1	INPUTAREA	DA	1
0 2	CUSTNO		0 3, 0 9
0 3			

A field might occupy only the low-order digit positions of a word, as does CUSTNO above. However, positions within the field itself are numbered starting with 0. If the programmer wishes to refer to the three high-order positions of CUSTNO, he would refer to positions 0, 1 and 2 by placing the entry CUSTNO (0, 2) in the operand. During assembly, the processor will convert the field definers relative to CUSTNO to field definers which refer to actual digit positions of a word. Thus, during assembly, the processor will convert CUSTNO (0, 2) to the actual digit positions 3, 4 and 5 of the word of which CUSTNO is a portion. If the storage location of this word is 1001, symbolic instructions using CUSTNO and their assembled equivalents might be as follows:

<i>Symbolic Instruction</i>	<i>Assembled Instruction</i>
ZAL CUSTNO(0, 2)	+1300351001
ZAL CUSTNO(4, 5)	+1300781001
ZAL CUSTNO	+1300391001

CUSTNO might instead be defined as a full word, as follows:

Line	Label	Operation	OPERAND
3	5	15	16 20 21 25 30 35 40 45
0 1	I,N,P,U,T,A,R,E,A	D,A	4
0 2	C,U,S,T,N,O		0 0, 0 9
0 3			

In this case, any field definition following CUSTNO would be both relative and actual. If the storage location of this word is 1001, symbolic instructions using CUSTNO and their assembled equivalents might be as follows:

<i>Symbolic Instruction</i>	<i>Assembled Instruction</i>
ZAL CUSTNO(0, 2)	+1300021001
ZAL CUSTNO(4, 5)	+1300451001
ZAL CUSTNO	+1300091001

In the example on page 41 illustrating how subfields are entered, the subfield DAY would be equivalent to DATE(2, 3).

*Format Indicators.* In addition to defining the position and length of a field, a DA subsequent entry may also specify format and characteristics for that field. It will be noted that Autocoder macro-instructions are capable of acting upon fields which bridge words, which exceed ten digits in length, or whose decimal points are not aligned. The macro generators will consider the formats for the fields concerned; they will then generate instructions to align decimal points while performing arithmetic and compare operations, and will convert a field from one format to another format.

Five different types of fields may be designated by writing the proper format indicators immediately after the ending digit position of the field. The following indicators may be used to indicate to the processor the format of the data that will occupy the field at object program time.

<i>Format Indicator</i>	<i>Meaning</i>
A	The field is to contain numerical data to be treated as an automatic-decimal number. Although automatic-decimal fields must not exceed 20 digits for arithmetic operations, longer integer fields may be defined, and are acceptable to some macro-instructions.
F	The field is to contain numerical data to be treated as a floating-decimal number. The field must be exactly 10 digits in length, beginning in digit position 0 of a word; it may <i>not</i> bridge words.

Format Indicator	Meaning
@	The field is to contain the double-digit representation of alphameric characters. The data is to be treated as neither an automatic-decimal nor a floating-decimal number. The field may be any desired length; the number of digits in the field must be evenly divisible by two because each character is represented by two digits. The field must start at an even-numbered digit position.
@A	The field is to contain the double-digit representation of data which is to be treated as an automatic-decimal number. The size of the field must not exceed 40 digits (20 characters). The number of digits in the field must be evenly divisible by two because each character is represented by two digits. The field must start at an even numbered digit position.
@F	The field is to contain the double-digit representation of data which is to be treated as a floating-decimal number. The size of the field must be exactly 20 digits, occupying two full words; it may not be split among three words.

If format indicators are omitted, the processor will assume that the field is to contain numerical data to be treated as a signed integer:

1. For fields less than or equal to 20 digits in length, this is equivalent to an automatic-decimal field.
2. For fields longer than 20 digits, see individual macro-instructions for the treatment of long integer fields.

Users are strongly urged, however, to furnish format indicators in DA subsequent entries.

*Decimal-Point Indicators.* In the case of automatic-decimal numbers, indicators may also be used to indicate where the decimal point would fall if it were included in the field. (It should be noted that the decimal point is thus implicit, or “understood”; it is not actually to be included in the field nor is a field position to be reserved for it.) The implicit positioning of the decimal point is done by placing an indicator of the form

a.b

immediately to the right of the A or @A format indicator. The “a” is replaced by the number of characters to the left of the implicit decimal point in the field described; the “b” is replaced by the number of characters to the right of the implicit decimal point. The sum of these indicators must not exceed 20 characters in length and must exactly equal the number of characters defined by the starting and ending positions for the field. For a numerical field, then, (a + b) must equal the number of digits defined for the field; for a double-digit representation, however, (a + b) must equal one half of the number of digits defined for the field.



For example, the operand

20,25A4.2

informs the processor that the field which occupies digit positions 20 through 25 should be operated upon as if it included six digits of numerical data and a decimal point fell between the four high-order digits and the two low-order digits.

However, the operand

20,25@A2.1

informs the processor that the field should be operated upon as if a decimal point fell between the two high-order *characters* and the low-order *character* of a *three-character* alphanumeric field. (It will be remembered that one character occupies two digit positions.)

When an integer is defined, the decimal point and the number of decimal positions may be omitted. Thus, the following entries would have the same effect:

20,29A10.0

20,29A10

If the A or @A format indicator is not followed by decimal-point indicators, the processor will assume that the field is to contain an integer. Users are strongly urged, however, to furnish format indicators and decimal-point indicators for automatic-decimal numbers in DA subsequent entries. Specification of these characteristics makes it possible for macro generators to perform diagnostic analysis of the source-program statements, and thus aid in minimizing problems in program testing.

Additional examples of the use of format and decimal indicators are given below:

Line	Label	Operation	OPERAND											
			3	5	6	15	16	20	21	25	30	35	40	45
01	INPUT	DA												
02	FIELD A													
03	FIELD B													
04	FIELD C													
05	FIELD D													
06	FIELD E													
07	FIELD F													
08	FIELD G													
09														

FIELD A reserves the six high-order positions of the second word of the record area for an automatic-decimal number which will have, at most, four integer places and two decimal places.

FIELD B reserves the six low-order positions of the third word of the record area for a maximum of three alphanumeric characters.

FIELD C reserves the entire fourth word of the record area for a floating-decimal number.

FIELD D reserves the last part of the fifth and the first part of the sixth words of the record area for an automatic-decimal number in double-digit form, having

at most three integer places and two decimal places. Ten digit positions must be reserved, however, because of the double-digit representation.

FIELD E reserves the seventh and eighth words of the record area for a floating-decimal number in double-digit representation.

FIELD F reserves the six high-order positions of the ninth word of the record area for an automatic-decimal integer.

FIELD G reserves the four low-order positions of the ninth word of the record area for the double-digit representation of an automatic-decimal integer.

Two or more symbolic names may be assigned to the same field and the same characteristics by listing both names as subsequent entries under the same DA and duplicating the operand desired, as in the following example:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1		DA	10, RDW, 0+INDEXWORD
0 2	FIELD,A		0,9A6,.4
0 3	A		0,9A6,.4
0 4	FIELD,B		20,29F
0 5	B		20,29F
0 6	FIELD,C		42,59@
0 7	C		42,59@
0 8			

In this example, A is made equivalent to FIELD A, B is made equivalent to FIELD B, and C is made equivalent to FIELD C. Any reference to A, B, or C will give precisely the same results as those which would result from reference to FIELD A, FIELD B, or FIELD C, respectively.

The contents of a given field may be treated as having different characteristics at various times by defining that field using more than one subsequent entry, as in the following example:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1		DA	10, RDW, 0+INDEXWORD
0 2	FIELD,A		0,9A6,.4
0 3	A		0,9A8,.2
0 4			

Each reference to the field must be to the label which is associated with the characteristics that apply to the contents of the field at that particular time in the object program.

## Processing Techniques

### Use of DA With Symbolic Machine Instructions

Although the programmer will usually find it more efficient to use macro-instructions to refer to DA header lines and subsequent entries, symbolic machine instructions may also be used for this purpose. When a symbolic machine instruc-

tion references a symbolic label, however, it will operate directly on a maximum of a single full word. Certain additional restrictions must also be considered; these are noted below.

*DA Header Line.* When a symbolic machine instruction references the label of a DA header line, it will operate on the entire first word reserved by the DA statement. Thus, if RDW generation is specified, the symbolic machine instruction will act on the entire word containing the first generated RDW. If RDW generation is *not* specified, the symbolic machine instruction will act on the entire first word of the first record defined. Symbolic machine instruction references to the label of a DA header line are not affected by the presence or absence of implicit indexing.

The above usages are illustrated by the examples on page 47.

*DA Subsequent Entry.* When a symbolic machine instruction references the label of a DA subsequent entry, it will act only on that portion of the field which falls within the word containing the starting digit position defined by the subsequent entry. The specification of implicit indexing will cause the symbolic machine instruction to address a word (or portion of a word) in the current record, as determined by the contents of the specified index word. If implicit indexing is not specified, the word (or portion of a word) addressed will be within the first record as defined by the subsequent entry.

The above usages are illustrated by the examples on page 47.

#### **Use of DA with Macro-instructions**

The Autocoder language is designed to allow record processing to be accomplished without regard to the number of records that may appear in one block of input or output. Thus, when a macro-instruction refers to the label of a DA header line, *one complete record* will be the maximum that is affected or considered. (It should be noted that, in some cases, only part of the record will be affected or considered.) When *N* is greater than 1, therefore, no single reference by a macro-instruction to the label of a DA header line will ever affect or consider the *entire* area reserved.

Regardless of the absence or presence of specified RDW generation in the DA header line, the effect of a macro-instruction reference to the label of a DA header line or subsequent entry will be as follows:

1. When a DA header line does *not* contain a relative address and implicit indexing,
  - (a) a macro-instruction reference to the label of the DA header line will cause the generation of coding which will affect or consider (as a maximum) the *first* record, as defined by the subsequent entries under the DA header line.
  - (b) a macro-instruction reference to the label of a subsequent entry will cause the generation of coding which will affect or consider (as a maximum) the named field within the *first* record, as defined by the subsequent entries under the DA header line.

This technique would normally be used when a work area is defined; i.e., *N* is 1 and the first record is, therefore, the only record defined.

PAGE AA	PROGRAM DA	7070 COMPILER SYSTEM VERSION QMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01 0002	*		EFFECTS OF SYMBOLIC MACHINE INSTRUCTION REFERENCE TO					
02 0003	*		LABEL OF DA HEADER LINE.					
03 0004	*							
04 0005	AREA1	DA	1				+0003250327	
05 0006			12,25			29	0326	0326
06 0007		ZA1	AREA1	00001			0328 +1300090325	
07 0008	*							
08 0009	AREA2	DA	2				+0003290334	
09 0010			12,25			29	0330	0330
10 0011		ZA1	AREA2	00002			0335 +1300090329	
11 0012	*							
12 0013	AREA3	DA	2,RDW				+0003360343	
13	X						0336 +0003380340	0336
14	X						0337 -0003410343	0337
15 0014			12,25			29	0339	0339
16 0015		ZA1	AREA3	00003			0344 +1300090336	
17 0016	*							
18 0017	AREA4	DA	2,RDW,0+INDEXWORD				+0003450352	
19	X						0345 +0003470349	0345
20	X						0346 -0003500352	0346
21 0018			12,25			29	0348	0001
22 0019		ZA1	AREA4	00004			0353 +1300090345	
23 0020	*							
24 0021	AREA5	DA	2,0+X15				+0003540359	
25 0022			12,25			29	0355	0001
26 0023		ZA1	AREA5	00005			0360 +1300090354	
27 0024	*							

PAGE AA	PROGRAM DA	7070 COMPILER SYSTEM VERSION QMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01 0028	*		EFFECTS OF SYMBOLIC MACHINE INSTRUCTION REFERENCE TO					
02 0029	*		LABEL OF DA SUBSEQUENT ENTRY.					
03 0030	*							
04 0031		DA	1				+0003250327	
05 0032	FIELDA		12,25			29	0326	0326
06 0033		ZA1	FIELDA	00001			0328 +1300290326	
07 0034	*							
08 0035		DA	2				+0003290334	
09 0036	FIELDB		12,25			29	0330	0330
10 0037		ZA1	FIELDB	00002			0335 +1300290330	
11 0038	*							
12 0039		DA	2,RDW				+0003360343	
13	X						0336 +0003380340	0336
14	X						0337 -0003410343	0337
15 0040	FIELDC		12,25			29	0339	0339
16 0041		ZA1	FIELDC	00003			0344 +1300290339	
17 0042	*							
18 0043		DA	2,RDW,0+INDEXWORD				+0003450352	
19	X						0345 +0003470349	0345
20	X						0346 -0003500352	0346
21 0044	FIELDDD		12,25			29	0348	0001
22 0045		ZA1	FIELDDD	00004			0353 +1301290001	
23 0046	*							
24 0047		DA	2,0+X15				+0003540359	
25 0048	FIELDDE		12,25			29	0355	0001
26 0049		ZA1	FIELDDE	00005			0360 +1315290001	
27 0050	*							

2. When a DA header line contains a relative address and implicit indexing,
  - (a) a macro-instruction reference to the label of the DA header line will cause the generation of coding which will affect or consider (as a maximum) the *current* record (as determined by the contents of the implicit index word at the time the instructions are executed).
  - (b) a macro-instruction reference to the label of a subsequent entry will cause the generation of a coding which will affect or consider (as a maximum) the named field within the *current* record (as determined by the contents of the implicit index word at the time the instructions are executed).

The generated instructions will include implicit indexing and references to the implicit index word wherever necessary.

Input/output macro-instructions may indirectly affect a record by operating on the RDWS defining that record; e.g., the PUTX macro affects records by interchanging RDWS. When the other macro-instructions (i.e., those not concerned with input/output) refer to the label of a DA header line which specifies RDW generation, however, they will not operate on the RDWS themselves. For example, when the ZERO macro is used in this way, it will not zero the RDWS; it will zero out one complete *record* area.

When input or output files are to be referred to by input/output macro-instructions, the following considerations apply:

1. Macro-instructions which are to operate on a *record* will reference the name of the file. This name must appear as the label of the DA header line *and* as the operand of the DTF for that particular file. Depending on the method of generating RDWS, this name may or may not be the same as the operand of the subsequent entry IORDWLST in that same DTF.
2. In the header line of the DA referenced, N and RDW specification will be prepared according to the procedures described in the 7070 Data Processing System Bulletin "IBM 7070 Input/Output Control System," form J28-6033-1.
3. A relative address of 0 will be specified in the DA header line.
4. Implicit indexing will be specified in the DA header line. The name of the index word specified must be the same as the operand of the subsequent entry INDXWRDA in the appropriate DTF.

When input or output files are to be referred to by the other Autocoder macro-instructions or by *both* the input/output and other macro-instructions, the following considerations apply:

1. If RDWS are generated by the DA header line, all macro-instructions can reference the name of the file. If RDWS are generated by a DRDW, the name of the file cannot be referenced; however, the area which the DRDW references may be used as an operand of any macro-instruction, where applicable.
2. In the header line of the DA referenced, N and RDW specification will be prepared according to the procedures described in the 7070 Data Processing System Bulletin "IBM 7070 Input/Output Control System," form J28-6033-1.
3. A relative address of 0 will be specified in the DA header line.
4. Implicit indexing will be specified in the DA header line. The name of the index word specified must be the same as the operand of the subsequent entry INDXWRDA in the appropriate DTF.

When a DA statement is used to define a work area, the following considerations will normally apply:

1. N will be 1.
2. RDW generation will be specified.
3. Neither a relative address nor implicit indexing will be specified.

### **Arrangement of Fields**

In general, the arrangement of fields defined by DA subsequent entries is determined by prior considerations, e.g., the record form and format. Thus, a tape record is usually arranged to take full advantage of the zero suppression of the IBM 7070/7074; i.e., numerical fields which most frequently contain leading zeros are each placed in the high-order positions of a word, but numerical fields which less frequently (or never) contain leading zeros are each placed in the low-order positions of a word. In the interests of tape capacity and speed, it is also desirable to hold the appearance(s) of the delta character to a minimum. In general, this means that all alphameric fields should appear together.

Since imperative statements frequently alter the sign of a word to reflect the sign of the last field entered, the programmer should insure that fields which occupy the same storage word are to be associated with the same sign. If this is not the case, inconsistent results may occur. For example, if alpha information is entered into a word in which another field contains numerical data, the sign of the word may become alpha; the numerical information would then become alpha, possibly consisting of invalid double-digit codes.

When the format indicator F is used, the field must occupy one complete word; when the format indicator @F is used, the field must occupy two complete words.

Within the considerations listed above, the use of macro-instructions allows the programmer extreme flexibility in arranging the fields within a record area. In fact, the use of macro-instructions permits the programmer to refer, with equal ease, to fields which overlap or bridge words and to fields which occupy all or part of a word. It should be noted, however, that the most efficient coding will result when the amount of word-bridging is held to a minimum.

### **Additional Examples**

Additional examples of the use of the DA statement in conjunction with other statements are included throughout the manual.

## CODE

CODE is a special type of declarative statement which may *only* appear as one of the subsequent entries under a DA header line. It has the following uses:

1. To name a field in which one or more code values may be present during the running of the object program.
2. To name and define these code values which are the condensed representations of various conditions, categories or classifications, etc.

The primary use of CODE is in connection with the LOGIC and DECOD macro-instructions. By means of these macro-instructions, a code field may be interrogated for the presence or absence of a specific code value. Various switches may be set or branches may be made depending upon the decision. When CODE is used, interrogation may be made by referencing the *symbolic name* of the code value for the condition to be tested, rather than the code value itself. CODE symbolic names may not be referred to in symbolic machine instructions.

### Source Program Format

The CODE entry is itself a header line, giving the position and format of the code field. Subsequent entries consist of the symbolic names of the conditions to be tested during the course of the program, each indented one space, with the actual code value indicated for each condition.

A CODE entry may be followed by subsequent fields of the DA or by another CODE entry. If CODE is the last DA entry, any other type of entry may follow.

Code fields may not be more than ten numerical digits or five alphameric characters in length and must be wholly contained within one word of the record defined by the DA.

### CODE Header Line

The basic format of the CODE header line is as follows:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
01	A.N.Y.L.A.B.E.L.	C.O.D.E.	2,0.,29,A8.,2					
02								

A symbolic label must always be supplied for the CODE entry and the letters CODE must be written in the operation column. The operand includes field definition as in any other DA subsequent entry, except that the area defined may not be more than a single storage word and it may not bridge words.

Two types of format indicators (shown below) are allowed. The desired indicator is written immediately after the terminal digit position of the field definition.

Code Format Indicators	Meaning
A	The field is to contain numerical data which is to be treated as an automatic-decimal number.

The format indicator may be followed by a decimal position indicator of the form "a.b" as described on page 43. In this case, however, the field size may not be greater than 10.

@

The field is to contain the double-digit representation of alphameric characters. The field size defined must not exceed 10 digits (5 characters); it must be evenly divisible by 2. The starting digit position must be an even number.

If a format indicator does not appear immediately after the ending digit position in the CODE header line, the processor will consider the field to contain an automatic-decimal integer.

Consider the following example:

Line	Label	Operation	OPERAND					
3	5 6	15 16	20 21	25	30	35	40	45
0 1		D.A	1					
0 2	ANY LABEL	C.O.D.E.	2 0 . 2 9 A 8 . 2					
0 3								

The CODE header line defines a field occupying the entire third word of the record area. At object program time, the field is to contain numerical data to be treated as an automatic-decimal number with 8 integer digits and 2 decimal digits.

The following example defines a field which is to contain one alphameric character.

Line	Label	Operation	OPERAND					
3	5 6	15 16	20 21	25	30	35	40	45
0 1		D.A	2					
0 2	ANY LABEL	C.O.D.E.	1 4 . 1 5 @					
0 3								

### CODE Subsequent Entries

A symbolic label is always required for a CODE subsequent entry so that the code value may be referenced by name. For this entry, Autocoder requires that the *label* be indented exactly one space. Thus, the maximum size of this type of label is nine characters. The operation column must be blank.

The operand is used to define a code value which may appear in the field during the running of the object program. A code value may not exceed the length of the field as defined by the CODE header line; the maximum length, therefore, is 10 numerical digits or 5 alphameric characters. Each code value must also be consistent with the format indicator(s) which appear in the CODE header line.

If the CODE header line specifies an automatic-decimal field, therefore, the operand of each CODE subsequent entry under that header line must contain a numerical



value. This value may be signed or unsigned. (The sign, however, is superfluous; the digits will be considered absolute in sign. The sign of the word will not be tested or changed.)

In the example below, the CODE header line defines a field which is to contain data to be treated as a one-digit automatic-decimal integer. Each subsequent entry, therefore, must define a one-digit numerical code value.

Line	Label	Operation	OPERAND
3	5 6	15 16 20 21	25 30 35 40 45
0 1	ANY LABEL	DA	1
0 2	COLOR	CODE	2 5 A 1
0 3	BLACK		0
0 4	GREEN		1
0 5	RED		5
0 6	WHITE		6
0 7	BLUE		8
0 8			

In this example, COLOR is the name of the defined field. In the source-language program, the field may be tested for the presence of a color such as BLACK. The processor, however, will generate coding which will cause the object program to test for the presence of a "0." If the field is to be tested for the presence of GREEN, the processor generates coding to test for a "1." Thus, the programmer may make reference by means of a macro-instruction to the condition or name of a code value and the processor will translate this into instructions which actually check for the presence of the code value itself.

In the example below, CLASS is the name of a field which is to contain a three-digit, automatic-decimal number. In this case, the number is to include two integer digits and one decimal digit. Each subsequent entry, therefore, must define a numerical code value containing the number of decimal digits specified by the CODE header line. This is done by inserting a decimal point at the proper position within the code value. The decimal point does not occupy a storage location; it merely indicates the treatment of the integer and decimal digits in the code value.

Line	Label	Operation	OPERAND
3	5 6	15 16 20 21	25 30 35 40 45
0 1	AREA NAME	DA	1
0 2	CLASS	CODE	1 0 1 . 1 0 3 A 2 . 1
0 3	A		1 1 . 9
0 4	B		1 0 . 9
0 5	C		9 . 9
0 6	D		8 . 9
0 7			

Although the code value must contain exactly the same number of decimal digits as defined in the CODE header line (in the above example, one decimal digit),

leading zeros may be omitted from the integer digits, as in the code values for C and D above.

When CODE defines an alphameric field, the actual code values pertinent to the various conditions should be indicated in the operand by writing the @ character, followed by the alphameric characters which comprise the code value, followed by another @ character. The code value may *not* include the @ character itself; the form @@@ or @A@B@C@ is thus prohibited. In addition, the record mark may *not* be used as a code value. Except for being restricted to one word, the operand of an alphameric CODE subsequent entry may be identical to that allowed for an alphameric literal.

Consider the following example:

Line	Label	Operation	OPERAND
3 5 6	15 16	20 21	25 30 35 40 45
0 1		DA	3, , 0 + INDEXWORD
0 2	STATE	CODE	1 4, 1 9 @
0 3	ALABAMA		@ALA@
0 4	GEORGIA		@GA @
0 5	FLORIDA		@FLA@
0 6			

If the field, STATE, is tested during the running of the object program and the contents at that time are "ALA" the STATE will be considered to be ALABAMA; if the contents are "GA" the STATE will be considered to be GEORGIA; etc. If leading blanks appear, they may be omitted. Thus, if the values within the fields have been right-justified (rather than left-justified, as in the example above) the value for GEORGIA could have been written @GA@ rather than @GA @.

## Processing Techniques

Since the order is irrelevant, CODE subsequent entries may appear in any sequence, provided that they follow a CODE header line. Only those code values which are to be tested need to be entered as a CODE subsequent entry, even when it is known that other data will appear in the field during the running of the object program.

The label of a CODE header line may be freely referenced in exactly the same ways that any other DA subsequent entry may be referenced; in each case the assigned location and the field control of its DA position will be compiled in the address portion of the referencing instruction. The data that will be operated on will be the data in the CODE field at that particular time in the running of the object program; a CODE subsequent entry will *not* cause a constant to be assigned to that area of storage. The names of the code values, which must be unique symbols and follow the usual rules for symbols, may not be referenced by symbolic machine instructions.

## Additional Examples

Additional examples of the use of the CODE statement in conjunction with the LOGIC and DECOD macro-instructions appear on pages 173 and 177.

## DC — Define Constant

The declarative statement, DC, may be used to enter numerical, alphameric and address constants (adcons) into the object program, and to assign names to constants for ease of reference. Like the DA statement, the DC statement consists of a header line and one or more subsequent entries with blank operation columns. Unlike the DA statement, a DC statement actually causes specified *data* to be compiled as a part of the object program. The DC header line causes the processor to assign storage locations to the constants which are defined by subsequent entries.

### Source Program Format

#### DC Header Line

The formats for a DC header line are as follows:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	ANY LABEL	DC		+ RDW				
0 2	ANY LABEL	DC		- RDW				
0 3	ANY LABEL	DC		RDW				
0 4	ANY LABEL	DC						
0 5								

ANYLABEL may be any symbolic label or it may be omitted. It may *not* be an actual address. The entries DC and RDW must be written exactly as shown.

If plus or minus RDW is written in the operand, the processor will generate a single RDW with the indicated sign. If no sign is indicated, a minus RDW will be generated. This RDW, which will define the area containing all of the constants included in the subsequent entries, will be stored in the location immediately preceding the constant area. If present, the label of the DC header line may also be used to refer to the RDW. If the operand is blank, no RDW will be generated and the DC label will refer to the first word of the constant area.

#### DC Subsequent Entries

Four principal classes of constants may be defined in a DC subsequent entry: automatic-decimal numbers, floating-decimal numbers, address constants, and alphameric constants. These constants may be entered in either of two ways. (1) They may be entered separately, with individual symbolic names assigned to each, or (2) several constants of the same type may be defined by one subsequent entry, with one symbolic name assigned to the first location of the series.

Each constant may be entered in a separate core storage word or, within the limitations described below, several constants may be "packed" into the same word by means of field definers. Each subsequent entry may appear with or without a symbolic label.

*In all examples of location assignments, it is assumed that the Autocoder assignment counter was at 1000 when the DC header line was encountered.*

*Automatic-Decimal Constants.* This type of constant, used to enter signed numbers when the constant is to be referred to by a macro-instruction, may be from one to twenty digits in length. A decimal point may be included to indicate the magnitude of the number according to ordinary usage; it is neither stored in the constant nor saved with it, but merely serves as an indicator to Autocoder of the desired decimal-point placement. If the decimal point falls to the right of the right-most digit of the number, it may be omitted; the number will be considered an integer or an ordinary numerical constant.

If the constant is to be referred to by a symbolic machine instruction, however, it may neither exceed ten digits in length nor bridge words. The decimal point will be ignored; a symbolic machine instruction will treat the constant as an integer.

When used as a DC subsequent entry, automatic-decimal numbers may be individually named (with or without field definition) or written in a series, with the first number named.

INDIVIDUAL NAMES. The following example illustrates how automatic-decimal constants may be entered with individual symbolic names:

Line	Label	Operation	OPERAND
3	56	15 16	20 21 25 30 35 40 45
0 1		D.C.	
0 2	CONSTANT1		+1,2,3,4,5.,6,7,8,9,0
0 3	CONSTANT2		-1.,2,3
0 4	CONSTANT3		-4,1,2,6,7,9,3,6,5,4,1.,3,9,2,1,7
0 5	CONSTANT4		-2
0 6			

The Autocoder processor would make the following actual location assignments:

Symbol	Field Definition	Location	Contents
CONSTANT1	0,9	1000	+1234567890
CONSTANT2	0,9	1001	-0000000123
CONSTANT3	0,9	1002	-0000412679
	0,9	1003	-3654139217
CONSTANT4	0,9	1004	-0000000002

Note that each constant is right-justified in a separate storage word (two words in the case of a constant in excess of 10 digits). Reference to CONSTANT3 by a *symbolic machine instruction*, therefore, would result in the consideration of digit positions 0,9 of location 1002 only.

The same constants could be packed into fewer words by using field definition in the subsequent entries, as in the following example:

Line	Label	Operation	OPERAND
3	56	15 16	20 21 25 30 35 40 45
0 1		D.C.	
0 2	CONSTANT1		+1,2,3,4,5,6,7,8,9,0(.0.,9.)
0 3	CONSTANT2		-1.,2,3(.0.,2.)
0 4	CONSTANT3		-4,1,2,6,7,9,3,6,5,4,1.,3,9,2,1,7(.3.,1,8.)
0 5	CONSTANT4		-2(.9.)
0 6			

Actual location assignments would then be as follows:

<i>Symbol</i>	<i>Field Definition</i>	<i>Location</i>	<i>Contents</i>
CONSTANT1	0, 9	1000	+1234567890
CONSTANT2	0, 2	1001	-1234126793
CONSTANT3	3, 9		
CONSTANT4	0, 8	1002	-6541392172
	9, 9		

When field definition is used with an automatic-decimal constant, it must define not less than the number of digits in the number and not more than 20 digits. Since the left field definition may be no greater than 9, the right field definition may be no greater than 28.

Constants written with field definition will be packed into words according to the following rules:

1. A change in sign will start a new word.
2. Field definition which would, in effect, result in overlapping in the same word will force the constant into a new word.

**SERIES.** The operand of a DC subsequent entry may include a series of signed automatic-decimal numbers, each without field control. All of the automatic-decimal numbers must share the same attributes, i.e., if the first number is written with a decimal point, they must *all* be written with a decimal point; if the first number has two decimal places, *all* must have two decimal places. (The number of integer positions, however, need not be the same.)

Since field definition is not allowed, each constant will be right-justified in one or two locations, depending on the length of the largest constant. Automatic-decimal constants 11-20 digits in length, therefore, should not be intermingled with constants 1-10 digits in length in the same subsequent entry since *every* constant would be assigned two locations.

The following entry will cause the indicated constants to be right-justified in separate words. (Note that commas are *not* used as separators.)

Line	Label	Operation	OPERAND							Basic										
			3	6	15	16	20	21	25		30	35	40	45	50					
01	CONAREA	D.C.																		
02																				
03																				

Autocoder will make the following assignments:

<i>Symbol</i>	<i>Field Definition</i>	<i>Location</i>	<i>Contents</i>
CONAREA	0, 9	1000	+0000000152
		1001	+0000000038
		1002	+0000000097
		1003	-0000000021
		1004	-0000000304
		1005	+0000016541

Since the operand of the DC header line is blank, no RDW is generated. Therefore, when a symbolic machine instruction references the label CONAREA, it will refer to the first constant, +0000000152; CONAREA+1 will refer to +0000000038; CONAREA(8,9)+2 will refer to +97; etc.

The following example illustrates the generation of a plus RDW to be associated with the constant area. Note also the use of integer constants.

Line	Label	Operation	OPERAND
3	5 6	15 16	20 21 25 30 35 40 45
0 1	CON AREA	D C	+RDW
0 2			+1,5,2+3,8+9,7-2,1-3,0,4+1,6,5,4,1
0 3			

Autocoder would make the following assignments:

Symbol	Field Definition	Location	Contents
CONAREA	0, 9	1000	+0010011006
		1001	+0000000152
		1002	+0000000038
		1003	+0000000097
		1004	-0000000021
		1005	-0000000304
		1006	+0000016541

The label CONAREA would now refer to the RDW defining the constant area; CONAREA(7,9)+1 will refer to the first constant, +152; CONAREA+2 will refer to +0000000038; etc.

*Floating-Decimal Numbers.* Floating-decimal numbers may be entered in the operand of a DC subsequent entry with the format

$$\pm nF \pm m$$

where  $\pm n$  is a decimal or integer number of not more than eight digits and  $\pm m$  is a one- or two-digit exponent. If the sign preceding  $m$  is omitted,  $m$  is considered to be positive. A decimal point may be used to indicate implied positioning of the decimal point in the number  $\pm n$ . If no decimal point appears, the number will be considered an integer. The value of the number is  $\pm n$  multiplied by  $10^m$ . Thus,  $-.98765432F+1$  would represent the number  $-.98765432 \times 10^1$ . The exponent  $m$  may be omitted if equal to 0.

The Autocoder processor will consider the signs, the value of  $n$ , and the value of  $m$  of the entry in the format,  $\pm nF \pm m$ . The processor will then generate a standard 7070 normalized floating-decimal word, which is of the form

$$\pm MMNNNNNNNN$$

where the sign is the sign of  $n$ . The number,  $\pm nF \pm m$ , is normalized by placing  $n$  between +1 and -1 and by adjusting the value of  $m$  accordingly. The value of  $MM$  equals 50 plus or minus the adjusted value of  $m$  and the value of  $NNNNNNNN$  is the normalized value of  $n$ . For example, the number  $-1860.723 \times 10^3$  would be represented by  $-1860.723F+3$ . Autocoder converts this number to

the standard 7070 normalized floating-decimal format by normalizing the number to  $-.1860723F+7$  and converting it to the form  $-571860723$ .

Since a floating-decimal number will always occupy ten digit positions and it may not bridge words, field definition is not allowed.

INDIVIDUAL NAMES. The following examples illustrate the use of individually named constants:

Line	Label	Operation	OPERAND						
3	5,6	15,16	20,21	25	30	35	40	45	
01		D,C							
02	FLTNUMBER1			-3	.	4	F		
03	FLTNUMBER2			+3	.	4	.	5,6,7,8,9,3F-2	
04	FLTNUMBER3			-3	.	1	.	9,2F+7	
05	FLTNUMBER4			-2	.	9,5,6,7	.	1F-3	
06	FLTNUMBER5			+1	.	2,5,4,6	F+1	5	
07									

Assignment would be made as follows:

Symbol	Field Definition	Location	Contents
FLTNUMBER1	0,9	1000	-5134000000
FLTNUMBER2	0,9	1001	+5034567893
FLTNUMBER3	0,9	1002	-5931920000
FLTNUMBER4	0,9	1003	-5229567100
FLTNUMBER5	0,9	1004	+7012546000

SERIES. The operands of DC subsequent entries may include a series of floating-decimal numbers, each written in the  $\pm nF \pm m$  format described above. Unlike a series of automatic-decimal constants, the floating-decimal numbers *must* be separated by commas and the decimal point *need not* appear in the same place in each number.

The floating-decimal numbers used in the previous example could also have been written as follows:

Line	Label	Operation	OPERAND							Autocoder					
3	5,6	15,16	20,21	25	30	35	40	65	70						
01		D,C													
02	FLTCON			-3	.	4	F	.	+3	.	4	.	5,6,7,8,9,3F-2		
03															

The assigned locations would be the same as for the previous example with the exception that the single label FLTCON would be used to refer to the first constant at 1000, FLTCON+1 would refer to the second, etc.

*Address Constants (Adcons).* An adcon is used to produce an address constant which is the numerical representation (usually 4 digits) of the location

to which a given symbolic label has been assigned. As noted in a previous section, an adcon may be used as a literal or as a dc subsequent entry. When used as either, the form is  $\pm$ SYMBOL. When used as a dc subsequent entry, adcons may be individually named (with or without field definition) or written in a series in which the first adcon may be named.

INDIVIDUAL NAMES. Adcons may be named individually; they may appear with or without field definition. Adcons written without field definition will be right-justified in separate words. Adcons written with field definition will be packed into words according to the following rules:

1. A change in sign will start a new word.
2. Field definition which would, in effect, result in overlapping in the same word will force the adcon into a new word.
3. If less than four digit positions are defined, the low-order digit(s) of the address will appear in the constant.
4. Adcons should not bridge words.

Consider the following example:

Line	Label	Operation	OPERAND											
			3	5	6	15	16	20	21	25	30	35	40	45
0.1		D.C.												
0.2	A.D.C.O.N.1													
0.3	A.D.C.O.N.2													
0.4	A.D.C.O.N.3													
0.5	A.D.C.O.N.4													
0.6														

Autocoder would first assign locations to ADDRESS1, ADDRESS2, ADDRESS3, and ADDRESS4. Assuming that these locations were 1125, 0025, 3542, and 6453, respectively, Autocoder would then assemble the adcons as follows:

Symbol	Field Definition	Location	Contents
ADCON1	0, 9	1000	-0000001125
ADCON2	0, 1	1001	+2535420000
ADCON3	2, 5		
ADCON4	6, 9	1002	-0000006453

SERIES. Several adcons may be written sequentially in a single dc subsequent entry as shown below. Field definition and address adjustment are not allowed; commas are not used between adcons.

Line	Label	Operation	OPERAND											
			3	5	6	15	16	20	21	25	30	35	40	45
0.1		D.C.												
0.2	A.D.C.O.N.S													
0.3														



Assuming that the labels have been assigned locations as in the previous example, Autocoder would assemble the adcons as follows:

Symbol	Field Definition	Location	Contents
ADCONS	0, 9	1000	-0000001125
		1001	+0000000025
		1002	+0000003542
		1003	-0000006453

The first adcon could be referred to by the label ADCONS; the second by ADCONS+1; etc.

*Alphameric Constants.* Alphameric information may be included in a constant area by using a DC subsequent entry similar to those in the following example:

Line	Label	Operation	OPERAND	B								
3	5	6	15	16	20	21	25	30	35	40	45	
0 1		DC										
0 2	ANYLABEL		@1 2 3 4 5 @									
0 3			@ CONSTANT, ALPHAMERIC, INFO, @									
0 4			@A @B @C @									
0 5			@ ABC @R									
0 6			@ @R									
0 7												

ANYLABEL may be any symbolic label, or it may be omitted. Each operand must begin and end with the @ character. All characters between the *initial* @ character and the *final* @ character will be converted to double-digit form and assembled successively, beginning in the high-order positions of each location assigned. The sign of each word used to contain the characters will be alpha.

Packing of sequentially listed constants is automatic; it continues from line to line. If packing is to be avoided, therefore, blanks must be inserted within constants to fill out each storage word. If the processor finds that the total number of characters assigned through the end of a DC subsequent entry is not a multiple of five and the next entry is not an alphameric constant, blanks (double-digit 00) will be assembled in the low-order positions of the last word assigned.

The constant may contain the digits 0 through 9, the characters of the alphabet, blanks, and the special characters

. □ & \$ \* - / , % # @

Note that the @ character may be included at any point in a constant field. It should also be noted that the processor will scan the *entire* operand for the *last* @ character. Therefore, this character *must not* appear in the remarks portion of the card if the constant field is to be properly defined.

A record mark may be defined by following the terminal @ with the character R, as explained below.

If symbolic machine instructions refer to the label of a DC header line which specifies the generation of an rdw, the assembled instruction will address the rdw. If, however, symbolic machine instructions refer to the label of a DC header line that does not specify rdw generation, or to the label of a DC subsequent entry defining an alphameric constant, they will act only on the first five characters of the area or the field, respectively. When defined for use by symbolic machine instructions, therefore, each subsequent entry should contain a maxi-

imum of five characters. A macro-instruction, however, may properly refer to the label of the header line or any subsequent entry and act on the entire constant area (exclusive of the RDW, if any) or field, regardless of length.

An alphameric constant may be used to define a message, as in the following example:

Line	Label	Operation	OPERAND
3	5 6	15 16 20 21	25 30 35 40 45
0 1	EOFMSG	D.C.	-RDW
0 2			@REMOVE TAPE ON UNIT, OO @
0 3			@AND MOUNT, NEXT @
0 4			

Autocoder will make the following assignments:

Symbol	Field Definition	Location	Contents
EOFMSG	0, 9	1000	-0010011008
		1001	R E M O V
		1002	E T A P
		1003	E O N
		1004	U N I T
		1005	O O A N
		1006	D M O U
		1007	N T N E
		1008	X T

Note that the label EOFMSG refers to the RDW defining the constant area.

In many instances, it may be desirable to name a portion of a long alphameric constant so that modification of the constant, such as entering variable tape addresses, codes, etc., may be readily accomplished. Thus, the constant in the previous example could have been defined as follows:

Line	Label	Operation	OPERAND
3	5 6	15 16 20 21	25 30 35 40 45
0 1	EOFMSG	D.C.	-RDW
0 2			@REMOVE TAPE ON UNIT, @
0 3	TAPENO		@OO@
0 4			@ AND MOUNT, NEXT @
0 5			

Autocoder will make the following assignments:

<i>Symbol</i>	<i>Field Definition</i>	<i>Location</i>	<i>Contents</i>
EOFMSG	0, 9	1000	-0010011008
		1001	R E M O V
		1002	E T A P
		1003	E O N
		1004	U N I T
TAPENO	0, 3	1005	0 0 A N
		1006	D M O U
		1007	N T N E
		1008	X T

Note that Autocoder will assign the proper field definition to each symbolic name given so that reference to TAPENO will result in appropriate field control. If a field is to be referred to by symbolic machine instructions, however, the programmer is responsible for insuring that any named portion of an alphameric constant does not bridge words.

The following are examples of the use of the @ character *within* a constant field:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1		D C	
0 2	C O N S T A N T 1		@ A @ B @ C @
0 3	C O N S T A N T 2		@ @ @
0 4			

Autocoder will make the following assignments:

<i>Symbol</i>	<i>Field Definition</i>	<i>Location</i>	<i>Contents</i>
CONSTANT1	0, 9	1000	A @ B @ C
CONSTANT2	0, 1	1001	@ b b b b

RECORD MARK. Since the record mark may not be entered through the IBM 7500 Card Reader, special provision is made for generating it in an alphameric constant by entering the letter R immediately following the terminal @ character.

If the last character generated by the current alphameric constant does not complete a word, the record mark will be assembled in positions 8 and 9 of that word, and any intervening positions will contain the double-digit representation of blanks. If the word is complete, the record mark will be assembled in positions 8 and 9 of the next word and positions 0 through 7 will contain the double-digit representation of blanks. Consider the following example:

Line	Label	Operation	OPERAND											
			3	5	6	15	16	20	21	25	30	35	40	45
0 1		D,C												
0 2	CONSTANT1													
0 3	CONSTANT2													
0 4	CONSTANT3													
0 5														

Autocoder would make the following assignments:

Symbol	Field Definition	Location	Contents
CONSTANT1	0, 5	1000	A B C D E
CONSTANT2	6, 9		F G b b †
CONSTANT3		1001	H I J K L
		1002	b b b b †
		1003	

(Note the packing of the first two characters of CONSTANT2 into the word which holds CONSTANT1.)

The following coding shows how a record mark may be entered and named so that field definers and/or address adjustment are not necessary when operating on it:

Line	Label	Operation	OPERAND											
			3	5	6	15	16	20	21	25	30	35	40	45
0 1		D,C												
0 2	CONSTANT1													
0 3	RECMARK1													
0 4	CONSTANT2													
0 5	RECMARK2													
0 6														

Autocoder will make the following assignments:

Symbol	Field Definition	Location	Contents
CONSTANT1	0, 9	1000	A B C D E
RECMARK1	8, 9	1001	b b b b †
CONSTANT2	0, 3	1002	F G b b †
RECMARK2	8, 9		

The first record mark may now be referred to by the label RECMARK1; the second by RECMARK2.

### Additional Examples

The coding on the following page is furnished to illustrate various operands which may occur in a DC statement. Note that while various types of constants may be generated in the same DC area, they may *not* be mixed in the same subsequent entry.



## DLINE — Define Line

The DLINE statement and subsequent entries provide the programmer with a convenient means of specifying:

1. The layout of a print-line area.
2. The editing of fields to be inserted in that print-line area.

Various formats used in the subsequent entries provide the ability to define fields in which the following information may appear:

1. Constant data which is always to be included in the print line.
2. Alphameric data which is to be inserted by the object program.
3. Data edited to floating-decimal print format which is to be inserted by the object program.
4. Data edited to numerical print format which is to be inserted by the object program.

Several print-line formats are usually required for the preparation of a single report, e.g., a heading line, a transaction line, a total line, etc. These lines may be defined by means of separate DLINE header lines, each followed by its own subsequent entries.

The DLINE header lines set up and name the print areas; the subsequent entries name the printing positions to be used, define the characteristics of the fields, and cause the generation of constants which are to appear in the print area. The use of DLINE, therefore, results in the establishment of an output print image in core storage in which two core storage digit positions are reserved for each print position included. This area may initially include constant information specified by the programmer, e.g., record marks, captions, etc. Fields to be inserted by the object program will normally be moved into the image by an EDMOV statement. This macro-instruction will perform the necessary editing, insert decimal points and commas, convert data into double-digit form, etc. It is also possible to insert unedited information by means of the MOVE macro-instruction.

### Source Program Format

#### DLINE Header Line

The basic format for the DLINE header line is as follows:

Line	Label	Operation	OPERAND					
3	3-6	15-16	20-21	25	30	35	40	45
01	ANYLABEL	DLINE						
02								

ANYLABEL is any symbolic label; it should not be omitted. It may *not* be specified as the operand of a DTF header line. DLINE must be written exactly as shown. The operand must be blank except for remarks.

#### DLINE Subsequent Entries

The header line must be followed by one or more subsequent entries, as in the example on the following page.

Line	Label	Operation	OPERAND
3	5	15	20
6	16	21	25 30 35 40 45
0 1	L I N E N A M E	D L I N E	
0 2			1 0 @ T O T A L @
0 3			1 8 @ \$ @
0 4	G R O S S A M T		1 9 X , , X X X . Z Z ) D R , C R
0 5	C H E C K A M T		6 0 \$ X , , X X Z . Z Z ) , C
0 6	I T E M N A M E		8 0 . 9 4
0 7	F L V A R		9 5 F
0 8			1 2 0 @ @ R
0 9			

The DLINE entry and the subsequent entries, other than those which define constant fields, must be labeled. The operand of each subsequent entry must begin with a number which indicates the print position for the left-most character of the field. Thus, the word "TOTAL" in the example above would appear in print positions 10 through 14, the dollar sign on line 03 would appear in print position 18, etc. Print position 0 may *not* be specified.

The processor will issue a warning message if a print position beyond 135 is named or included in a DLINE area; however, the line may be of any desired length.

Descriptions and formats follow for each of the subsequent entries used to define the various types of data fields that may appear in a DLINE.

*Constant Data.* Constant information may be included in a DLINE area by using subsequent entries similar to those in the following examples:

Line	Label	Operation	OPERAND
3	5	15	20
6	16	21	25 30 35 40 45
0 1	L I N E N A M E	D L I N E	
0 2	A N Y L A B E L		7 @ C O N S T A N T , I N F O R M A T I O N @
0 3			3 0 @ B L A N K S , A R E , I N C L U D E D @
0 4			1 0 0 @ \$ @
0 5			1 0 5 @ @ @
0 6			1 0 9 @ A @ B @ C @ D @
0 7			

A subsequent entry of this type may be written with a symbolic label, as in line 02 above. However, *unlike* the other types of fields which may be defined by subsequent entries, a constant field may appear *without* a label.

The operand must begin with the number of the print position for the left-most character in the field, immediately followed by the @ character. The @ character indicates that all *following* characters, up to (but not including) the *final* @ character, are to appear in the print area. The constant may contain numbers, characters of the alphabet, blanks, and any special characters (other than the record mark) acceptable to the input device used.

Note that the @ character may be included at *any* point in a constant field. Thus, line 05 would cause a single @ character to print in position 105 and line 06 would cause the following to print in positions 109-115:

A@B@C@D

A record mark may be entered by following the final @ with an "R" as shown in the following example:

Line	Label	Operation	OPERAND					
3	5 6	15 16	20 21	25	30	35	40	45
0 1	L I N E N A M E	D L I N E						
0 2				1	@	A	B	C @ R
0 3				1	5	@	A @ R	
0 4				1	1	7	@ @ R	
0 5								

The record mark will be positioned according to the following rules:

1. When constant information appears between the initial and final @ character, the record mark will be placed in the next available print position ending in "0" or "5," with blanks inserted in the intervening print positions. The entry in line 02 above would cause a blank to appear in print position 4 and a record mark to appear in print position 5. The entry in line 03 above would cause blanks to appear in print positions 16 through 19 and a record mark to appear in print position 20.
2. The appearance of @@R will cause a *single* record mark to appear in the print position *named*. The entry in line 04 above would cause a record mark to appear in print position 117.

Constant fields, unlike the other fields which may be specified, will be initially loaded with the information specified. No provision is made for regenerating this information during the running of the program; if the object program includes coding to alter these constants, it must also include any coding necessary to restore them. Print positions for constant fields must appear in ascending sequence.

*Alphameric Data Field.* When data is to be inserted into an alphameric field in the print line, the field may be indicated by using subsequent entries similar to those in the following example:

Line	Label	Operation	OPERAND					
3	5 6	15 16	20 21	25	30	35	40	45
0 1	L I N E N A M E	D L I N E						
0 2	A N Y L A B E L 1			1				
0 3	A N Y L A B E L 2			4	5			
0 4	A N Y L A B E L 3			2	1	1	0	0
0 5								

Any symbolic label may be used with an entry of this type; it should *not* be omitted, however. The operand must begin with the number of the print position for



the first or left-most character of the field. If more than one print position is to be included in the field, the first number must be followed immediately by a comma and the number of the print position for the last or right-most character in the field as in lines 03 and 04 above.

The object program will usually insert unedited alphameric data into a DLINE field of this type by means of a MOVE macro-instruction. If numerical data is to be edited to alphameric form, however, the EDMOV macro-instruction may be used. It is also possible to use symbolic machine instructions if the programmer is careful to consider word boundaries.

*Data Edited To Floating-Decimal Print Format.* When the object program uses the EDMOV macro-instruction to cause data to appear in the print area, edited to floating-decimal print format, the field in which the number is to appear must be defined by means of a subsequent entry under a DLINE statement. An entry of this type would be written as follows:

Line	Label	Operation	OPERAND					
3	5 6	15 16	20 21	25	30	35	40	45
0 1	L I N E N A M E	D L I N E						
0 2	A N Y L A B E L 1		1 3	F				
0 3	A N Y L A B E L 2		1 0 1	F				
0 4								

Any symbolic label may be used with an entry of this type; the label should *not* be omitted, however. The operand consists of the number of the print position for the left-most character, followed immediately by the symbol F. A field of this type requires a total of thirteen print positions.

Floating-decimal numbers will be edited to the printing notation,

$$\pm MM \pm .NNNNNNNN$$

where  $\pm MM$  is a two-digit exponent and  $\pm .NNNNNNNN$  is a normal, eight-digit number (see page 57). The value of the number is  $\pm .NNNNNNNN$  multiplied by  $10^{\pm MM}$ . For example, @9591999897@9695949372 (which in single-digit form would be  $-5198765432$ ) will be printed as  $\pm 01-.98765432$  representing the number  $-.98765432 \times 10^1$ .

*Data Edited To Numerical Print Format.* The combined use of DLINE and the EDMOV macro-instruction will provide a convenient way to do extensive editing of data to a numerical print format. This editing may include the insertion of the comma and decimal characters, the retention or elimination of leading zeros and commas, and the "floating" of a dollar sign. In addition, explanatory characters may be printed to identify plus or minus numbers.

A field of this type may be indicated by using subsequent entries similar to those in the following example:

Line	Label	Operation	OPERAND
3	5	15	20 21 25 30 35 40 45 50
01	LINENAME	DLINE	
02	ANYLABEL1		11X,XXX.ZZ
03	ANYLABEL2		19\$XXX,XXX.XX
04	ANYLABEL3		30\$XX,XXZ.ZZ)PLUS,MINUS
05	ANYLABEL4		45\$ZZZ.ZZ)PLUS
06	ANYLABEL5		56\$XXX,XXX,XXX,XXX,ZZ),MINUS
07	ANYLABEL6		80\$XXX,XXX,XXX,XXX)PLUS,MINUS
08			

Any symbolic label may be used with an entry of this type; it should *not* be omitted, however. The operand must begin with the number of the print position for the first or leftmost character of the field. The rest of the operand is used to specify the desired format of the edited field, and (optionally) to specify the explanatory characters which are to identify plus or minus amounts.

Reference to a field of this type by an EDMOV macro-instruction will cause the generation of instructions which edit the data and insert dollar signs, commas, decimals, and explanatory characters. The decimal point in automatic-decimal variables will be aligned with the decimal point indicated in the DLINE field. The following characters may be used to indicate the format of the field:

Character	Explanation
X	This character represents a digit that is to be replaced by a blank if it is a high-order zero (i.e., a zero not preceded by a significant digit) to the left of the decimal point (if any). It will also be replaced by a blank, regardless of its position, if all digits are represented by Xs (as in lines 03 and 07 above) and all digits are zero.
Z	This character represents a digit that is always to be printed, even when it is a high-order zero.
\$	This character represents a fixed dollar sign when followed by a Z (as in line 05 above) or a floating dollar sign when followed by one or more Xs (as in line 03 above). The \$ need not appear (see line 02 above), but, when it is used, it must immediately follow the print position number.
	A floating dollar sign will print immediately to the left of the first printed digit. However, it will not float any farther to the right than the print position immediately to the left of the deci-

mal point (if one appears). Thus, the \$ in line 03 might print in position 19, 20, 21, 23, 24, 25, or 26, depending on the placement of the leftmost significant digit in each amount printed. When all digit positions are replaced by blanks, the \$ will also be replaced by a blank.

A fixed dollar sign will print in the print position indicated. The \$ in line 05 will always print in position 45, since the code character X does not appear to its right. (An illustration of another way to denote a fixed dollar sign appears on line 03, page 66, where a dollar sign is entered in a constant field.)

,

This character represents a comma when used *within* the format of the field. Commas may appear to the left or right of the decimal point. The comma will print only when digits are printed to its left. If the code character X allows the suppression of all high-order zeros to the left of the comma, the comma will also be suppressed. When the floating dollar sign is used, therefore, it will also float over commas that are to be suppressed. Since zeros to the right of the decimal point cannot be suppressed, a comma to the right of the decimal point will always appear.

.

This character represents the decimal point. Only one decimal point may appear in any given field. When all digits are replaced by blanks, the decimal will also be replaced by a blank.

)

A closing parenthesis may follow the rightmost position of the edited field. This code character does *not* occupy a print position; it is necessary only if explanatory characters follow, e.g., the PLUS and MINUS in lines 04, 05, 06, and 07, above. The use of explanatory characters will be explained below.

When the above characters are used to indicate the format of an edited field, at least one X or Z must appear, but not more than 20. When all digits are represented by an X and no significant digit appears in the amount, the entire field (including the dollar sign, commas, and decimal point positions) will be replaced by blanks.

When a field of this type is considered by the EDMOV macro generator, any explanatory characters which appear beyond the closing parenthesis (if this option is elected) are treated as follows:

1. Two consecutive blanks will signal the end of the operand.
2. Any characters which appear between the closing parenthesis and the comma to the right (or two consecutive blanks) will be inserted in the print line when the edited amount field is positive or blank.

Any characters to the right of the comma will be inserted if the field is negative. The comma may be omitted if blanks are to be inserted for a negative field.

3. The explanatory characters may include numbers, letters of the alphabet, a *single* blank, and special characters (other than comma and record mark) acceptable to the input device used. A comma would signal the end of a set of characters; the maximum size for a set is one word.
4. The number of print positions required for explanatory characters will be equal to the largest of the two sets which may appear. Blanks will be inserted to the right of the shorter set when it appears.

Consider the following examples:

Line	Label	Operation	OPERAND																			
			3	5	6	15	16	20	21	25	30	35	40	45								
01	LINE,NAME	D,LINE																				
02	FIELD,ONE		17	\$	X	,	X	X	.	Z	Z	)	+	,	-							
03	FIELD,TWO		30	\$	X	,	X	X	.	Z	Z	)	D	R	,	C	R					
04	FIELD,THREE		50	\$	X	,	X	X	.	Z	Z	)	P	L	U	S	,	M	I	N	U	S
05	FIELD,FOUR		70	\$	X	,	X	X	.	Z	Z	)	,	C	R							
06																						

A positive and a negative amount, respectively, would print as follows:

FIELDONE, print positions 17 through 26.

\$1,234.56+

\$1,234.56-

FIELDTWO, print positions 30 through 40.

\$1,234.56DR

\$1,234.56CR

FIELDTHREE, print positions 50 through 63.

\$1,234.56PLUS

\$1,234.56MINUS

FIELDFOUR, print positions 70 through 80.

\$1,234.56

\$1,234.56CR

A *single* blank may appear at one or more points within a set of explanatory characters, as in the following example:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	L I N E N A M E	D L I N E						
0 2	F I E L D N A M E		2 7	\$ X X Z . Z Z	+		-	**
0 3								

A positive amount and a negative amount, respectively, would print as follows in print positions 27 through 38:

\$123.45+  
\$123.45- \*\*

## Processing Techniques

### General

At object program time, DLINE areas are initialized by condensed load cards which set all print positions within the area to blanks. These cards are followed by condensed load cards which initialize the constant data fields to the contents specified.

Since data fields will not necessarily occupy complete words, these cards will also cause blanks to be loaded into print positions which occupy the remainder (if any) of a word in which constant data falls. Constant data fields, therefore, should be entered in ascending order. If the processor encounters a constant data field which (1) is lower in sequence than the previous print position reserved, and (2) specifies a print position in a word in which a print position has been previously specified as constant data, the resultant condensed load card will overlay the previously specified print positions, thus yielding inconsistent results.

Although constant data fields are restricted to ascending sequence, the other subsequent entries may appear in any order. Thus, they may be used to specify print positions which cause different labels (and different characteristics, if desired) to be assigned to fields which occupy the same or overlapping print positions as in the following example:

Line	Label	Operation	OPERAND					
3	56	1516	2021	25	30	35	40	45
0 1	P R I N T L I N E	D L I N E						
0 2	F I E L D O N E		2 0	\$ X . X X X . Z Z				
0 3	F I E L D A		2 0	. 2 8				
0 4	F I E L D T W O		3 0	\$ X X X . Z Z				
0 5	F I E L D B		3 0	\$ X X X . Z Z				
0 6	F I E L D T H R E E		4 0	\$ X . X X X . Z Z				
0 7	F I E L D C		4 2	. 5 0				
0 8								

It should be noted that a unique problem arises if the two fields overlap in the manner indicated by lines 06 and 07 above. Reference to either field will cause that field to be inserted as specified. However, if `FIELD C` is inserted, data may remain in print positions 40 and 41 from a previous use of `FIELD THREE` and, when `FIELD THREE` is inserted, data may remain in print positions 49 and 50 from a previous use of `FIELD C`. Thus, the object program must include additional coding to blank out data remaining from use of the other field.

Since all print positions are initially set to blanks, it is not necessary for the programmer to provide coding to blank out the print positions between fields. However, if the object program includes coding which alters these undefined print positions, it must also include coding to restore them to blanks when this becomes desirable.

If, in some instances, certain fields in a line are to be blank, they should be blanked out by a `ZERO` macro-instruction which references the label of the appropriate `DLINE` subsequent entry. If the `ZERO` macro-instruction refers to the label of a `DLINE` header line, however, it will permanently blank out constant data fields (if any) in the line as well as clear the variable fields. It will not, however, affect the dollar signs, commas, and decimals in the edited fields; these are inserted by coding which is generated by the `EDMOV` macro-instruction.

### **On-line Printing**

During the running of the object program, data will be edited and moved into the `DLINE` area which serves as an image of the line to be printed. When all data is included in the area, the object program will issue an output statement which will transmit data from the `DLINE` area to the printer. This would normally be accomplished by a print command which refers to one or more `RDW`s defining the area(s) to be printed.

If the standard `IBM 7400` printer panel wiring is used, line spacing and carriage spacing may be determined by digits 6, 7 and 9 of a Control Information Word as described in the 7070 Data Processing System Bulletin "IBM 7070 Utility Control Panels," form J28-6095, pages 8 and 9. Thus, the programmer is responsible for establishing a constant to fit the requirements for the Control Information Word as described in that bulletin. The print command must then reference one or more `RDW`(s) which causes the Control Information Word to be the first word sent to the printer, followed by the data in the `DLINE` area. Note that the Control Information Word will not be printed.

It should also be noted that the use of the standard `IBM 7400` utility control panel provides for printing from typewheels 1 through 75 on one print cycle and from typewheels 76 through 120 on another print cycle. Thus, while a single `DLINE` statement may be used to describe a print line which uses both sets of typewheels, a single print command may *not* be used to print this entire line.

The following coding, however, would accomplish this purpose:

Line	Label	Operation	OPERAND
3 5 6	15 16	20 21	25 30 35 40 45
01		UW	2, F, I, R, S, T, H, A, L, F
02			
03			
04		UW	2, S, E, C, O, N, D, H, A, L, F
05			
06			
07	F, I, R, S, T, H, A, L, F	DRDW	+C, I, W, 1, C, I, W, 1
08		DRDW	-A, B
09	S, E, C, O, N, D, H, A, L, F	DRDW	+C, I, W, 2, C, I, W, 2
10		DRDW	-C, D
11			

The coding on line 07 above defines the first Control Information Word. Words 1-15 of the DLINE area are defined on line 08. The coding on line 09 defines the second Control Information Word. Words 16-24 of the DLINE area are defined on line 10.

The first Control Information Word (ciw1) would contain a digit "0" in position 5 to cause printing from typewheels 1 through 75. The second Control Information Word (ciw2) would contain a digit "1" in position 5 to cause printing from typewheels 76 through 120; a digit "6" in position 6 would cause spacing to be suppressed before printing.

### Off-line Printing

When a tape is to be prepared for off-line printing, the following characteristics of the printer must be considered:

1. The first character of each line may be used to control carriage skipping and spacing.
2. The number of print-line records per tape block and the size of each record must not exceed that which the printer will accept.
3. A record mark may be required to separate the print lines within a tape block.

Thus, when the DLINE subsequent entries are written, the programmer should first refer to the appropriate printer manual to determine the allowable tape and printer formats.

When the first character of the print line is to be used to control carriage skipping and spacing, the other DLINE fields will print from the type wheel to the left of the position named. Therefore, if the character "1" on line 02 of the following example is used for carriage control, the characters "HEADER" will print in type wheels 1 through 6.

Line 3	Label 5 6	Operation 15 16 20 21	OPERAND					
			25	30	35	40	45	
0 1	PRINTLINE	DLINE						
0 2	CARRIAGE		1	@	1	@		
0 3			2	@	HEAD	ER	@	
0 4								

Thus, by defining the proper constant as "print position" 1 of each DLINE area the programmer can predetermine the carriage control for the various lines on the output listing; e.g., a header line would have a "1" in position 1, detail lines might have a "0" for double spacing, etc. Another method would be to insert the proper alphameric control character through programming, e.g., by reference to the label, CARRIAGE, above.

When the data has been edited and moved to the DLINE area, normal procedure would be to insert the assembled print line in an output file, where it is blocked under control of the Input/Output Control System. The following macro-instruction would accomplish this function:

Line 3	Label 5 6	Operation 15 16 20 21	OPERAND					
			25	30	35	40	45	
0 1		PUT	PRINTLINE	IN	OUTPUTFILE			
0 2								

The following coding might also be used:

Line 3	Label 5 6	Operation 15 16 20 21	OPERAND					
			25	30	35	40	45	
0 1		PUT	OUTPUTFILE					
0 2								
0 3								
0 4		MOVE	PRINTLINE	TO	OUTPUTFILE			
0 5								

In each case, PRINTLINE is the label of the DLINE header and OUTPUTFILE is defined as a tape output file. PUT and MOVE are explained in the 7070 Data Processing System Bulletin, "IBM 7070 Input/Output Control System," form J28-6033-1.

As soon as the PUT (or PUT and MOVE) macro-instruction(s) have placed the DLINE area in the output file, the object program may proceed to prepare the next print line in the DLINE area.



Note that the `DLINE` area itself may *not* be defined as a file, i.e., the operand of a `DTF` header line may *not* be the label of a `DLINE` header line.

Particular consideration must be given to the use of the record mark in tapes prepared for off-line printing on the `IBM 720`. If more than one print-line record is to appear in a tape block, each record must end with a record mark, *with the exception of the last record in each block*. The object program, therefore, must include coding to count the lines that have been moved by the `PUT` macro-instruction and must cause a record mark to be inserted at the end of each record *except* the last one in a block. The end-of-job routine must also include coding to insure that no record mark appears at the end of the final print line.

### **Additional Examples**

Additional examples of the use of the `DLINE` statement in conjunction with the `EDMOV` macro-instruction appear on pages 210 and 212.

## DRDW — Define Record Definition Word

The declarative statement, DRDW, may be used to generate an RDW defining any area of storage specified by the programmer. It may also be used to cause the generation of one or more RDWs associated with an area defined by a DA or DC statement in some other part of storage, i.e., not immediately preceding the DA or DC area.

### Source Program Format

#### Single RDW

When used to generate a single RDW for a given area, the format of the DRDW statement is as follows:

Line	Label	Operation	OPERAND
3	5 6	15 16	20 21      25      30      35      40      45
0 1	ANY LABEL	DRDW	+ADDRESS 1 , ADDRESS 2
0 2	ANY LABEL	DRDW	-ADDRESS 1 , ADDRESS 2
0 3	ANY LABEL	DRDW	ADDRESS 1 , ADDRESS 2
0 4			

ANYLABEL may be any symbolic name or it may be omitted; it may *not* be an actual address. DRDW must be written exactly as shown. ADDRESS1 and ADDRESS2 are the limits of the area to be defined by the generated RDW. Either ADDRESS1 or ADDRESS2 (or both) may be an actual, \*, or symbolic address. The \* or symbolic addresses may appear with or without address adjustment. The sign of the generated RDW is determined by the sign preceding ADDRESS1; if a sign does not appear, however, the generated RDW will be signed minus. For example, if ADDRESS1 had been assigned to location 4372 and ADDRESS2 had been assigned to location 4408, the processor would have generated the RDW, -0043774400, at the point where the following DRDW was encountered:

Line	Label	Operation	OPERAND
3	5 6	15 16	20 21      25      30      35      40      45
0 1	ANY LABEL	DRDW	ADDRESS 1 + 5 , ADDRESS 2 - 8
0 2			

The addresses in the operand of a DRDW statement may be the same, as in the following examples:

Line	Label	Operation	OPERAND						
3	5/6	15/16	20/21	25	30	35	40	45	
01	ANY LABEL 1	DRDW	+	*	*				
02	ANY LABEL 2	DRDW	-	3	2	4	3	2	4
03	ANY LABEL 3	DRDW	ADDRESS, ADDRESS						
04									

When this is the case, the generated RDW will define a one-word area. Thus, line 02 above would cause the generation of the RDW, -0003240324.

### Multiple RDWs

As explained under "DA-Define Area" and "DC-Define Constant," writing "RDW" on a DA or DC header line will cause the processor to generate an RDW(s) associated with the DA or DC entry and to assign it a storage location immediately preceding the defined area. Sometimes, however, it may be advantageous to cause the RDW(s) associated with a DA or DC statement to be generated in some other portion of storage, i.e., *not* immediately preceding the DA or DC area. This may be accomplished by using a DRDW statement with one of the following formats:

Line	Label	Operation	OPERAND					
3	5/6	15/16	20/21	25	30	35	40	45
01	ANY LABEL	DRDW	+	HEADRLABEL				
02	ANY LABEL	DRDW	-	HEADRLABEL				
03	ANY LABEL	DRDW	HEADRLABEL					
04								

ANYLABEL may be any symbolic name or it may be omitted; it may *not* be an actual address. DRDW must be written exactly as shown. HEADRLABEL must be the label of a DA or DC header line which appears as an entry in the program sequence. When this format is used, address adjustment is not allowed. If HEADRLABEL is the label of a DA header line, the number of RDWs generated will be the same as the number of areas designated by the DA header line. One RDW will be generated if HEADRLABEL is the label of a DC header line. A plus sign preceding HEADRLABEL will cause all RDWs generated to be signed plus; a minus sign will cause all RDWs generated to be signed minus. If no sign is shown, as in the format on line 03 above, all RDWs generated will be signed plus except the last which will be signed minus. When using symbolic machine statements, the programmer may make reference to the first RDW by referring to the label of the DRDW; address adjustment may be used to refer to the subsequent RDWs.

### Additional Examples

The coding on the following page illustrates some of the operand forms which might appear in a DRDW statement.

Line	Label	Operation	OPERAND							
			3	6	15	16	20	21	25	30
01	ANY LABEL	DRDW	HEADR LABEL							
02		DRDW	+HEADR LABEL							
03		DRDW	-HEADR LABEL							
04		DRDW	+1,99							
05		DRDW	-325,4999							
06		DRDW	0,9989							
07		DRDW	+*,*							
08		DRDW	-* -1, *							
09		DRDW	+*,*+2							
10		DRDW	-*+1,*+7							
11		DRDW	*-1,*+2							
12		DRDW	+ADDRESS1, ADDRESS2							
13		DRDW	-ADDRESS1+1, ADDRESS2							
14		DRDW	+ADDRESS1, ADDRESS2+5							
15		DRDW	-ADDRESS1-10, ADDRESS2-11							
16		DRDW	ADDRESS1-2, ADDRESS2+5							
17		DRDW	+*, ADDRESS							
18		DRDW	-ADDRESS, *							
19		DRDW	+ADDRESS+1, *+1							
20		DRDW	+0, *							
21		DRDW	-* , 4999							
22		DRDW	+325, ADDRESS+2							
23		DRDW	-ADDRESS, 9989							
24		DRDW	0, ADDRESS							
25		DRDW	*, ADDRESS							
		DRDW	ADDRESS, 9989							

## DSW — Define Switch

The primary function of the dsw declarative statement is to provide from one to ten digital switches which may be considered singly, or as a group, by the SETSW and LOGIC macro-instructions. SETSW and LOGIC will treat these switches as logically equivalent to electronic switches, although processed in a slightly different fashion. (The switches may *not* be referred to by electronic switch commands, e.g., ESN, ESF, etc.)

Each switch that is generated by the dsw statement occupies one digit position of a word and is considered OFF if its digit value is zero and ON if its digit value is other than zero, regardless of the sign of the word. Since the switches are generated at the point where the dsw statement is encountered, this statement should not appear within a series of machine instructions.

### Source Program Format

The format for a dsw entry is as follows:

Line	Label	Operation	OPERAND							Base	
3	5	15	16	20	21	25	30	35	40	45	50
01	ANYLABEL	DSW	SWITCH1, -SWITCH2, +SWITCH3, etc.								
02											

ANYLABEL may be a symbolic name or it may be omitted. The entry dsw must be written exactly as shown. As many as ten symbolic switch names may appear, with a comma inserted between names. Continuation cards may be used, if necessary. The name of each switch must be unique; i.e., it must not be defined elsewhere as the label of another item. The initial setting of a switch is determined by the following:

1. If a plus sign, or no sign, precedes the name of a switch, the switch will be considered ON and set to "1."
2. If a minus sign precedes the name of a switch, the switch will be considered OFF and set to "0."

### Processing Techniques

When the dsw statement is encountered, the processor will construct a one-word, positive, numerical field. The leftmost position will contain either "1" or "0," depending on whether the first-named switch in the dsw is to be initially ON or OFF. Succeeding digit positions will indicate the status of the remaining switches, in the order they are named. If less than ten switches are named, the remaining digits are set to zero. It should be noted that, while the switches are initially loaded as described, the programmer must provide additional coding (e.g., a SETSW macro-instruction) to reinitialize the switches if they are to be utilized in multi-pass programs.

Reference to the label of the dsw statement by symbolic machine instructions or

macro-instructions will result in reference to the *entire* word used to contain the switch settings. Consequently, if a label is supplied for a dsw, the entire set of switches may be tested or altered by LOGIC and SETSW statements.

The programmer is warned against trying to initialize electronic switches by using a dsw with an ORIGIN to 101, 102, or 103. The switches would be treated as digit switches, ignoring the fact that they are electronic switches.

### Additional Examples

The coding on the following page illustrates various operands which might appear in a dsw statement. An additional example appears on page 187, where the defined switches are referred to by the SETSW macro-instruction.

Line 3	Label 56	Operation 15 16	OPERAND 20 21 25 30 35 40 45	Basic Autocoder			Autocoder	
				50	55	60	65	70
0 1		D SW	SW I T C H A					
0 2	ANY LABEL 1	D SW	- SW I T C H B					
0 3		D SW	+ SW I T C H C					
0 4	ANY LABEL 2	D SW	SW I T C H D , SW I T C H E					
0 5	ANY LABEL 3	D SW	+ SW I T C H F , - SW I T C H G , SW I T C H H					
0 6	ANY LABEL 4	D SW	- SW I T C H 1 , SW I T C H 2 , SW I T C H 3 , + SW I T C H 4 , - SW I T C H 5 , + SW I T C H 6 ,					
0 7			SW I T C H 7 , - SW I T C H 8 , SW I T C H 9 , - SW I T C H 1 0					
0 8								

## EQU — Equate

The EQU statement may be used to equate a symbol to:

1. An actual or symbolic address.
2. An index word or electronic switch number.
3. A channel, tape unit, combined channel and tape unit, combined arm and file, unit record synchronizer, inquiry synchronizer, or alteration switch number.

Thus, the EQU statement provides a convenient way to cause one or more symbols to be assigned to an actual location or to machine hardware. In this way, the same item may be referred to by different names in different parts of a program. Meaningful and easily remembered symbols may be used throughout the program, rather than the actual machine numbers which might be required in the operand of some Autocoder entries. In addition, when it is necessary to change the actual location or machine number, it is more convenient to change a single EQU statement than to alter each Autocoder statement which might otherwise contain the actual number.

### Source Program Format

The general formats for an EQU statement are as follows:

Line	Label	Operation	OPERAND
3	5 6	15 16	20 21 25 30 35 40 45
0 1	ANY LABEL	EQU	ADDRESS
0 2	ANY LABEL	EQU	ADDRESS, Y
0 3	ANY LABEL	EQU	, Y
0 4			

The entry EQU must appear exactly as written. ANYLABEL is the symbolic name which is to be equated to ADDRESS. ANYLABEL may *not* be defined elsewhere in the same source program and may *not* be omitted.

ADDRESS may be any of the following:

1. An actual address, with or without field definition.
2. An index word or electronic switch number.  
address adjustment. This symbolic entry *must* appear as the label of an Autocoder entry elsewhere in the source program (not necessarily previously). It may *not* appear as the *label* of another EQU statement.
3. The number of an index word, with or without field definition.
4. The number of an electronic switch.
5. The number of a channel, tape unit, combined channel and tape unit, combined arm and file, unit record synchronizer, inquiry synchronizer, or alteration switch.
6. Omitted, in which case the first character of the operand must be a comma, followed by an X, S, T, etc., as described on page 86.

Y is replaced by a one- or two-character code which identifies ADDRESS as a particular piece of machine hardware.

Since EQU statements do not actually occupy core storage locations in the object program, they may be inserted at any point in the source program, provided that they are not intermingled with the subsequent entries under a DA, DC, or DLINE header line. (This is in contrast to the other declarative statements, which must be separated from the program instruction area.)

## Processing Techniques

The method of coding each of the general uses of the EQU statement is described below.

### Actual or Symbolic Address

A symbolic name may be made equivalent to an actual or symbolic address. The symbolic name to be equated is written in the label columns. The operand may contain an actual or symbolic address, with or without field definers. Address adjustment may also be used with a symbolic address.

Line	Label	Operation	OPERAND
3	56	15 16 20 21	25 30 35 40 45
0 1		DA	1
0 2	C U S T N O		2 , 9
0 3		.	
0 4		.	
0 5		.	
0 6		.	
0 7		.	
0 8	C L A S S	EQU	C U S T N O ( 0 , 1 )
0 9			

The above entries will cause the processor to assign the same location to the symbol CLASS as was previously assigned to CUSTNO. In addition, CLASS will be given field definers denoting the two high-order positions of CUSTNO, i.e., 2, 3.

#### EXAMPLE:

Line	Label	Operation	OPERAND
3	56	15 16 20 21	25 30 35 40 45
0 1		DA	1
0 2	R A T E		2 , 6 A 2 , 3
0 3		.	
0 4		.	
0 5		.	
0 6		.	
0 7		.	
0 8	C L A S S	EQU	R A T E S E E W A R N I N G B E L O W
0 9			



The above entries will cause the processor to assign the same location to the symbolic name, CLASS, as was previously assigned to RATE. In addition, CLASS will be given the same field definers as RATE, i.e., 2, 6.

If address adjustment had been used in the preceding examples, only the assignment of a location would have been affected; the field definition would have been derived in the same manner as before. It is important to note, therefore, that the programmer is responsible for insuring that the field definition will actually be that desired for the new location.

The additional characteristics defined in the operand of a declarative statement will *not* be assigned to the name in the label field of an EQU statement. Thus, in the preceding example, CLASS will *not* be identified as a numerical field containing an automatic-decimal number. This creates no difficulty with symbolic machine instructions, since they do not use these characteristics, but limits the utility of the equated symbols in macro-instruction operands. When a symbol is equated to another symbol, a macro generator will treat the symbol in the label of the EQU statement as if it had the characteristics of a single whole word. Unless specifically desired, therefore, the use of equated symbols in a macro-instruction operand may cause program errors and should be avoided.

Two other methods are suggested for assigning two (or more) different symbolic names to the same field and the same characteristics. One method is to list both names as subsequent entries under the same DA, repeating the starting and ending digit positions and format indicators of the field, as in the example on page 45.

Another method of assigning two different symbolic names to the same location is to place the second name (with format indicators as desired) under a separate DA that is made equivalent to the first through the use of an Origin Control statement as explained on page 90.

Note that the EQU statement does not allow the following transitive relation:

Line	Label	Operation	OPERAND					
3	5 6	15 16	20 21	25	30	35	40	45
0 1	A	E Q U	B					
0 2			.					
0 3			.					
0 4	B	E Q U	C					
0 5								

Statements of this form are not acceptable; they will result in an incorrect location assignment. (The desired effect may, however, be obtained by writing: A EQU C; B EQU C.)

**Index Word or Electronic Switch Number**

A symbolic name may be made equivalent to an index word or electronic switch. The symbolic name to be equated is written in the label columns. The operand contains the one- or two-digit word (1-99) or electronic switch number (1-30), followed by a comma and the letter X or S, respectively. An index word may be field-defined as illustrated below. The index word or electronic switch that is equated to the symbolic name will be reserved during compilation; i.e., it will be passed over when Autocoder assigns symbolic index words and electronic switches to actual locations during Phase III.

EXAMPLES:

Line	Label	Operation	OPERAND										
			3	5	6	15	16	20	21	25	30	35	40
0 1	L O O P C O U N T	E Q U	1	.	X								
0 2				.									
0 3				.									
0 4	I X W O R D	E Q U	5	2	(	2	,	5	)	.	X		
0 5				.									
0 6				.									
0 7	S W I T C H A	E Q U	2	5	.	S							
0 8													

The first entry will assign the name LOOPCOUNT to index word 1. The second entry will assign the name IXWORD to index word 52, with field definition (where applicable) of 2, 5. The third entry will assign the name SWITCHA to electronic switch 25.

**Input/Output Units and Alteration Switches**

A symbolic name may be made equivalent to a particular piece of machine hardware. The symbolic name to be equated is written in the label columns. The operand contains the number or value of the item, followed by a comma and an explanatory code character. The explanatory code characters used in EQU statement operands are as follows:

Item	Code
Tape Channel and Unit	CU
Tape Channel	C
Tape Unit	U
Disk Storage Arm and Unit	AF
Index Word	X
Electronic Switch	S
Alteration Switch	SN
Unit Record Latch	I
Unit Record Synchronizer	
Reader	R
Printer	W
Punch	P
Inquiry Synchronizer	Q
Typewriter	T

To illustrate, the following entry will cause the processor to assign the name RESTART to alteration switch 1:

Line	Label	Operation	OPERAND										
			3	5	6	15	16	20	21	25	30	35	40
0 1	R E S T A R T	E Q U	1	.	S	N							
0 2													

This entry would make it possible for the programmer to write the more meaningful entry on line 01 below, rather than the entry on line 02:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
01		BAS	RESTART, RESTORE
02		BAS	1, RESTORE
03			

This particular type of EQU statement may also be used for the specific purpose of identifying a particular item of hardware for the benefit of a macro generator (see "SNAP," example 1, page 228). When used for this purpose, the first entry in the operand may be omitted, as in the following examples:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
01	REGNAME	EQU	, X
02	SWITCH	EQU	, S
03	TYPEWRITER	EQU	, T
04			

## Additional Examples

The coding on the following page illustrates operands which might appear in an EQU statement.

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
01	ANYLABEL	EQU	1000
02	ANYLABEL1	EQU	325(6,9)
03	ANYLABEL2	EQU	OTHERLABEL
04	ANYLABEL3	EQU	OTHERLABEL(0,4)
05	ANYLABEL4	EQU	OTHERLABEL+2
06	ANYLABEL5	EQU	OTHERLABEL(2,3)+1
07	XWORDNAME	EQU	1,X
08	XWORD	EQU	77(2,5),X
09	SWITCH	EQU	29,S
10	XWORDA	EQU	,X
11	SWITCHA	EQU	,S
12	CHANNEL	EQU	1,C
13	ACTIONTAPE	EQU	5,U
14	MASTERCU	EQU	21,CU
15	DISK	EQU	11,AF
16	ALTSWITCH	EQU	3,SN
17	UNITLATCH	EQU	2,I
18	READERNAME	EQU	1,R
19	PRINTNAME	EQU	2,W
20	PUNCHNAME	EQU	3,P
21	INQUIRY	EQU	1,Q
22	TYPEWRITER	EQU	,T
23			

## Control Statements

Control statements are, in effect, orders to the processor which give the programmer control over portions of the assembly process. Thus, `ORIGIN` and `LITORIGIN` statements give the programmer control over the placement of his program in core storage. `BRANCH` statements cause the processor to produce execute cards containing unconditional Branches to locations specified by the programmer. An `END` statement will cause the processor to compile all remaining generated material and then produce an execute card containing an unconditional Branch to a location specified by the programmer. Control over the assignment of locations to symbolic index words and electronic switches is maintained through the use of `XRESERVE`, `SRESERVE`, `XRELEASE` and `SRELEASE` statements.

The formats and detailed descriptions of the use of these control statements are presented below. In all cases, the operation will be `CNTRL`. The labels must be prepared exactly as shown: `ORIGIN`, `BRANCH`, etc. The operand may vary as described for each control statement.

## **ORIGIN Control and LITORIGIN Control**

ORIGIN statements order the processor to override its automatic assignment of storage locations and to begin the assignment of succeeding entries at the particular location specified by the programmer. Thus, they enable the programmer to control storage assignments of source-language input such as area definitions, constants, and instructions (including those generated "in-line" by macro-instructions). If an ORIGIN statement does not appear before the first such entry in a source program, the processor will begin the assignment of storage locations at an address specified to the Compiler Systems Tape. This address is originally 0325, but it may be altered as described in the IBM 7070/7074 Data Processing System Bulletin "IBM 7070/7074 Compiler Systems: Operating Procedure," form J28-6105.

LITORIGIN statements are used:

1. To partition or "segment" a program in order to enable the correct loading of a multi-phase program by causing the immediate compilation of all remaining material generated "out-of-line" in each segment (i.e., since the last previous LITORIGIN, or since the beginning of the program, if no LITORIGIN appears).
2. To regulate the placement of this material.

Material generated "out-of-line" in each segment includes generated constants, generated area definitions, generated symbolic subroutines, all literals, and all adcons. This material will normally be assigned locations immediately following the highest location assigned to the source program. The use of LITORIGIN, however, makes it possible to assemble this material at the end of each section or phase of a program so that it may be loaded with that section or phase. In addition, LITORIGIN is used to specify the beginning core-storage location at which the generated material is to be assigned.

The assignment of actual program locations is effected by means of location counters which may be named symbolically by the programmer and used by the Autocoder processor. The programmer has complete control over which counter is to be used while assigning locations to a given section of the program. In addition, he also controls the setting and resetting of the counters as desired. When in current use, however, a counter will be advanced automatically by the processor as locations are assigned; thus, after each assignment, it will always contain the address of the next location to be assigned. Provision is also made for "remembering" the minimum and maximum value attained by each counter. In certain cases, e.g., if the processor finds that addresses are to be assigned before the programmer has named the first symbolic counter, use is made of an internal (and, therefore, unnamed) counter.

### **Source Program Format**

The ORIGIN and LITORIGIN statements have identical basic formats, written as follows:

Line	Label	Operation	OPERAND
3	5 6	15 16 20 21	25 30 35 40 45
0 1	ORIGIN	CNTRL	NAMEONE, NAMETWO
0 2	ORIGIN	CNTRL	NAMEONE
0 3	LITORIGIN	CNTRL	NAMEONE, NAMETWO
0 4	LITORIGIN	CNTRL	NAMEONE
0 5			

The entries ORIGIN, LITORIGIN, and CNTRL must be written exactly as shown. NAMEONE is an entry which supplies the initial value to the location counter, NAMETWO.

NAMEONE may be any of the following types:

1. A symbolic label, with or without address adjustment, which has appeared previously.
2. The character \*, with or without address adjustment. It will be considered to have the value (as adjusted) of the location counter in use at the time this statement is encountered.
3. An absolute machine address.
4. The symbolic name of a location counter which has been established in some previous ORIGIN or LITORIGIN statement, with or without address adjustment.
5. From 1 to 97 entries of the preceding types, separated by commas, enclosed by parentheses, and preceded by the characters MAX, as shown in the example on page 91. The largest (adjusted) value in this set will be used as the value to be established. If any of the entries in the set are the names of location counters, the *highest* value it has attained (rather than the last value) will be the value for comparison. Continuation cards, described on page 17, may be used as necessary.

In each of the five cases above, the value of NAMEONE will be placed in the counter NAMETWO, which will then be used to assign subsequent locations. If NAMETWO is omitted, an unnamed counter will be used for subsequent location assignments. NAMETWO may be the symbolic name of a previously defined location counter or the symbolic name of a new location counter. In either case, no address adjustment is allowed.

It is recommended that programs should normally be written with consistent NAMETWO usage: either (1) always spell out a NAMETWO counter, or (2) always omit any reference to NAMETWO. The latter case will mean that the unnamed counter is used throughout a program; this will often be suitable when a program is simple and straightforward, with little segmentation or overlaying of program areas. It can be seen that if NAMETWO is used consistently in a program, omitting a NAMETWO in some statement may cause the counter in use to be changed unwittingly.

When macro generators create an ORIGIN or LITORIGIN statement, there is an exception in its processing; i.e., an omitted NAMETWO will *not* cause the counter in use to be changed. The current counter will continue to be the one used for subsequent assignment because there is no way for a set of generated coding to refer to the counter that is in current use, but is obvious that this counter must remain the effective one.

To be compatible with Four-Tape Autocoder, an ORIGIN or LITORIGIN CNTRL with a blank operand has a special function. Each location counter, other than the counter named "S," is examined to determine the highest previous location (not necessarily the current value) assigned by *any* location counter. The value obtained is placed in the unnamed counter for subsequent assignment.

If statements are assigned locations in the index word area by means of an ORIGIN to an actual address, the corresponding index words will be reserved as they are encountered *during* the assignment pass. This will normally be done early in assembly to avoid duplicate assignment of these words.

Not more than 25 LITORIGIN statements may appear in one source-language program.

## Processing Techniques

In its simplest form, the Autocoder ORIGIN statement is used to indicate the initial location which is to be used in assigning locations to a program. Suppose, therefore, that a program to be processed by Autocoder begins with the entry below and that no other ORIGIN or LITORIGIN statements are present:

Line	Label	Operation	OPERAND			
3	56	1516	2021	25	30	35 40 45
0 1	O R I G I N	C N T R L	5 0 0	C O U N T A		
0 2						

In this case, Autocoder will establish a single location counter, COUNTA, with an initial value of 500. The entire program will be located in sequential locations beginning with 500, and all of the generated material will be assigned locations following the other programming entries.

The ORIGIN statement may also be used to:

1. Assign the same area of storage to several sections of a program.
2. Partition a program into several sections.
3. Assign program sections relative to the size and/or placement of other sections of the program.

Consider, therefore, a program that is to begin in location 1000. In this program, records that are read into a certain area of storage may have three different formats; therefore, three different sets of symbolic names and field definers may be desirable. The programmer may use separate DAS to define the three record formats that are to appear. Then, by proper use of ORIGIN statements, he may cause the processor to assign all three DAS to the same area of storage.

The ORIGIN statements for this program might appear as follows:

Line	Label	Operation	OPERAND						Basic Addr																									
			3	5	6	15	16	20		21	25	30	35	40	45	50																		
0.1	ORIGIN	CNTRL	1	0	0	,	C	O	U	N	T	A																						
0.2		.																																
0.3		.																																
0.4		.																																
0.5	FORMAT1	DA	1	0	,	,	0	+	X	W	O	R	D																					
0.6		.																																
0.7		.																																
0.8		.																																
0.9	ORIGIN	CNTRL	F	O	R	M	A	T	1	,	C	O	U	N	T	B																		
1.0	FORMAT2	DA	1	0	,	,	0	+	X	W	O	R	D																					
1.1		.																																
1.2		.																																
1.3		.																																
1.4	ORIGIN	CNTRL	F	O	R	M	A	T	1	,	C	O	U	N	T	C																		
1.5	FORMAT3	DA	1	0	,	,	0	+	X	W	O	R	D																					
1.6		.																																
1.7		.																																
1.8		.																																
1.9	ORIGIN	CNTRL	M	A	X	(	C	O	U	N	T	A	,	C	O	U	N	T	B	,	C	O	U	N	T	C	)	,	C	O	U	N	T	D
2.0																																		

The last ORIGIN statement insures that succeeding entries (and/or generated material) will be assigned locations beyond the *longest* DA. If all three areas are the same length, or if the longest area is the last one to be defined, the last ORIGIN statement would not be necessary.

In the preceding examples, all generated material would have been assigned locations beyond the last program entry. However, if ORIGIN statements are used to cause subroutines or phases of a program to be assigned to the same or overlapping areas of storage, it may be desirable to include, in each section of the program, the generated material which it has produced. This may be accomplished through the proper use of the LITORIGIN statement as follows:





If subroutines 2 and 3 were to start 25 words after the beginning of subroutine 1, the ORIGIN statements preceding them would have been written:

14		.	
15	ORIGIN	CNTRL	COUNTA+25,COUNTC
16		.	
17		.	
18	ORIGIN	CNTRL	COUNTA+25,COUNTD
19		.	
20			

If an ORIGIN statement to an *actual* address causes succeeding entries to be assigned locations in the index word area, the corresponding index words will be reserved and, if the entries are labeled, the index words will be named. The following example would cause index words 50 through 54 to be reserved and named and index word 55 to be reserved.

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
01	ORIGIN	CNTRL	50
02		DA	1
03	NAMEONE		0,9
04	NAMETWO		10,19
05	NAMETHREE		20,29
06	NAMEFOUR		30,39
07			40,49
08	ORIGIN	CNTRL	
09			

Since the index words are reserved only from the point at which the ORIGIN statement appears, this usage should normally appear at the beginning of the program or following a LITORIGIN statement.

### Additional Examples

The coding on the following two pages illustrates some of the ways in which the assignment of locations may be manipulated through the use of ORIGIN and LITORIGIN control statements.

PAGE AA	PROGRAM CNTRL	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.				PAGE AA	
LN CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF
01	1002	*	ILLUSTRATION OF ORIGIN AND LITORIGIN CONTROL STATEMENTS.				
02	1003		NOP			ASSIGNED BY UNNAMED COUNTER.	
03	1004	ORIGIN	CNTRL 700	00001	0325	-0100090000	
04	1005		NOP			ASSIGNED BY UNNAMED COUNTER.	
05	1006	0400	NOP	00002	0700	-0100090000	
06	1007		NOP			ACTUAL ADDRESS IN LABEL.	
07	1008	ORIGIN	CNTRL 800,COUNTER1	00003	0400	-0100090000	
08	1009		ZAI +1			ASSIGNED BY UNNAMED COUNTER.	
09	1010	LITORIGIN	CNTRL COUNTER1+10,COUNTER2	00004	0701	-0100090000	
						COUNTER1 SET TO 800.	
						ASSIGNED BY COUNTER1.	
						COUNTER2 ASSIGNS LITERAL.	
			LITERALS				
10	X		+1	00005	0800	+1300000811	
11	1011		ZAI +2	00006	00	0811 +1	0811
12	1012	ORIGIN	CNTRL **20,COUNTER2			0812 +1300000923	
13	1013	LABEL	ZAI +2			ASSIGNED BY COUNTER2.	
14	1014	LITORIGIN	CNTRL LABEL+90,COUNTER3	00007	0833	+1300000923	
						COUNTER2 INCREASED BY 20.	
						ASSIGNED BY COUNTER2.	
						COUNTER3 ASSIGNS LITERAL.	
			LITERALS				
15	X		+2	00008	00	0923 +2	0923
16	1015		ZAI +3			0924 +1300001876	
17	1016	ORIGIN	CNTRL COUNTER3-50,COUNTER4			ASSIGNED BY COUNTER3.	
18	1017		ZAI +4	00009	0875	+1300111876	
19	1018	LITORIGIN	CNTRL **1000,S			*SEE NOTE BELOW.	
						ASSIGNED BY COUNTER4.	
						SPECIAL COUNTER USED TO ASSIGN LITERALS.	
			LITERALS				
20	X		+3	00010	00	1876 +3	1876
21	X		+4		11	1876 + 4	1876
22	1019		ZAI +5			1877 +1300000927	
23	1020	ORIGIN	CNTRL COUNTER3,COUNTER3			ASSIGNED BY SPECIAL COUNTER.	
24	1021		ZAI +5	00011	0925	+1300000927	
25	1022	ORIGIN	CNTRL MAX(800,LABEL,COUNTER1,COUNTER2,COUNTER4,*),COUNTER5			*BACK TO COUNTER3, NOTE BELOW.	
26	1023		ZAI +5			ASSIGNED BY COUNTER3.	
27	1024	LITORIGIN	CNTRL SET UNNAMED COUNTER TO MAX ALL BUT S COUNTER.			ASSIGNED BY COUNTER5.	
						ASSIGNED BY COUNTER5.	
						COUNTER5 ASSIGNS LITERAL.	
			LITERALS				
28	X		+5	00	0927	+5	0927
29	1025		NOP			0928 -0100090000	
30	1026	ORIGIN	CNTRL *,COUNTER5			ASSIGNED BY UNNAMED COUNTER.	
31	1027		ZAI +5			VALUE OF UNNAMED COUNTER IN COUNTER5	
32	1028	ORIGIN	CNTRL MAX(COUNTER5,S)			ASSIGNED BY COUNTER5.	
33	1029		ZAI +5			0929 +1300008000	
34	1030	LITORIGIN	CNTRL 8000,COUNTER5	00012	1878	+1300008000	
						SET COUNTER5 TO MAX NAMED.	
						ASSIGNED BY COUNTER5.	
						COUNTER5 ASSIGNS LITERAL.	
			LITERALS				
35	X		+5	00013	00	8000 +5	8000
36	1031	*	NOTE THAT THE USE OF ADDRESS ADJUSTMENT IN LINE 17 DID NOT AFFECT				
37	1032	*	COUNTER3				
38	1033	*					

ORIGIN COUNTER	INITIAL VALUE	LAST VALUE	HIGHEST VALUE	LOWEST VALUE
*UNNAMED*	0325	0928	0928	0325
COUNTER1	0800	0800	0800	0800
COUNTER2	0811	0833	0833	0811
COUNTER3	0923	0924	0924	0923
COUNTER4	0875	0875	0875	0875
COUNTER5	0926	1878	1878	0926
S	1876	1877	1877	1876

## BRANCH Control

The **BRANCH** statement will cause the processor to produce an execute card containing an unconditional Branch instruction. When encountered during the loading of the object program, this instruction will cause the normal loading process to stop and a branch to be executed to a specified location.

### Source Program Format

The **BRANCH** statement should be written as follows:

Line	Label	Operation	OPERAND					
3	5/6	15/16	20/21	25	30	35	40	45
01	B,R,A,N,C,H,	C,N,T,R,L	A,D,D,R,E,S,S,					
02								

**BRANCH** and **CNTRL** must be written exactly as shown. **ADDRESS** is the actual, \*, or symbolic address to which the branch is to be made after the preceding portion of the program has been loaded into storage. If an \* or symbolic address is used, address adjustment and/or indexing may be specified. If an actual address is used, indexing *only* may be specified.

### Processing Techniques

A **BRANCH** statement may be used in conjunction with an **ORIGIN** statement to execute portions of a program already loaded into storage and to overlap these with other instructions, as in the following example:

Line	Label	Operation	OPERAND					
3	5/6	15/16	20/21	25	30	35	40	45
01	S,T,A,R,T,		(F,I,R,S,T	I,N,S,T,R,U,C,T,I,O,N,)				
02								
03								
04								
05								
06	X,Y,Z,	B	L,O,A,D,P,R,O,G					
07	B,R,A,N,C,H,	C,N,T,R,L	S,T,A,R,T,					
08	O,R,I,G,I,N,	C,N,T,R,L	S,T,A,R,T,C,O,U,N,T,E,R,					
09								

When the resultant object program is being loaded, the loading operation will be interrupted by a branch to the location assigned to the symbol **START**, followed by the execution of the instructions from **START** through the instruction located at **xyz**. The instruction located at **xyz** will cause a branch back to the load program, which will then resume the loading of the remainder of the object program. Since the following entries were assigned locations beginning again with the location **START**, they will overlap the instructions which have already been executed. (In this example, it is assumed that the starting location of the load program is symbolic location **LOADPROG**.)

## Additional Examples

The following coding illustrates various operands which might appear in a BRANCH statement:

Line	Label	Operation	OPERAND											
			3	5	6	15	16	20	21	25	30	35	40	45
01	BRANCH	CNTRL	*											
02	BRANCH	CNTRL	*-2											
03	BRANCH	CNTRL	*+1+X,1											
04	BRANCH	CNTRL	*+XWORD											
05	BRANCH	CNTRL	32.5											
06	BRANCH	CNTRL	O+X,1											
07	BRANCH	CNTRL	O+XWORD											
08	BRANCH	CNTRL	LABEL											
09	BRANCH	CNTRL	LABEL+1.0											
10	BRANCH	CNTRL	LABEL+X,8.4											
11	BRANCH	CNTRL	LABEL+2+XWORD											
12														

## END Control

The END statement is an indication to the processor that the end of the source-language program has been reached. When this statement is encountered, the processor will assign locations to all material generated since the last previous LITORIGIN statement or since the beginning of the program if LITORIGIN was not used. In addition, it then produces an execute card containing an unconditional Branch instruction which, when encountered at the end of the loading of the object program, will cause a branch to a specified location.

### Source Program Format

The END statement may be written as follows:

Line	Label	Operation	OPERAND					
3	56	15 16	20 21	25	30	35	40	45
0 1	E.N.D.	C.N.T.R.L	A.D.D.R.E.S.S.					
0 2	E.N.D.	C.N.T.R.L						
0 3								

END and CNTRL must be written exactly as shown. ADDRESS is the symbolic, actual, or \* location to which the branch is to be made after the entire object program has been loaded into storage. It should be noted that a BRANCH statement may not ordinarily be used for this purpose since it would cause the branch to occur before the generated material (if any) was loaded. If an \* or symbolic address is used, address adjustment and/or indexing may be specified. If an actual address is used, indexing *only* may be specified.

### Processing Techniques

If the END statement is used, it must be the last entry in the program. If it is not used, or if it is used with a blank operand, the processor will assign locations to all generated material and then generate an unconditional branch to an address specified to the Compiler Systems Tape. This address is originally 0325, but it may be altered as described in the 7070/7074 Data Processing System Bulletin "IBM 7070/7074 Compiler Systems: Operating Procedure," form J28-6105.

### Additional Examples

The coding on the following page illustrates various operands which might appear in an END statement.

Line	Label	Operation	OPERAND
3	56	1516 2021	25 30 35 40 45
01	END	CNTRL	*
02	END	CNTRL	*+1.0
03	END	CNTRL	*-5+XWORD
04	END	CNTRL	*+X.15
05	END	CNTRL	3.25
06	END	CNTRL	0+XWORD
07	END	CNTRL	0+X.11
08	END	CNTRL	LABEL
09	END	CNTRL	LABEL-.10
10	END	CNTRL	LABEL+XWORD
11	END	CNTRL	LABEL+.5+X.45
12	END	CNTRL	
13			

## XRESERVE Control and SRESERVE Control

References to the actual address of an index word or electronic switch may be made at the discretion of the programmer. Normally, however, he will use symbolic names which the compiler will then assign to actual addresses. XRESERVE and SRESERVE statements cause the processor to reserve index words and electronic switches, respectively, so that they will be "unavailable" when this assignment is made; i.e., symbolic names will not be assigned to the address(es) reserved.

### Source Program Format

The XRESERVE and SRESERVE statements may be written as follows:

Line	Label	Operation	OPERAND
3	5,6	15,16	20,21 25 30 35 40 45
0 1	X,R,E,S,E,R,V,E	C,N,T,R,L	N <sub>1</sub> , N <sub>2</sub> - N <sub>3</sub> , . . . e . t . c .
0 2	X,R,E,S,E,R,V,E	C,N,T,R,L	
0 3			
0 4	S,R,E,S,E,R,V,E	C,N,T,R,L	N <sub>1</sub> , N <sub>2</sub> - N <sub>3</sub> , . . . e . t . c .
0 5	S,R,E,S,E,R,V,E	C,N,T,R,L	
0 6			

The entries XRESERVE, SRESERVE and CNTRL must be written exactly as shown.

The operand field contains the one- or two-digit number of the index word(s) or electronic switch(es) to be reserved, or is blank. Continuation cards, described on page 17, may be used as necessary. The entry N<sub>2</sub>-N<sub>3</sub> reserves all of the index words or switches between and including N<sub>2</sub> and N<sub>3</sub>. When this form is used, N<sub>3</sub> must be a number *greater* than N<sub>2</sub>. The two forms of operand entries may be used exclusively or intermixed freely on each card. A blank operand will cause all of the index words or electronic switches to be reserved.

### Processing Techniques

As described on page 19, the XRESERVE and SRESERVE statements provide one of several methods of making index words or electronic switches unavailable for assignment to symbolic names. Unlike some of the other methods described, XRESERVE and SRESERVE statements affect the availability table only when they are encountered in the statement-by-statement processing during the Phase III pass which assigns index word and electronic switch addresses.

XRESERVE and SRESERVE statements are usually placed at the beginning of a source program. However, they may also appear at a later point(s) in the program. When this is the case, it is possible that a previous symbolic index word or electronic switch name may already have been assigned to the address specified. The programmer must therefore be cautious in the use and placement of these statements.



### Additional Examples

The following coding illustrates various operands which might appear in an XRESERVE statement:

Line	Label	Operation	OPERAND
3	5/6	15/16	20/21 25 30 35 40 45
01	XRESERVE	CNTRL	1
02	XRESERVE	CNTRL	90
03	XRESERVE	CNTRL	3, 4, 7, 9, 2
04	XRESERVE	CNTRL	9-21
05	XRESERVE	CNTRL	1-4, 7-8, 79-80, 45-47
06	XRESERVE	CNTRL	1, 3-4, 7, 9, 11, 15, 13, 17-19, 21
07			2, 3, 25, 27, 29, 31-34, 37-39, 50
08	XRESERVE	CNTRL	
09			

The following coding illustrates various operands which might appear in an SRESERVE statement:

Line	Label	Operation	OPERAND
3	5/6	15/16	20/21 25 30 35 40 45
01	SRESERVE	CNTRL	2
02	SRESERVE	CNTRL	29
03	SRESERVE	CNTRL	3, 4, 15, 7
04	SRESERVE	CNTRL	9-21
05	SRESERVE	CNTRL	1-4, 7-8
06	SRESERVE	CNTRL	1, 3-5, 7, 9, 11, 13, 15,
07			21, 25-27
08	SRESERVE	CNTRL	
09			



results. These statements should be used only where the programmer is certain that the index word or electronic switch is no longer required in its previous role, e.g., at the end of a given phase in a multi-phase program, perhaps following a LITORIGIN statement. It should be noted that where more than one name is made equivalent to a given index word or electronic switch, an XRELEASE or SRELEASE statement referring to any one of the names (or to the actual address itself) has the effect of a statement referring to *all* of the equivalent names.

### Additional Examples

The coding on the following page illustrates various operands which might appear in XRELEASE and SRELEASE statements.

Line	Label	Operation	OPERAND	Basic Autocoder	Autocoder
3	56	15	16 20 21 25 30 35 40 45 50 55 60	65	70 75
01	XRELEASE	CNTRL	1		
02	XRELEASE	CNTRL	9,0		
03	XRELEASE	CNTRL	3,4,7,9,2		
04	XRELEASE	CNTRL	9-21		
05	XRELEASE	CNTRL	1-4, XWORDA		
06	XRELEASE	CNTRL	XWORDB, XWORDC		
07	XRELEASE	CNTRL	XWORDD, 10-20, 30-35		
08	XRELEASE	CNTRL	XWORDE, 6-8, XWORDF, 7, 9, 10-12		
09	XRELEASE	CNTRL	1, 10, XWORDNAMEA, 11-13, XWORDNAMEB, XWORDNAMEI, XWORDNAMED,		
10			XWORDNAMEE, XWORDNAMEF, XWORDNAMEH, 25, XWORDNAMEG, 29-31,		
11			XWORDNAMEL		
12	XRELEASE	CNTRL	XWORDA		
13					
14	SRELEASE	CNTRL	2		
15	SRELEASE	CNTRL	2,9		
16	SRELEASE	CNTRL	3,4,15,7		
17	SRELEASE	CNTRL	9-21		
18	SRELEASE	CNTRL	1-4, SWITCHA		
19	SRELEASE	CNTRL	SWITCHB, SWITCHC,		
20	SRELEASE	CNTRL	SWITCHD, 10-20, 27-29,		
21	SRELEASE	CNTRL	SWITCHE, 2-7, SWITCHF, 2, 9, 11-15		
22	SRELEASE	CNTRL	1, 3, SWITCHI, SWITCHF, 7-11, SWITCHC, SWITCHD, SWITCHE, 2, 0,		
23			SWITCHM, SWITCHG, SWITCHH, SWITCHA, SWITCHB, SWITCHK,		
24			SWITCHL, 2, 9,		
25	SRELEASE	CNTRL	SWITCHM		



The IBM Reference Manual "7070 Data Processing System," form A22-7003-2, contains numerous illustrations of how symbolic machine instructions are written and gives examples of the machine language instructions which will be assembled from them.

The following is a section of a payroll routine which further illustrates the use of symbolic machine instructions. Note that remarks may appear anywhere in the operand, provided two blank spaces separate the remarks from the operand of the instruction.

Line	Label	Operation	OPERAND								
			5	10	15	20	25	30	35	40	45
0.1	CALCTAX	Z A 3	TAXCLASS	DETERMINE	IF						
0.2		M	-1300	PAY	IS						
0.3		A 2	GROSS	TAXABLE							
0.4		B M 2	NOTAX	BRANCH	IF NO TAX						
0.5		Z A 3	9992								
0.6		M	+18	CALCULATE	TAX						
0.7		S R R	2								
0.8		B	*+2								
0.9	NOTAX	S 2	9992	CLEAR	ACCUM2						
1.0		S T 2	WITHHOLD	STORE	TAX AMOUNT						
1.1											

The use of macro-instructions described, below, will make possible the coding of a source program with only limited use of symbolic machine instructions.

## Macro-Instructions

A macro-instruction represents a single operation entry on the Autocoder coding sheet that is converted during assembly into a sequence of machine instructions. Autocoder macro-instructions, including those used for input/output operations, free the programmer to a large extent from attention to machine details such as location assignment and data flow in carrying out the operations required by frequently encountered problems. The user may extend Autocoder by adding appropriate macro generators to the system to process newly-created macro-instructions.

The macro-instructions fall into several categories, as follows:

Category	Macro-Instructions
Input/Output:	OPEN, GET, PUT, PUTX, CLOSE, END
Arithmetic:	ARITH
Decision Making:	COMP, CYCLE and RECYC, DECOD, LOGIC, ZSIGN
Initialization:	SETSW, ZERO, FILL
Data Movement:	EDMOV, MOVE, SHIFT
Reference:	SNAP

A full discussion of the input/output macro-instructions can be found in the 7070 Data Processing System Bulletin "IBM 7070 Input/Output Control System," form J28-6033-1. The programmer should not attempt to use these macro-instructions without careful study of the material contained in that bulletin, especially the DIOCS, DTF, and DUF entries. For convenient reference, however, the formats and brief descriptions of six of the principal input/output macro-instructions are included in this manual.

Programming features common to all macro-instructions are outlined immediately below. A detailed description and the format of each individual macro-instruction follow. Examples are included showing typical source-program statements and the coding generated from them by the macro generators.

*Label.* A macro-instruction which will be referred to elsewhere in the program is written with a label. In all other cases, the label column is blank. The label may be any symbolic label acceptable for a symbolic machine instruction, i.e., beginning with a letter, possibly followed by any combination of up to nine letters or numbers (no special characters). This label will reference the first instruction generated from the macro-instruction.

Certain labels are forbidden. The following sets of characters have special meanings when used in the operands of the macro-instructions specified. They should not, therefore, be simultaneously employed as labels in programs using these macro-instructions:

<i>Characters</i>	<i>Macro-Instructions</i>
ABS	ARITH, COMP
AND	LOGIC, FILL
E	LOGIC
G	LOGIC
IN	PUT, PUTX
L	LOGIC, SHIFT
LC	SHIFT
LS	SHIFT
NOT	LOGIC
NOZERO	ZSIGN
OR	LOGIC
R	SHIFT
RR	SHIFT
RS	SHIFT
TO	EDMOV, GET, MOVE
WITH	FILL
The name of any function in the Macro Table	ARITH

*Actual Addresses.* The operand may not include actual addresses, except where specifically allowed by individual macro-instructions.

*Field Characteristics.* When a program is assembled, it is not known what the contents of the various fields will be at object program time. Therefore, all macro generators must proceed on the assumption that these contents will conform to the characteristics of the field as defined by some declarative statement. If a field has been defined by a DA subsequent entry as alphameric, instructions will be generated to treat the contents as if they were alphameric, and difficulties may arise if, during the object program, numerical information has been stored instead. The same is true if alphameric information is put into numerical fields, or if the numerical modes are not distinguished from each other.

*Asterisks (\*).* When a macro-instruction is written, the programmer does not know how many machine instructions will be generated in its place; therefore,

addresses using the asterisk symbol are not acceptable. For example, if a macro-instruction must reference the next following instruction of the source program, this should be done by means of a label attached to that instruction, not by the address  $* + 1$ . The same considerations make it impossible to do address arithmetic on a macro-instruction label or on any other label by amounts that would carry the address across a macro-instruction.

*Alphameric Literals.* Alphameric literals may appear in macro-instruction operands, provided they are meaningful to the specific instruction involved. It is not possible, however, to write an alphameric literal that includes @ as one of its characters. The scan for macro-instructions, sensing an operand like @AAAAA@AAAA@, would take the second @ to signal the termination of the literal. If such a constant is required by the nature of the program, it must be entered as a subsequent entry under a DC header line; it may then be referenced in the operand of the macro-instruction by its symbolic name.

*Address Modification.* Although field definition is not permitted within macro-instruction operands (except through the device of referencing a label of a DA, DC, or DLINE subsequent entry specifying field definition), indexing and address adjustment are permitted. The conventions for writing these, however, differ from those applicable to symbolic machine instructions. In particular, indexing *precedes* address adjustment if both are present.

Indexing may be specified by referencing the symbolic name of index words or their actual one- or two-digit number. When actual indexing is used, the number of the index word may not be signed, nor should it be preceded by an x. Address adjustment must be signed plus or minus; it is this which distinguishes one- or two-digit adjustments from indexing.

Address modifiers, both index words and address adjustment, must be enclosed by parentheses. The left parenthesis should be in the column immediately following the last character of the address being modified. No blanks should occur anywhere within the entire address modification. If both indexing and address adjustment are used, indexing *precedes* address adjustment, and one set of parentheses should enclose them both.

The following are acceptable examples of modification for macro-instruction operands:

<i>Operand</i>	<i>Explanation</i>
TABLE(ROW)	The symbolic address TABLE is indexed by the index word ROW.
LIST(+34)	The symbolic address LIST is incremented by 34.
LIST(34)	The symbolic address LIST is indexed by index word 34.
CHART(LINE-10)	The symbolic address CHART is decremented by 10; then the address is indexed by the index word LINE.
ARRAY(29+17)	The symbolic address ARRAY is incremented by 17; then the address is indexed by the index word 29.

*Continuation Cards.* Since macro-instructions frequently have long operands, it will often be necessary to use continuation cards to accommodate all the required parameters. All macro-instructions are limited to five cards, i.e., four continuation cards in addition to the header card.

On a continuation card the label and operation columns must be blank, and the continuation of the operand portion must begin in column 21; i.e., it must be left-justified in the operand column of the coding sheet. The entire operand, columns 21-75, may be used.

In the header card or continuation cards (other than the last), each operand need not extend across the entire operand column; it may end with the comma following any parameter and continue on the succeeding card.

In two cases, however, the operand can not be broken off in the middle to be distributed over two cards. These exceptions are:

1. *Alphameric Literals.* Alphameric literals may be up to 120 characters in length and may, in consequence, have to be distributed over two or more cards. This is permitted, provided that the operand columns of all cards but the last are filled to the very end, i.e., through column 75, and that the continuation begins in column 21 of the next card. Blanks which are to be included in the alphameric literal are regarded as part of the literal and may appear in *any* column, including column 75 and column 21. Any extraneous blanks or remarks which appear at the right end of any card before the terminal @ character will be regarded as part of the literal and these characters will be incorporated into the literal.
2. *Address Modification.* When a parameter is modified by indexing and/or address adjustment, separation at the end of a card may be made as follows: The name of the parameter is placed on the first card. The address modification which follows may be on the same card or it may be broken off at any point following the left parenthesis (except *within* the actual or symbolic addresses involved) and continued on the next card beginning in column 21. If the last character of a parameter falls in column 75, then the left parenthesis may be placed in column 21 of the next card.

The examples on the following page illustrate the method of separating alphameric literals and address-modified parameters on continuation cards. Examples 1 and 2 illustrate alphameric literals; examples 3-7 illustrate address modification.

*Punctuation and Spacing.* In general, each macro statement has its own conventions of punctuation, which are stated in detail in the descriptions of the individual instructions. Illegal characters or other faulty punctuation will be regarded as an error condition. For all macro-instructions, any entry in an operand portion of the coding sheets that is preceded by two blanks will not be processed, unless the blanks are inside an alphameric literal. As a general rule, it is recommended that no blanks be written, especially following commas, equal signs, or other punctuation marks, unless a specific demand is made under "Source Program Format" in the description of the individual macro-instruction (e.g., in the *MOVE* statement on either side of the operator *TO*) or when blanks appear in address modification as a result of the use of continuation cards.

*Remarks.* Remark entries may be made at the end of the operand portion of the coding sheet just as with symbolic machine instructions. At least two blanks must precede the first character of a remark entry. These entries, which may include the @ character, will be listed but not processed. If a macro-instruction requires continuation cards, remarks are not necessarily confined to the last card; parameter entries, except for alphameric literals, may be terminated wherever the programmer desires, subject to the rules stated under "Continuation Cards," and continued on the next card, leaving room for remarks at the right end. Attention is called to the fact that macro-instructions are limited to a total of five cards for each instruction; if an operand is very long, remark entries on the header or con-



tinuation cards may waste needed space. In such a case, remarks may be entered on separate comments cards immediately preceding or following those containing the macro-instruction proper.

*Error Conditions.* If the macro generator detects an error condition in analyzing the source statement, an error or warning message will be issued. Warning messages inform the programmer that the machine instructions being generated may have certain unintended effects in the object program, such as accumulator overflow in ARITH, branching to the same location regardless of the outcome of the test in ZSIGN, etc. The programmer should recheck his use of the macro-instruction to make sure that he has employed it correctly and that the special condition will either not arise in his program or that it is intentional. Error messages are issued if a programming error has made it impossible for the macro generator to generate meaningful instructions on the basis of the source statement; in such a case, a NOP will be generated to aid in patching. Assembly will not be interrupted.

Frequently a macro generator will pass a portion of its work on to another generator by putting out what is called a "lower-level" macro-instruction. This will be automatically assembled by the processor, with all generated instructions properly sequenced. The possibility exists, however, that a parameter passed to a lower-level macro generator for processing could bring about an error condition in that generator. In that case, the error message issued would be one from the lower-level generator. The following list shows some cases in which macro generators may call others; it may aid the programmer in interpreting such error messages:

<i>Source Statement</i>	<i>Lower-Level Macro-Instruction</i>
ARITH	Any function in the Macro Table
COMP	ARITH
EDMOV	MOVE
LOGIC	COMP, ZSIGN

Thus, if writing a LOGIC statement has resulted in an error message that is not listed or explained in the LOGIC macro-instruction description, the programmer should consult the descriptions of COMP and ZSIGN, and thereafter, since COMP in turn may have called ARITH, that of ARITH, etc.

*Hardware Usage.* Instructions generated from a macro-instruction may affect the following: (1) the contents of the three accumulators, (2) index words 93 and 94, in addition to such other index words as may be required (these will be assigned in the same fashion as are other symbolic index words), (3) latches as implied by the intent of the macro-instruction (Low, Equal, High for COMP, Accumulator Overflow, Field Overflow, and Sign Change for ARITH, etc.), (4) certain temporary storage areas reserved for working space, and (5) those fields or switches on which the macro-instruction is to operate, i.e., that are specifically named in the operand of the instruction in question. Wherever possible, these fields and switches will be treated non-destructively; thus logical variables in LOGIC, input fields in ARITH, the "from" fields in MOVE, etc. will preserve their contents during the execution of the generated instructions unless the contrary is indicated (e.g., the "to" field and the "from" field in SHIFT are identical). Such fields or switches as are intended to be affected (the result field in ARITH, switches in SETSW, etc.) will, of course, generally lose their previous contents, SNAP will leave the priority mask set to "allow," regardless of its former condition.



## OPEN

OPEN generates instructions to initialize input or output tape files for processing.

### Source Program Format

The basic format for the OPEN statement is as follows:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1	ANY LABEL	OPEN	FILE 1, FILE 2, FILE 3, etc.
0 2			

ANYLABEL is any symbolic label; it may be omitted. The entry OPEN must be written exactly as shown. The operand must contain the name(s) of one or more tape files to be processed. Each name must be the same as the name which appears in the operand of the DTF entry which defines the file. As many tape files as desired may be named in the operand of an OPEN statement, provided the operand does not extend over more than four continuation cards. The names of the tape files must be separated by commas.

### Processing Techniques

The first instruction generated by the processor as a result of an OPEN statement is as follows:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1	ANY LABEL	BLX	IOC.SIXG, IOC.IOPEN
0 2			

Following this, the processor will generate a branch constant containing the address of the first word of the File Specifications Table for each file named in the operand. A NOP will be generated after the last branch constant.

This calling sequence and the subroutine IOC.IOPEN (normally included as a result of a DIOCS statement) will perform the following operations:

1. Furnish details about the file to the File Scheduler routine.
2. Check on the availability of the file to the program.
3. Rewind the tape, if necessary.
4. Process the tape label, if any.
5. Mark the file as "active."

These operations will be performed automatically for subsequent reels of multi-reel files. In addition, end-of-reel operations (rewinding, writing of tape marks, and writing of trailer labels (if any)) will be instituted.

## Error Messages

The following error messages will be produced during assembly under the conditions specified:

### OPERAND BLANK

If the operand is blank, a NOP will be generated instead of the calling sequence.

### PARAM $nn$ NOT A FILE

If an operand parameter (the number is indicated by  $nn$ , above) is not defined by a DTF entry, a NOP will be generated at the point in the calling sequence where a branch constant would normally be included. Since the IOC.IOPEN subroutine would consider this NOP to be the end of the list of files to be initialized by OPEN, a manual correction must be made before the object program is run. If corrections are not made and the file is not the last one named in the operand of the OPEN statement, the object program will execute the next branch constant as a true Branch instruction, thus transferring control (in error) to the first word of the corresponding DTF.

# GET

GET generates instructions to obtain a record for processing.

## Source Program Format

The basic formats for the GET statement are as follows:

Line	Label	Operation	OPERAND
3	56	15 16	20 21 25 30 35 40 45
0 1	ANY LABEL	GET	TAPEFILE
0 2	ANY LABEL	GET	TAPEFILE TO WORKAREA
0 3	ANY LABEL	GET	CARDFILE
0 4	ANY LABEL	GET	CARDFILE TO WORKAREA
0 5			

ANYLABEL is any symbolic label; it may be omitted. The entries GET and TO must be written exactly as shown. The first item in the operand must be the name of a tape file or unit record file. This name must be the same as the name which appears in the operand of the DTF or DUF entry which defines the file.

If the second or fourth format is used, the second operand item must be the word TO, preceded and followed by a single blank character. The third operand item must then be a name which appears as the label of a DRDW statement, the label of a DA header line, or the label of a DA subsequent entry.

## Processing Techniques

The calling sequence generated by each GET statement, in conjunction with the File Schedulers and other subroutines of the Input/Output Control System, make it possible for the object program to obtain each input record one at a time, regardless of the form for the input. The record will be made available in the input area and, if specified, moved to a work area. When the third item in the operand is the label of a DRDW, the record will be moved to the area defined by the generated RDW. When the third item in the operand is the label of a DA header line, the record will be moved into the area defined by the DA. If the DA header line does not specify the generation of an RDW, the processor will generate (elsewhere) an RDW to be used by the GET statement. When the third item in the operand is the label of a DA subsequent entry, the input record will be moved into the area defined by that entry. In this case, the processor will always generate an RDW to be used by the GET statement.

For the second and fourth formats, the amount of the input record moved is completely dependent on the size of the area defined by the third item in the operand. No warning message will be issued if the size of the work area is not equal to the size of the input record. The programmer must define the work area to equal the amount of the input record that is to be moved.

## Card Files

For card files, the GET statement using the third format causes the generation of the following instructions:

ANYLABEL	GET	CARDFILE
ANYLABEL	BLX	IOCSIXH, IOC.DN

The GET statement using the fourth format causes the generation of the following instructions:

ANYLABEL	GET	CARDFILE TO WORKAREA
ANYLABEL	BLX	IOCSIXH, IOC.CN
	B	WORKAREA

IOC.CN and IOC.DN are the labels of the two entry points to the unit record routine which is generated by the DUF entry that defines the file.

## Tape Files

For tape files, the use of the first or second format will cause the address of the first word of the next record to be placed into the indexing portion of an index word specified in the File Specifications Table of the input file. Processing may then be done by using instructions referring to fields within the record as defined by a DA subsequent entry relative to 0000 plus this index word. This indexing will be assigned automatically by writing the name of the index word in the DA header line operand as described on page 35. If processing is to be done by using non-indexed instructions, the record may be moved to a work area by means of the MOVE macro-instruction. However, the second format above will initialize the index word *and* move the record to the work area. With either the first or second format, reading in of the next block of records from tape when all records in the input area have been processed is automatic.

## Error and Warning Messages

The following error and warning messages will be produced during assembly under the conditions specified:

### OPERAND BLANK

If the operand is blank, a NOP will be generated instead of the calling sequence.

### OPERAND HAS TWO PARAMETERS

If there are two parameters in the operand, the calling sequence for the GET INPUTFILE form will be generated, followed by a NOP.

### PARAM 01 NOT A FILE

If the first item in the operand is not the name of a tape or card file, a NOP will be generated instead of the calling sequence.

### PARAM 01 NOT INPUT FILE

If a unit record file is named in the operand and if it is not an input file, this message will be produced. A NOP will be generated instead of the calling sequence.

### PARAM 03 IS A FILE. PARAM 03 IGNORED

If the third item in the operand is the name of a tape or unit record file, the calling sequence for the GET INPUTFILE form will be generated, followed by a NOP.

PARAM 03 NOT DEFINED

If the third item in the operand is not defined by a DA header line, DA subsequent entry, or a DRDW statement, the calling sequence for the GET INPUTFILE form will be generated, followed by a NOP.

SRBFORM4 BLANK, ASSUMED 10

If the file named in the operand specifies Form 4 records, and if the subrecord blocking factor is not specified, a subrecord blocking factor of 10 will be assumed. The calling sequence, however, will be generated in the normal manner.

WARNING—PARAM 01 NOT INPUT FILE

If a tape file is named in the operand and if it is not an input file, this message will be produced. However, the calling sequence will be generated in the normal manner.

WARNING—PARAM 02 IS NOT -TO-

If there are three parameters in the operand, and if the second parameter is not the word “to,” this message will be produced. The calling sequence will be generated as if the second parameter *had* been “to.”

WARNING — RLIFORM3 BLANK

If the file named in the operand specifies Form 3 records, and if the record length indicator is not specified, this message will be produced. It will be assumed that the record length indicator is located in position 0 of word 0 of the record. In all other respects, the calling sequence will be generated in the normal manner.

## PUT

PUT causes the generation of instructions that will provide the address of the next available word in the output area and, if desired, include a processed record in an output file by moving the record to the output area.

### Source Program Format

The basic formats for the PUT statement are as follows:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1	ANY LABEL	PUT	OUTAPEFILE
0 2	ANY LABEL	PUT	TAPEFILE IN OUTAPEFILE
0 3	ANY LABEL	PUT	WORKAREA IN OUTAPEFILE
0 4	ANY LABEL	PUT	FIELDNAME IN OUTAPEFILE
0 5	ANY LABEL	PUT	CARDFILE IN OUTAPEFILE
0 6	ANY LABEL	PUT	OUTCARDFILE
0 7	ANY LABEL	PUT	TAPEFILE IN OUTCARDFILE
0 8	ANY LABEL	PUT	WORKAREA IN OUTCARDFILE
0 9	ANY LABEL	PUT	CARDFILE IN OUTCARDFILE
1 0			

ANYLABEL is any symbolic label; it may be omitted. The entries PUT and IN must be written exactly as shown. The operand may contain either one or three parameters, as follows:

1. If the operand contains one parameter it must be the name of an output tape or unit record file; it must appear as the operand of the DTF or DUF entry which defines the file.
2. If the operand contains three parameters, the third parameter must be the name of an output tape or unit record file; it must appear as the operand of the DTF or DUF entry which defines the file. The second parameter must be the word IN, preceded and followed by a single blank character. The first parameter may be defined by appearing as a name in any of the following:
  - a. Operand of a DTF (the first format above).
  - b. Operand of a DUF (the sixth format above).
  - c. Label of a DA header line.
  - d. Label of a DA subsequent entry.
  - e. Label of a DC header line.
  - f. Label of a DC subsequent entry.
  - g. Label of a DLINE header line.
  - h. Label of a DLINE subsequent entry.
  - i. Label of a DRDW. (The area defined will be included in the output file.)

### Processing Techniques

The calling sequences generated by each PUT statement, in conjunction with the File Schedulers and other subroutines of the Input/Output Control System, make it possible for the object program to cause the inclusion of each processed



record or field in the output file one at a time, regardless of the output blocking factor. If the name of the item to be included is defined by a declarative statement, and, if an RDW(s) is not specified for that item, the processor will generate (elsewhere) an RDW to be used by the PUT statement.

The use of the first or sixth format will cause the address of the next available word in the output area to be placed in the indexing portion of an index word specified in the File Specifications Table of the output file. Data may then be included in the output area by means of a later MOVE statement. (It may also be processed there until the next PUT occurs.) The other formats, however, not only cause the address of the next available word to be placed in the index word specified, but also cause automatic inclusion of the record in the output file and updating of the proper index words.

Writing of the output area will occur automatically when the area is full.

## **Error and Warning Messages**

The following error and warning messages will be produced during compilation under the conditions specified:

### **OPERAND BLANK**

If the operand is blank, a NOP will be generated instead of the calling sequence.

### **OPERAND HAS TWO PARAMETERS**

If there are two parameters in the operand, a NOP will be generated instead of the calling sequence.

### **OUTPUT SRBFORM4 BLANK, ASSUMED 10**

If the output tape file named in the operand specifies Form 4 records, and if the subrecord blocking factor is not specified, a subrecord blocking factor of 10 will be assumed. The calling sequence, however, will be generated in the normal manner.

### **PARAM 01 (03) NOT A FILE**

If the parameter named as the output file is not a tape or unit record file, a NOP will be generated instead of the calling sequence.

### **PARAM 01 (03) NOT OUTPUT FILE**

If the parameter named as the output file is a unit record file, and if it is not defined by its DUF as an output file, this message will be produced. A NOP will be generated instead of the calling sequence.

### **PARAM 01 UNDEFINED**

If there are three parameters in the operand, and if the first parameter is not defined by one of the nine entries named under "Source Program Formats," this message will be produced. The calling sequence for the PUT OUTPUTFILE form will be generated, followed by a NOP.

### **WARNING — OUTPUT RLIFORM3 BLANK**

If the output tape file named in the operand specifies Form 3 records, and if the record length indicator is not specified, this message will be produced. It will be assumed that the record length indicator is located in position 0 of word 0 of the record, and the calling sequence will otherwise be generated in the normal manner.

**WARNING — PARAM 01 (03) NOT OUTPUT FILE**

If the parameter named as the output file is a tape file, and if it is not defined by its DTF as an output file, this message will be produced. However, the calling sequence will be generated in the normal manner.

**WARNING — PARAM 02 IS NOT -IN-**

If there are three parameters in the operand, and if the second parameter is not the word "IN," this message will be produced. The calling sequence will be generated as if the second parameter had been "IN."

# PUTX

PUTX causes the generation of instructions that will include a processed record in an output file by exchanging `rdws` rather than by moving the record.

## Source Program Format

The basic format for the PUTX statement is as follows:

Line	Label	Operation	OPERAND
3	5	15	20
3	6	16	21
		25	30
		35	40
			45
0,1	ANYLABEL	PUTX	INPUTFILE IN OUTPUTFILE
0,2			

ANYLABEL is any symbolic label; it may be omitted. The entries PUTX and IN must be written exactly as shown.

The first item in the operand must be the name of an input tape file. This name must appear in the operand of the DTF entry which defines the file.

The second item in the operand must be the word IN, preceded and followed by a single blank character.

The third item in the operand must be the name of an output tape file. This name must appear in the operand of the DTF entry which defines the file.

## Processing Techniques

The calling sequences generated by each PUTX statement, in conjunction with the File Schedulers and other subroutines of the Input/Output Control System, make it possible for the object program to cause the inclusion of each processed record in the output file, one at a time, regardless of the output blocking factor. Unlike the PUT statement, however, the PUTX statement causes this inclusion by the interchange of `rdws`; the record itself is not moved. Thus, an `rdw` describing the record to be written is placed in the list of output `rdws` and the `rdw` which was previously at that point in the list is placed back in the input list, replacing the original `rdw`.

The form of the records in the files places the following restrictions on the use of the PUTX macro-instructions:

1. Form 3 records *can not* be processed with a PUTX macro-instruction.
2. The combination of input record form and output record form must be one of the following:

<i>Input File Record Form</i>	<i>Output File Record Form</i>
1	1
1	2
2	1
2	2
4	4

3. The length of fixed length records or the maximum length of variable length records must be identical for both the input and output files.
4. For Form 4 records, the number of sections in input and output records must be the same and the maximum number of words in the corresponding sections of each input and output record must be identical.

There is no restriction on the blocking factor of the input and output files. The blocking factor, i.e., the number of records in one block, may be different for each of the files provided that the other restrictions listed above are observed.

After `RDWS` have been exchanged by a `PUTX` macro-instruction, the input record is no longer available for processing; the programmer must be certain that all processing requiring the input data is completed before issuing the command.

When `PUTX` is to be used, processing should be done in the input area using indexed instructions which refer to fields within the record as defined by a `DA` entry relative to 0000. If input data is moved to a work area for processing, do not use `PUTX`. If `PUTX` is used, the original input data rather than the results of processing will appear in the output file.

The automatic function of writing blocks of records on tape is the same for the `PUTX` macro-instruction as for the `PUT` macro-instruction.

## Error and Warning Messages

The following error and warning messages will be produced during compilation under the conditions specified:

### IMPROPER OPERAND

If the operand does not contain three parameters, a `NOP` will be generated instead of the calling sequence.

### PARAM 01 (03) FILE FORM INVALID

If the files named in the operand do not conform to the restrictions regarding record form which are listed under "Processing Techniques," a `NOP` will be generated instead of the calling sequence.

### PARAM 01 (03) NOT A FILE

If either the first or third item in the operand is not defined by a `DIF`, a `NOP` will be generated instead of the calling sequence.

### PARAM 03 — SRBFORM4 BLANK, ASSUMED 10

If the output tape file named in the operand specifies Form 4 records, and if the subrecord blocking factor is not specified, a subrecord blocking factor of 10 will be generated. The calling sequence, however, will be generated in the normal manner.

### RECLENGTHS UNEQUAL

If the record lengths specified for the tape files named in the operand are not equal, this message will be produced. The calling sequence will be generated in the normal manner, however.

### SRBFORM4 UNEQUAL. OUTPUT SRB USED

If the tape files named in the operand specify Form 4 records, and if the subrecord blocking factors of the files are not equal, this message will be produced. The subrecord blocking factor of the output file will be used in the calling sequence.

# CLOSE

CLOSE generates instructions to remove tape files from use.

## Source Program Format

The basic format of the CLOSE statement is as follows:

Line	Label	Operation	OPERAND
3	5	15	16 20 21 25 30 35 40 45
0.1	ANY LABEL	CLOSE	FILEA, FILEB, FILEC, etc.
0.2			

ANYLABEL is any symbolic label; it may be omitted. The entry CLOSE must be written exactly as shown. The operand must contain the name(s) of one or more tape files to be removed from processing use. Each name must be the same as the name which appears in the operand of the DTF entry which defines the file. As many tape files as desired may be named in the operand of a CLOSE statement, subject only to the restriction that the operand may not be extended over more than four continuation cards. The names of the tape files must be separated by commas.

## Processing Techniques

The first instruction generated by the processor as a result of a CLOSE statement is as follows:

Line	Label	Operation	OPERAND
3	5	15	16 20 21 25 30 35 40 45
0.1	ANY LABEL	BLX	IOCSIXG, IOC.ICLOSE
0.2			

Following this, the processor will generate a branch constant containing the address of the first word of the File Specifications Table for each file named in the operand. A NOP will be generated after the last branch.

This calling sequence and the subroutine IOC.ICLOSE (normally included as a result of a DIOCS entry) will perform the following for each output file:

1. Write out remaining records in the output area(s).
2. Write a tape mark.
3. Write end-of-file trailer labels (if desired).

In addition, the following will be performed for all files:

1. Rewind if necessary.
2. Mark the file as "inactive."

The CLOSE statement will normally be used to cause these operations to be performed for the *last* reel of each file.

## Error Messages

The following error messages will be produced during assembly under the conditions specified:

### OPERAND BLANK

If the operand is blank, a NOP will be generated instead of the calling sequence.

### PARAM $nn$ NOT A FILE

If an operand parameter (the number is indicated by  $nn$ , above) is not defined by a DTF entry, a NOP will be generated at the point in the calling sequence where a branch constant would normally be included. Since the IOC.ICLOSE subroutine would consider this NOP to be the end of the list of files to be removed by CLOSE, a manual correction must be made before the object program is run. If corrections are not made and the file is not the last one named in the operand of the CLOSE statement, the object program will execute the next branch constant as a true Branch instruction, thus transferring control (in error) to the first word of the corresponding DTF.

## END

END generates instructions to remove tapes from use, type an end-of-job message, and then branch, halt, or permit SPOOL operations.

### Source Program Format

The basic format for the END statement is as follows:

Line	Label	Operation	OPERAND
3	5	15	16 20 21 25 30 35 40 45
0.1	ANY LABEL	END	BRANCHADDR
0.2			

ANYLABEL is any symbolic label; it may be omitted. The entry END must be written exactly as shown. The operand may be blank, or it may contain an actual or symbolic address.

### Processing Techniques

The END statement will cause the generation of one of the following calling sequences:

```
ANYLABEL    END
ANYLABEL    BLX    IOCSIXG, IOC.IEND
              NOP    0
              OR
ANYLABEL    END
ANYLABEL    BLX    IOCSIXG, IOC.IEND
              B      BRANCHADDR
```

These instructions and the IOC.IEND subroutine (normally included as a result of a DIOCS statement) will initiate the operations generally performed by the CLOSE macro-instruction for all files for which this has not yet been done, type an end-of-job message, and cause a branch to BRANCHADDR (if one is named).

If the branch address is omitted, the generated instructions will perform the necessary CLOSE operations; thereafter, the instructions will cause one of the following:

1. A halt if no SPOOL program is run in conjunction with the main program.
2. A program loop to be entered to permit SPOOL programs to continue if any are being run. Loading of another main program can then be initiated on signal from the SPOOL routine.

## ARITH — Arithmetic Operations

ARITH generates instructions to compute the value of an arithmetic expression and to store the result in any desired field.

### Source Program Format

The basic formats for the ARITH statement in the source program are as follows:

Line	Label	Operation	OPERAND
3	56	1516 2021	25 30 35 40 45
0 1	ANY LABEL	ARITH	RESULT = EXPRESSION
0 2	ANY LABEL	ARITH	RESULT = EXPRESSION, OVERFLOWBR
0 3			

ANYLABEL is any symbolic label; it may be omitted. The entry ARITH and the equal sign must be written exactly as shown. The equal sign indicates that the RESULT field is to be set equal to the value of the arithmetic EXPRESSION. RESULT may be any symbolic name; the various allowable forms of the EXPRESSION are described under "Arithmetic Expressions," below. OVERFLOWBR is the symbolic label of the first instruction of a routine to which the program is to branch in the event of an overflow.

### Arithmetic Expressions

Arithmetic expressions are formed from *numerical quantities* which may be in one of several *modes*. In addition, *arithmetic operators* or *functions* may operate on the quantities and *arithmetic punctuation* may establish an execution sequence.

### Numerical Quantities

The numerical quantities that enter the computation may be of several kinds, namely symbolic fields (representing variables), literals, or defined constants. Fields referenced by symbolic names will generally be defined under a DA or DC header line, with their mode and (where applicable) format specified. Literals may be signed or unsigned; if unsigned, they will be interpreted as positive. Adcons will be treated as four-place integers in the automatic-decimal mode.

### Arithmetic Operators

The ARITH statement interprets two types of arithmetic operators, unary and binary. A unary operator identifies or changes the sign of *one* numerical quantity; a binary operator indicates an operation to be performed upon *two* quantities to form another.

The following two *unary* operators are written preceding the quantity upon which they are to operate:

Operator	Operation
& or +	The ampersand or plus sign preceding a literal identifies it as positive. When either precedes another quantity, it does not affect its value.
—	A minus sign preceding a positive quantity changes its sign to minus; a minus sign preceding a negative quantity changes its sign to plus

The following five *binary* operators are written between the two quantities upon which they are to operate:



<i>Operator</i>	<i>Operation</i>
& or +	The ampersand or plus sign indicates that the two numbers are to be added.
—	The minus sign indicates that the second number is to be subtracted from the first number.
*	The asterisk indicates that the two numbers are to be multiplied.
/	The slash indicates that the first number is to be divided by the second number.
**	Double asterisks indicate exponentiation; i.e., the first number is to be raised to the power indicated by the second number.

The use of arithmetic operators preserves the mode of the operands.

## **Modes**

Computation can be carried out in either automatic-decimal or floating-decimal mode. All numerical quantities appearing in *EXPRESSION* should be in the same mode. Two exceptions (4 and 5, below) are included in the following rules which apply to modes within *EXPRESSION*:

1. If A is a numerical quantity, then +A, -A, and ABS(A) are quantities of the same mode as A.
2. Enclosing a quantity(ies) in parentheses does not change the mode of the quantity(ies)
3. Adding, subtracting, multiplying, or dividing two quantities of the same mode will give a result in the same mode.
4. If A and B are numerical quantities, then A\*\*B will be a permissible expression of the same mode as A, regardless of the mode of B. The permissible combinations of A and B and the resultant mode of A\*\*B are indicated in the following table:

<i>Mode of A</i>	<i>Mode of B</i>	<i>Mode of A**B</i>
automatic-decimal	automatic-decimal	automatic-decimal
automatic-decimal	floating-decimal	automatic-decimal
floating-decimal	automatic-decimal	floating-decimal
floating-decimal	floating-decimal	floating-decimal

5. Arguments of functions may be in a different mode from that of the rest of *EXPRESSION*. (For details, see "Functions," below.)

The *RESULT* field may be in a different mode from that of *EXPRESSION*. In such a case, coding will be generated to edit the answer to the mode and format of the *RESULT* field.

The *ARITH* macro generator will not accept fields in the alphameric mode as operands. When an alphameric field, either literal or symbolic, is encountered, an error message will be issued and a *NOF* instruction generated.

## **Functions**

During the computation of an arithmetic *EXPRESSION*, functions may be evaluated provided that appropriate subroutines are available. Such subroutines must be in the Program Library, even when they are specially provided by the programmer. (The function names used in the various examples of the *ARITH* macro-instruction are intended for illustration only. Their use should not imply that subroutines evaluating these functions are being furnished.)

The maximum number of arguments (independent variables) allowed for a function is equal to 69 minus the *total* number of parameters in the source statement. If a function is included in the source statement and this function with its arguments are the *only* parameters in the statement, then a maximum of 34 arguments would be allowed for the function. This is deduced since 34 arguments plus 1 function name equals 35 parameters and 69 minus 35 equals 34 arguments that will be accommodated. Under the same circumstances, if the function had 35 arguments, then 35 plus the 1 function name equals 36 parameters and 69 minus 36 equals only 33 arguments which can be accommodated. The last 2 arguments would be ignored and a message would be issued. If the source statement contains the maximum number of parameters permitted, i.e., 50, then a function included in the statement would be allowed only 19 arguments. Arguments may be considered to be either single variables (A, b, X, y) or expressions (a+b, 2\*c/d-e, X\*\*Y).

The arguments to be used by a function are enclosed by a set of parentheses and written following the function name. A comma is written following each argument but the last. An example of a function followed by a series of five arguments is as follows:

```
FUNCTION(A, b-2, X*Y, C**d, z)
```

ARITH will generate a call for the subroutine specified by the symbolic name FUNCTION, furnish the arguments to this subroutine and generate instructions for the computation to be continued, if necessary, after the function has been evaluated.

Function arguments constitute one of the exceptions to the rule that all fields within EXPRESSION must be in the same mode. For functions, this rule is modified as follows:

1. The argument(s) of any function must be in the proper mode in order for the specific subroutine called in to be able to operate upon it.
2. The function value must be in the proper mode so that the next computational step can be carried out.

The second condition states specifically that function values need not be in the mode of the EXPRESSION as a whole, since it is possible that the value obtained from one function will serve as an argument for another. When functions are nested in this manner, only the value of the outermost function must be in the mode of the EXPRESSION; the inner functions must produce values in a mode acceptable to the next function towards the outside. For example, if the expression  $A = \text{ARCTAN}(\text{LOG}_e(x))$  is to be evaluated, an Arctan function routine must be available whose output is in the same mode as A. A Log<sub>e</sub> function must be available whose output is in a mode that the Arctan function can accept as an argument, and x must be in a mode that Log<sub>e</sub> can accept as an argument. The result will be in the mode of A.

If a function has an expression as an argument, coding will be generated to compute this expression and to furnish the result to the function subroutine. Such computation, as well as arithmetic operations on function values prior to their being used as arguments of other functions, will preserve the original mode. For example, if  $3 - \text{SIN}(3.1415 + \text{ARCTAN}(1 - x))$  is to be computed, x must be in the automatic-decimal mode because 1 is in the automatic-decimal mode. The Arctan function must take arguments and yield function values in the same mode, and the same must be true of the Sine function. The result will then be automatic-decimal.

The absolute value function is a special function which is provided on the Compiler Systems Tape. This function is the only one whose instructions are generated in-line. The ABS function preceding a quantity indicates that the sign of the quantity is to be plus. Thus, a negative quantity is changed to a positive quantity; a positive quantity is unaffected. The quantity in question must be enclosed by parentheses.

### **Arithmetic Punctuation and Execution Sequence**

Arithmetic expressions using more than one operator may be ambiguous unless the order in which the operations are to be carried out is indicated. For example,  $12-2-3$  yields 7 if the expression is interpreted to be  $(12-2)-3$  and yields 13 if interpreted to be  $12-(2-3)$ . It is therefore necessary to establish an execution sequence of arithmetic operations by means of certain punctuation rules.

The ARITH statement will use parentheses in the customary way: expressions enclosed by parentheses are to be computed before they can be used as a component in the next operation. For nested parentheses (parentheses within parentheses) the same rule holds; the expression(s) enclosed by a greater number of parentheses will be computed before those enclosed by fewer parentheses.

To reduce the great number of parentheses which might be introduced by the explicit punctuation of all EXPRESSIONS, certain conventions are introduced. The following order of preference for operations is established:

1. Exponentiation
2. Unary Operators
3. Multiplication and Division
4. Addition and Subtraction

A sequence of operations will be carried out in this order unless parenthesization intervenes or directs otherwise.

If several operators of the same order are present, they will be executed from left to right. For example,  $A-B-C-D$  is treated as  $((A-B)-C)-D$ ; similarly, for multiplication and division,  $A/B/C*D$  is treated as  $((A/B)/C)*D$ .

Repeated exponentiation is also executed from left to right unless punctuated otherwise. For example,

$2^3$  is interpreted to mean  $(2^3)^2$  or  $8^2$ , not  $2^{(3^2)}$  or  $2^9$ .

In other words,  $A**B**C$  is treated as  $(A**B)**C$  by the ARITH macro generator.

The only exception to this "left-to-right" rule applies to the addition and/or subtraction of automatic-decimal fields. The fields are taken in order of increasing number of decimal places rather than from left to right. This eliminates shifting the accumulator to the right which avoids loss of digits in this direction. Consequently, shifting will be only to the left; possible error is standardized to leftward overflow, which can be dealt with according to the procedures described under "Overflow Branch."

This execution order for operations is followed for expressions enclosed by the most parentheses, in order of decreasing parentheses, until the value of the entire EXPRESSION has been computed.

If these conventions are taken into account, a large number of parentheses may be omitted.

The generator will also accept statements in which unnecessary but correctly placed parentheses exist, but it is in the interest of storage economy to leave out parentheses wherever possible without introducing ambiguity. In no case may the number of parentheses (other than those used in address modification to enclose any one numerical field) exceed fifty. Since this limitation is identical to the number of permitted parameters, enough parentheses are allowed to punctuate any admissible expression. If the maximum parenthesis level is exceeded and the generator is unable to process the macro-instruction, an error message will be issued.

The following examples illustrate the use of the punctuation rules and the resultant sequence of execution:

EXPRESSION:	$A-B*C$
Interpretation:	$A+((-B)*C)$
Execution Sequence:	$-B$ and $C$ are multiplied; the result is added to $A$ .
EXPRESSION:	$(A-B)*C$
Interpretation:	Same as EXPRESSION.
Execution Sequence:	$B$ is subtracted from $A$ ; the result is multiplied by $C$ .
EXPRESSION:	$ABS(A)-B**C$
Interpretation:	$(ABS(A))- (B**C)$
Execution Sequence:	$B$ is raised to the power $C$ ; the absolute value of $A$ is taken. Finally, the first result is subtracted from the second.
EXPRESSION:	$A+B*C-D/3$
Interpretation:	$A+(B*C)-(D/3)$
Execution Sequence:	$B$ and $C$ are multiplied; $D$ is divided by $3$ ; the necessary addition and subtraction operations are performed.
EXPRESSION:	$A*1.2-SIN(X+3)/3+Y**4$
Interpretation:	$(A*1.2)-((SIN(X+3))/3)+(Y**4)$
Execution Sequence:	$X$ and $3$ are added and the sine of the sum taken. $Y$ is raised to the 4th power; $A$ is multiplied by $1.2$ ; the sine is divided by $3$ . Finally, the necessary addition and subtraction of the intermediate results are performed.

## Processing Techniques

### Limitations on Length

The number of permissible parameters is fifty. Parameters are considered to be the following:

1. Numerical fields (including function arguments).
2. Names of functions, including `ABS`.
3. Overflow branch, if specified.

Arithmetic operators and punctuation are *not* counted as parameters. An attempt to write more than the permitted number of parameters in the operand of an `ARITH` statement will be intercepted by the processor and no coding will be generated.

### Spacing and Punctuation

No blanks should appear in the operand of an `ARITH` macro-instruction. Only one equal sign may appear, positioned as shown under "Source Program Format."

Commas must separate function arguments; one additional comma is required to separate `EXPRESSION` from `OVERFLOWBR` if the second format is chosen. The appearance of illegal characters in the operand will cause an error message to be issued.

### Address Modification

Symbolic addresses may be modified by indexing and address adjustment.

### Overflow Branch

Since the conditions under which accumulator overflow could take place during a computation are extremely varied, no general routine for dealing with overflow has been provided. Instead, opportunity is given to the programmer to supply his own correction routine, which can be suited to his particular `ARITH` statement. If an overflow branch is specified, the generated instructions will contain a `BLX` instruction to cause a transfer to the overflow routine under the following conditions:

1. *Automatic-Decimal Computations:* Overflow resulting from addition, subtraction, exponentiation, store, and add-to-storage operations. (Transfer will *not* be caused by multiplication and division overflow. However, warning messages will be issued.)
2. *Floating-Decimal Computations:* Exceeding the maximum value for floating-decimal numbers as a result of *any* operation. This transfer will be made whether the object program machine has floating-decimal hardware or not. In the first case, the floating-decimal overflow indicator is tested; in the second, this test is simulated along with the floating-decimal arithmetic procedures.

Following the overflow routine, the program will return to the instruction in the object program following the `BLX`, provided that the last instruction in the overflow routine is an unconditional Branch to location `0000 + X94`.

If no overflow branch is indicated, warning messages will be issued during automatic-decimal arithmetic assemblies pointing out the possibility of overflow.

Since the possibility of overflow in floating-decimal computations cannot be detected on the basis of format alone, warning messages will never be issued during floating-decimal arithmetic assemblies.

If `OVERFLOWBR` is indicated in the operand of the `ARITH` statement but, because of the input formats involved, overflow cannot possibly occur, then no overflow branch will be generated. This is illustrated in example 4 under "Examples."

### Mode Size for Automatic-Decimal Computations

In handling automatic-decimal numbers, the processor establishes a mode size which is either ten or twenty digits in length and is based on the input formats of all the fields in the `ARITH` source statement. Then, based on the mode size which has already been determined and on the input formats, a "computation mask" or standard format is established. The mask indicates the decimal point placement and the maximum number of digits which may appear to the left and to the right of the decimal point. No intermediate results will be permitted to exceed this format.

Symbols have been established for this discussion and are defined as follows:

<i>Symbol</i>	<i>Definition</i>
<code>MS</code>	The mode size for automatic-decimal computations.
<code>MDL</code>	The maximum number of digits to the left of the decimal point (i.e., integer digits) as specified by the input formats in the <code>ARITH</code> source statement (both <code>EXPRESSION</code> and <code>RESULT</code> fields).
<code>MDR</code>	The maximum number of digits to the right of the decimal point (i.e., decimal digits) as specified by the input formats in the <code>ARITH</code> source statement (both <code>EXPRESSION</code> and <code>RESULT</code> fields).

<i>Symbol</i>	<i>Definition</i>
DL	The number of digits to the left of the decimal point (integer digits) in a mask.
DR	The number of digits to the right of the decimal point (decimal digits) in a mask.

The *mode size* for automatic-decimal computation is established as follows. If MDL exceeds 20, i.e., if any of the input fields has more than twenty integer digits, an error condition results; error message N 19 is issued and no coding other than a NOP will be generated.

If MDL is less than or equal to 20, the mode size depends on the sum of MDL and MDR, as follows:

$MDL+MDR$	$MS$
$\leq 10$	10
$> 10$	20

Once mode size has been established, the *mask* is defined. The procedure followed again depends on the value of MDL+MDR:

1. If MDL+MDR is less than or equal to 20, DR will be set equal to MDR, and DL will be set equal to MS-DR.
2. If MDL+MDR exceeds 20, DL will be set equal to MDL; DR will be set equal to MS-DL. Enough decimal digits will be truncated (without rounding) to reduce the overall length to twenty digits. A warning message (W 18) will include the computation mask in four-digit form, the first two representing the number of integers, the last two the number of decimals.

In short, if the total number of digits in all the input fields, with their decimal points aligned, is less than the computed mode size, the extra capacity is applied to the integer side. If the total exceeds 20, high-order digits are protected and decimal places truncated.

The following examples illustrate the methods of establishing computation masks:

1. Ten-digit mask with extra integer capacity.

<i>Input Formats</i>	MDL = 6	DR=MDR=3
5.2	MDR = 3	DL=MS-DR=10-3=7
6.1	Sum = 9	
3.3	MS = 10	<i>Mask: 7.3</i>

2. Ten-digit mask without extra capacity.

<i>Input Formats</i>	MDL = 8	DR=MDR=2
7.1	MDR = 2	DL=MS-DR=10-2=8
8.1	Sum = 10	
3.2	MS = 10	<i>Mask: 8.2</i>

3. Twenty-digit mask with extra integer capacity; input fields do not exceed ten digits.

<i>Input Formats</i>	MDL = 8	DR=MDR=7
2.3	MDR = 7	DL=MS-DR=20-7=13
8.2	Sum = 15	
1.7	MS = 20	<i>Mask: 13.7</i>

4. Twenty-digit mask with extra integer capacity; input fields exceed ten digits.

<i>Input Formats</i>	MDL = 9	DR=MDR=8
9.2	MDR = 8	DL=MS-DR=20-8=12
7.8	Sum = 17	
3.1	MS = 20	Mask: 12.8

5. Twenty-digit mask without extra capacity; input fields exceed ten digits. (This case could also occur with input fields not exceeding ten digits if two of them had formats of 10.0 and 0.10, respectively; the mask would then be 10.10.)

<i>Input Formats</i>	MDL = 4	DR=MDR=16
3.16	MDR = 16	DL=MS-DR=20-16=4
4.9	Sum = 20	
1.4	MS = 20	Mask: 4.16

6. Twenty-digit mask with decimal digits truncated.

<i>Input Formats</i>	MDL = 13	DL=MDL=13
13.2	MDR = 9	DR=MS-DL=20-13=7
11.8	Sum = 22	
4.9	MS = 20	Mask: 13.7

Two decimal places will be lost. Warning message W 18 will give the computation mask in the form 1307.

During computation, all intermediate results will be confined to the mask. Excess decimal digits developed will be truncated without rounding or warning. Excess integer digits may also be lost; if this becomes possible, the consequences will depend on the type of operation that caused the difficulty.

*Addition or Subtraction.* If an overflow branch has been specified, coding will be generated to transfer the object program to this branch if necessary. If no overflow branch is indicated, warning message W 21 or W 22 will be issued during assembly. Warning messages are issued if overflow or digit loss is *possible*, as determined on the basis of field *format* alone; the transfer to the overflow branch takes place only when these conditions become *actual* due to the specific object-time contents of the fields.

*Multiplication or Division.* If integer digits may possibly be lost during multiplication and division operations, warning message W 20 will be issued during assembly. Overflow resulting from multiplication or division will not cause a transfer to the overflow branch.

If the divisor field in a division operation is defined as having integer digits, the high-order digit is significant. In other words, if the automatic-decimal format of a divisor is 2.4, the generator will proceed on the assumption that the contents of the field at object program time will be *at least* 10.0000. If this condition is not satisfied, intermediate results may exceed the computation mask without a warning at assembly time.

If the divisor has an automatic-decimal format of the type 0.n, it will merely be assumed that the nth decimal digit contains at least a 1; this assumption is, of course, non-restrictive, since zero divisors lead to special procedures as described under "Zero Divisors."

When a warning message is issued, the programmer should check whether his intermediate results can exceed the mask on the left; this will depend on his

actual data as well as the defined field formats and the specified arithmetic operations. If overflow can occur, an input field can be redefined so as to (1) change a ten-digit computation to twenty-digit mode size, or (2) increase the number of integer digits at the expense of decimal digits. To avoid recurrence of the same problem after redefinition, it is generally advisable to modify the defined format of the result field where the program objectives permit.

*Exponentiation.* The exponentiation of all numbers (automatic-decimal or floating-decimal bases having either automatic-decimal or floating-decimal exponents) is carried out by means of floating-decimal routines. Automatic-decimal bases and exponents are converted to floating-decimal numbers and exponentiation is carried out by means of a subroutine. The result will be converted to an automatic-decimal number, if required. Three exceptions to this process are as follows:

1. If the base is an automatic-decimal integer and the exponent is an automatic-decimal integer less than ten digits in length, exponentiation is carried out by means of a subroutine which generates the required multiplication instructions.
2. If the base is an automatic-decimal integer and the exponent is a literal 2 or 3, Multiply instructions are generated in-line.
3. If the base is a floating-decimal number and the exponent is a literal 2 or 3, Floating Multiply instructions are generated in-line.

If the result of the exponentiation should exceed the computation mask format on the left, transfer will be made to an overflow branch when one is specified. Otherwise, the overflow latch for Accumulator 1 will be set ON. No warning message will be issued at assembly time (except for cases 1 and 2, above), since the size of the result cannot be predicted on the basis of floating-decimal field formats.

### **Zero Divisors**

Before a division is executed, the divisor is tested for zero. If the divisor is zero, the overflow latch for Accumulator 3 is set ON. Subsequent procedures depend upon the mode of computation, as follows:

1. *Automatic-Decimal Computations:* The accumulator(s) containing the quotient will be filled with 9s and given the proper sign determined by the signs of the dividend and divisor. Computation will then continue as usual.
2. *Floating-Decimal Computations:* The division is ignored; i.e., the dividend is used as quotient. Computation will then continue as usual.

No other operation initiated by the ARITH macro-instruction alters the overflow latch for Accumulator 3, with the possible exception of function subroutines, whose individual specifications may be consulted. This latch, therefore, provides a certain test as to whether division by zero has been attempted. If such a test is desired, the latch must be set OFF before the ARITH macro-instruction is executed, since no automatic provision is made for this by the generator.

### **Final Storage**

If the RESULT field is in a mode different from that of the computation, editing will be necessary before final storage. The following rules apply:

1. In editing from automatic-decimal to floating-decimal mode, the result will appear in normalized form. Only the first eight significant digits will be converted; further digits to the right will be truncated without rounding.
2. In editing from floating-decimal to automatic-decimal mode, the following three cases are distinguished:



- a. If, after conversion, the first significant digit falls to the left of the high-order digit of the `RESULT` field, an overflow condition exists. No warning message can be issued at assembly time since this condition cannot be predicted on the basis of floating-decimal field format alone. At object time, a message "SHIFT OUT OF RANGE, FLT TO DECI" will be typed out. Transfer will be made to the overflow branch if one is specified; such digits as can be accommodated in their proper places will be stored. If no overflow branch is specified, the overflow latch of Accumulator 1 will be set ON.
- b. The "normal" case exists if, after conversion, the first significant digit falls into one of the digits of the `RESULT` field. Excess decimal digits developed beyond the capacity of the `RESULT` field will be truncated after rounding.
- c. If, after conversion, the first significant digit falls to the right of the low-order digit of the `RESULT` field, then the decimal value of the answer is too small to register in the established format. At object time, a message "SHIFT OUT OF RANGE, FLT TO DECI" will be typed out. The `RESULT` field will be set to zero.

If the result of an automatic-decimal computation is to be stored in an automatic-decimal `RESULT` field of smaller format, the following procedures are followed:

1. Excess decimal digits are truncated *after* rounding; in other words, if the digit immediately to the right of the point of truncation contains a value of 5 or more, the next digit to the left is increased by 1.
2. Excess integer digits are lost and an overflow condition results. If an overflow branch is specified, transfer will be made to the overflow routine at object program time. If no overflow branch is indicated, warning message W 23 is issued during assembly.

### Setting Overflow Lights

The overflow lights for Accumulators 1, 2, and 3 should be set as follows for both automatic-decimal and floating-decimal computations. The exponent overflow light should be set as indicated for floating-decimal computations if floating hardware is used.

*Accumulator 1.* If the accumulator 1 overflow light is *not* ON and Accumulator 1 overflows, the machine will stop.

If the accumulator 1 overflow light is ON, one of the following provisions should be made:

1. An overflow branch in the operand of the `ARITH` macro-instruction.
2. A `bv1` (Branch if Overflow in Accumulator 1) instruction following each `ARITH` without an overflow branch.

If neither of the above provisions is made and the accumulator overflows, the condition will be carried and an error will be introduced.

*Accumulator 2.* The accumulator 2 overflow light must be ON during object program time.

*Accumulator 3.* The accumulator 3 overflow light must be ON if the detection of an attempt to divide by zero is desired.

*Exponent.* If the exponent overflow light is *not* ON and the exponent overflows, the machine will stop.

## Error and Warning Messages

If the exponent overflow light is ON, one of the following provisions should be made:

1. An overflow branch in the operand of the ARITH macro-instruction.
2. A FBV (Floating Branch Overflow) instruction following each ARITH without an overflow branch.

If neither of the above provisions is made and the exponent overflows, the condition will be carried and an error will be introduced.

Under the conditions specified, the ARITH macro generator will issue the following error and warning messages during assembly. Unlike the other macro generators, ARITH does not give the text of the message. Only the message code letter and number, possibly supplemented by a parameter number or computation-mask format, are given. The programmer must refer to the list below for text and interpretation. The code letters are to be interpreted as follows:

<i>Code</i>	<i>Interpretation</i>
N	An error condition exists that makes further coding impossible; a NOP has been generated.
W	A warning that either an unusual condition or the possibility of error exists; generation continues.
X	An error condition exists; generation will continue on the special assumptions stated in the message.

### N 01 NO OPERAND

No numerical field has been specified upon which an operation is to be performed.

### X 02 NO EQUAL SIGN — WILL NOT STORE RESULT

The RESULT field and the equal sign have both been omitted. The generator has produced instructions to compute the value of the arithmetic EXPRESSION, but not to store the result.

### X 03 NO RESULT FIELD — CANNOT STORE RESULT

The operand portion of the macro-instruction begins with the equal sign. The generator has produced instructions to compute the value of the arithmetic EXPRESSION, but not to store the result.

### N 04 ALPHA FIELD UNACCEPTABLE. PARAMETER xx

The parameter number of the alphameric field has been included in the message.

### X 07 INCOMPLETE — WILL PROCESS TO PARAMETER xx

The text of the input statement appears to be broken off; e.g., it ends with a left parenthesis. This condition may also occur if one blank precedes an entry that should be processed on the same card. The number of the last parameter processed has been included in the message.

### X 08 TEXT ENDS WITH OPERATOR — WILL IGNORE

The last operator does not have an operand on its right and has been ignored.

### X 09 CONSECUTIVE OPERATORS — WILL ACCEPT FIRST ONLY BEFORE PARAMETER xx

Two successive arithmetic operators have been detected by the scan, pre-

ceding the parameter whose number has been included in the message. The second of these operators has been ignored.

X 10 NO PUNCTUATION — WILL IGNORE PARAMETER *xx*

Two consecutive numerical fields have been detected, without an intervening operator. The second field, whose parameter number has been included in the message, has been ignored.

X 11 ILLEGAL CHARACTER — WILL BE IGNORED. BEFORE PARAMETER *xx*

The scan has detected a character that is not one of the allowable punctuation marks or arithmetic operators. The number of the next following parameter has been included in the message to aid in locating the faulty entry.

X 12 EXCESS RIGHT PARENTHESES — WILL IGNORE

An attempt has been made to close more parentheses than had been opened. All right parentheses unmatched by left parentheses have been ignored.

X 13 END OF SCAN — PARENTHESES DO NOT MATCH — WILL SUPPLY

Some parentheses have been left open. Generation of instructions have proceeded on the assumption that they are all to be closed at the right end of the EXPRESSION.

X 16 FUNCTION WITHOUT ARGUMENT — PROCESSING STOPS BEFORE PARAMETER *xx*

A function symbol, whose parameter number has been included in the message, has not been followed by any argument. Neither the function nor any subsequent entries have been processed.

X 17 FUNCTION HAS TOO MANY ARGUMENTS — WILL IGNORE EXCESS. PARAMETER *xx*

An attempt has been made to write a function, whose parameter number has been included in the message, with more than the maximum number of arguments which can be handled by the statement. Only the arguments which can be accommodated have been passed on to the function subroutine.

W 18 DECIMAL DIGITS TRUNCATED. COMPUTATION MASK IS *IDD*

In order to accommodate sufficient integer digits in computing intermediate results, some of the decimal places of the input fields must be truncated. The message includes the four-digit computation mask; the first two digits indicate the maximum number of integers, the last two digits the maximum number of decimal digits. If an input field has any decimal digits in excess of this maximum, the digits will be truncated without rounding.

N 19 AUTO-DECIMAL FIELD HAS MORE THAN 20 INTEGERS

Since the maximum mode size for automatic-decimal computation is twenty digits, input fields having more than twenty integers would lead to meaningless results. A NOP has been generated.

W 20 INTEGER DIGITS MAY EXCEED COMPUTATION MASK *IDD*

An automatic-decimal computation may develop more integer digits than are provided for in the maximum format for intermediate results. This format has been included in the message in four-digit form, the first two digits representing the maximum number of integer digits that can be accommodated. (For details, remedies, etc., see "Mode Size for Automatic-Decimal Computations.")

W 21 SHIFT LEFT MAY LOSE HIGH-ORDER DIGIT(S)

Shifting left in the accumulator(s) to accommodate automatic-decimal num-

bers with more decimal places for addition or subtraction may have caused digit loss on the left. This message has been issued because no overflow branch has been specified.

**W 22 ACCUMULATOR OVERFLOW POSSIBLE**

An operation of addition or subtraction has been performed in the automatic-decimal mode that may have led to accumulator overflow. This message has been issued because no overflow branch has been specified.

**W 23 OVERFLOW POSSIBLE IN RESULT FIELD**

The number of integer digits that may have been developed in evaluating the arithmetic `EXPRESSION` exceeds the number available in the `RESULT` field. This message has been issued because no overflow branch has been specified. (Excess decimal digits will have been rounded off without warning.)

**N 24 FIELD SIZE EQUAL TO 0**

The parameter record has been incorrectly constructed by a higher level macro-instruction. A NOP has been generated.

**X 25 MISSING OPERATOR AFTER FLOATING POINT LITERAL**

Operand following that literal has been ignored.

**X 26 UNUSUAL PARAMETER**

The parameter was not usable to `ARITH` in its original form. The generator has produced instructions to treat the parameter as a 10-digit integer.

**N 27 EQUAL SIGN IN MIDDLE OF EXPRESSION**

Only one equal sign is allowed in a statement. A NOP has been generated.

**W 28 INVALID REFERENCE TO A CODE SUBSEQUENT ENTRY**

The `RESULT` field is a symbolic label of a `CODE` subsequent entry. Coding has been generated to store the result in the indicated literal.

**Examples**

The following are examples of acceptable coding for the `ARITH` macro-instruction. For each, the associated source-program entries are given, followed by the `ARITH` statement, coding generated in-line and (where applicable) coding generated out-of-line.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01 0054	*		ARITH EXAMPLE 1						
02 0055		DA	1				+0003250327		
03 0056	INPUT		12,19A5.3			29 0326		0326	
04 0057	DELTA		22,29A4.4			29 0327		0327	
05 0058	*								
06 0059	ANYLABEL	ARITH	INPUT=DELTA						
07	X ANYLABEL	ZA2	DELTA(0,7)	00001		0328	+2300290327		
08	X	SRR2	1			0329	+5000002101		
09	X	ST2	INPUT(0,7)			0330	+2200290326		
10 0060	*								

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01 0064	*		ARITH EXAMPLE 2						
02 0065		DA	1				+0003250325		
03 0066	X		1,3A3.0			13 0325		0325	
04 0067	*								
05 0068	ANYLABEL	ARITH	X=X+1						
06	X ANYLABEL	ZA2	X(0,2)	00001		0326	+2300130325	M	
07	X	A2	+1			0327	+2400000329		
08	X	ST2	X(0,2)			0328	+2200130325		
09 0069	*								
10 0070	*		THE FOLLOWING IS GENERATED OUT OF LINE						
11 0071	*		LITERALS						
12	X		+1			00 0329 +1		0329	

## ERROR MESSAGE LIST

PG/LN MESSAGE

AA 05 ARITH W 20 0300

ARITH Example 2

Warning message W 20 has been produced because an automatic-decimal computation might develop more integer digits than can be accommodated by the mask 3.0, which has been included in the message in the form 0300.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	1469	*		ARITH EXAMPLE 3					
02	1471		DA	1				+0003250330	
03	1472	A		11,18A4.4		18	0326		0326
04	1473	B		23,29A4.3		39	0327		0327
05	1474	C		30,39A4.6		09	0328		0328
06	1475	D		40,48A3.6		08	0329		0329
07	1476	X		50,59A4.6		09	0330		0330
08	14761	OVERFLOWBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0331	-0100090000	
09	1477	*							
10	1478	ANYLABEL	ARITH	X=A-B+ABS(C-D),OVERFLOWBR					
11	X	ANYLABEL	ZS2	D(0,3)			0332	-2300030329	
12	X		SL2	6			0333	+5000002206	
13	X		A2	C(0,9)			0334	+2400090328	
14	X		MSP	9992			0335	-0300919992	
15	X		ST2	COMAREA.A(0,9)+1	00002		0336	+2200090348	
16	X		ZS2	B(0,3)			0337	-2300360327	
17	X		SL2	4			0338	+5000002204	
18	X		A2	A(0,7)			0339	+2400180326	
19	X		SL2	2			0340	+5000002202	
20	X		A2	COMAREA.A(0,9)+1	00003		0341	+2400090348	
21	X		BV2	M.1			0342	+2100090345	
22	X		BV3	M.1			0343	+3100090345	
23	X		B	**2			0344	+0100090346	
24	X	M.1	BLX	93,OVERFLOWBR			0345	+0200930331	
25	X		ST2	X(0,9)	00004		0346	+2200090330	
26	1479	*							
27	1480	*		THE FOLLOWING IS GENERATED OUT OF LINE					
28	1481	*							
29	X	COMAREA.A	DA					+0003470348	

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY22, CHANGE LEVEL 00001.				PAGE AA		
LN	CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF
01	0075	*		ARITH EXAMPLE 4.				
02	0076	OVERFLOWBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001	0325	-0100090000	
03	0077		DA	1			+0003260330	
04	0078	A		11,19A5.4		19	0327	0327
05	0079	B		23,29A4.3		39	0328	0328
06	0080	C		30,39A4.6		09	0329	0329
07	0081	X		42,49A6.2		29	0330	0330
08	0082	*						
09	0083	ANYLABEL	ARITH	X=A+B-C,OVERFLOWBR				
10	X	ANYLABEL	ZA1	+0	00002	0331	+1300000366	
11	X		ZA2	B(0,6)		0332	+2300390328	
12	X		BLX	94,LINK.A		0333	+0200940343	
13	X		SL	1		0334	-5000000201	
14	X		A2	A(0,8)		0335	+2400190327	
15	X		BLX	94,LINK.A	00003	0336	+0200940343	
16	X		SL	2		0337	-5000000202	
17	X		S2	C(0,9)		0338	-2400090329	
18	X		BLX	94,LINK.A		0339	+0200940343	
19	X		SRR	4		0340	-5000000104	
20	X		ST2	X(0,7)	00004	0341	+2200290330	
21	0084	*						
22	0085	*		THE FOLLOWING IS GENERATED OUT OF LINE				
23	0086	*						
24	X		B	LINK8.A+1		0342	+0100090365	
25	X	LINK.A	BV2	LINK1.A IS THERE A CARRY		0343	+2100090351	
26	X	LINK3.A	BZ2	LINK4.A		0344	+2000090348	
			BZ1	LINK5.A		0345	+1000090349	
						0346	-1000090356	
45	X		A2	+1 CHANGE 9992		0362	+2400090365	
46	X	LINK8.A	B	0+X94		0363	+2400110366	
				LITERALS		0364	+0194090000	
47	X			+9999999999		09	0365	+9999999999
48	X			+0	00009	00	0366	+0
01	X			+1		11	0366	+1
								0366

## ARITH Example 4

Since  $MDL+MDR$  is greater than 10 in the fields defined in this example,  
the computation will be in twenty-digit mode, with a mask of 14.6.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	1496	*		ARITH EXAMPLE 5					
02	1497		DA	1				+0003250328	
03	1498	FA		00,09F		09	0325		0325
04	1499	FB		10,19F		09	0326		0326
05	1500	FC		20,29F		09	0327		0327
06	1501	FRESULT		30,39F		09	0328		0328
07	1502	*							
08	1503	ANYLABEL	ARITH	FRESULT=SQRT(FA+FB)+FC**2					M
09	X	ANYLABEL	ZA1	FA(0,9)	00001		0329	+1300090325	
10	X		A2	FB(0,3)			0330	+2400030326	
11	X		ST1	COMAREA.A(0,9)+1			0331	+1200090397	
12	X		ZA1	COMAREA.A(0,9)+1			0332	+1300090397	
13	X		BLX	94,SQRT.A			0333	+0200940354	
14	X		ST1	COMAREA.A(0,9)+1	00002		0334	+1200090397	
15	X		ZA2	FC(0,3)			0335	+2300030327	
16	X		ZA3	9992			0336	+3300099992	
17	X		M	9993			0337	+5300099993	
18	X		SLC	MACREG.1			0338	-5000010300	
19	X		ZA3	MACREG.1(4,5)	00003		0339	+3300450001	
20	X		C3	+0000000016			0340	+3500090399	
21	X		BL	**4			0341	+4000090345	
22	X		XS	MACREG.1,0			0342	-4700010000	
23	X		SRR	0+MACREG.1			0343	-5001000100	
24	X		B	**5	00004		0344	+0100090349	
25	X		ZA1	+9999999999			0345	+1300090401	
26	X		ZA2	+9999999999			0346	+2300090401	
27	X		A1	+5000000000			0347	+1400090400	
28	X		A1	+5000000000			0348	+1400090400	
29	X		ST2	COMAREA.A(0,9)+2	00005		0349	+2200090398	
30	X		ZA1	COMAREA.A(0,9)+1			0350	+1300090397	
31	X		A2	COMAREA.A(0,9)+2			0351	+2400090398	
32	X		ST1	FRESULT(0,9)			0352	+1200090328	
33	1504	*							
34	15041	*		THE FOLLOWING IS GENERATED OUT OF LINE					
35	15042	*							
36	X		B	M*10+1			0353	+0100090396	
37	X	SQRT.A	BZ1	0+X94	00006		0354	+1094090000	
38	X		BM1	M*4			0355	-1000090383	

25	X	M.7		+5000000000		09	0387		
26	X	M.8		+6250000001		09	0388	+6250000001	0388
27	X	M.9		+0101501025		09	0389	+0101501025	0389
28	X			+0201930800		09	0390	+0201930800	0390
29	X			+0702750550	00013	09	0391	+0702750550	0391
30	X			+1704350350		09	0392	+1704350350	0392
31	X			+2906180250		09	0393	+2906180250	0393
32	X			+5907600200		09	0394	+5907600200	0394
33	X	M.10		+9911230138		09	0395	+9911230138	0395
34	X	COMAREA.A	DA					+0003960398	
			LITERALS						
35	X			+0000000016	00014	09	0399	+0000000016	0399
36	X			+5000000000		09	0400	+5000000000	0400
37	X			+9999999999		09	0401	+9999999999	0401

ERROR MESSAGE LIST

PG/LN MESSAGE

AA 08 ARITH W 20 0400



PAGE	AA	PROGRAM	LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
			01	0088	*		ARITH EXAMPLE 6.						
			02	0089		DA	1				+0003250328		
			03	0090	PRESSURE		00,07A4.4			07	0325	0325	
			04	0091	VOLUME		10,15A2.4			05	0326	0326	
			05	0092	CENTIGRADE		20,25A2.4			05	0327	0327	
			06	0093	CONSTANT		30,39A6.4			09	0328	0328	
			07	0094	*								
			08	0095	ANYLABEL	ARITH	CONSTANT=PRESSURE*VOLUME/CENTIGRADE						
			09		X ANYLABEL	ZA3	PRESSURE(0,7)	00001		0329	+3300070325		
			10		X	M	VOLUME(0,5)			0330	+5300050326		
			11		X	SR	4			0331	-5000000004		
			12		X	ZA3	CENTIGRADE(0,5)			0332	+3300050327		
			13		X	BLX	93, DIV1.A			0333	+0200930338		
			14		X	XA	MACREG.1,0+5	00002		0334	+4700010005		
			15		X	B	DIV2.A			0335	+0100090346		
			16		X	ST2	CONSTANT(0,9)			0336	+2200090328		
			17	0096	*								
			18	0097	*		THE FOLLOWING IS GENERATED OUT OF LINE						
			19	0098	*								
			20		X	B	DIV3.A+6 TEMP BRANCH FOR GENOS			0337	+0100090370		
			21		X DIV1.A	BZ3	DIV3.A			0338	+3000090364		
			22		X	ZA1	+0	00003		0339	+1300000372		
			23		X	SLC	MACREG.1			0340	-5000010300		
			24		X	SR	1			0341	-5000000001		
			25		X	SLC3	MACREG.2			0342	+5000023300		
			26		X	XS	MACREG.1,10+MACREG.2			0343	-4702010010		
			27		X	D	9993	00004		0344	-5300099993		
			28		X	B	0+X93			0345	+0193090000		
			29		X DIV2.A	BXM	MACREG.1, DIV4.A			0346	-4400010355		
			30		X	ZA1	MACREG.1(4,5)			0347	+1300450001		
			31		X	C1	+11			0348	+1500120372		
			32		X	BL	*+3	00005		0349	+4000090352		
			33		X	SR2	10			0350	+5000002010		
			34		X	B	2+X93			0351	+0193090002		
			35		X	STD1	*(8,9)+1			0352	-1200890353		
			36		X	SRR2	10			0353	+5000002110		
			37		X	B	2+X93	00006		0354	+0193090002		
			38		X DIV4.A	SLC2	MACREG.2			0355	+5000022300		
			39		X	MSP	MACREG.1			0356	-0300910001		
						XS	MACREG.2,0+MACREG.1			0357	-4701020000		
			01		X	SL	0			0366	-5000000200		
			02		X	ZA2	+9999999999			0367	+2300090371		
			03		X	SR	0			0368	-5000000000		
			04		X	B	2+X93	00009		0369	+0193090002		
						LITERALS							
			05		X		+5000000000			09	0370	+5000000000	0370
			06		X		+9999999999			09	0371	+9999999999	0371
			07		X		+0			00	0372	+0	0372
			08		X		+11			12	0372	+ 11	0.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	1520	*		ARITH EXAMPLE 7					
02	1521		DA	1					
03	1522	A		00,07A4.4				+0003250327	
04	1523	B		08,15A4.4			07	0325	0325
05	1524	X		20,29F			89	0325	0325
06	1525	*					09	0327	0327
07	1526	ANYLABEL	ARITH	X=A-B+15/(A-B)					
08	X	ANYLABEL	ZS2	B(0,1)	00001		0328	-2300890325	M
09	X		SL2	2			0329	+5000002202	
10	X		S2	B(2,3)			0330	-2400010326	
11	X		SL2	4			0331	+5000002204	
12	X		A2	A(0,7)			0332	+2400070325	
13	X		ST2	COMAREA.A(0,9)+1	00002		0333	+2200090392	
14	X		ZA2	+15			0334	+2300340396	
15	X		ZA3	COMAREA.A(0,9)+1			0335	+3300090392	
16	X		BLX	93,DIV1.A			0336	+0200930391	
17	X		XA	MACREG.1,0+1			0337	+4700010001	
18	X		B	DIV2.A	00003		0338	+0100090359	
19	X		ST2	COMAREA.A(0,9)+1			0339	+2200090392	
20	X		ZS2	B(0,1)			0340	-2300890325	
21	X		SL2	2			0341	+5000002202	
22	X		S2	B(2,3)			0342	-2400010326	
23	X		SL2	4	00004		0343	+5000002204	
24	X		A2	A(0,7)			0344	+2400070325	
25	X		A2	COMAREA.A(0,9)+1			0345	+2400090392	
26	X		ZA1	+0			0346	+1300000396	
27	X		ZA3	+0000000070			0347	+3300090393	
28	X		BLX	94,FLOT2.A	00005		0348	+0200940385	
29	X		ST1	X(0,9)			0349	+1200090327	
30	1527	*							
31	15281	*		THE FOLLOWING IS GENERATED OUT OF LINE					
32	15282	*							
33	X		B	DIV3.A+6 TEMP BRANCH FOR GENOS			0350	+0100090383	
34	X	DIV3.A	BZ3	DIV3.A			0351	+3000090377	
35	X		ZA1	+0			0352	+1300000396	
	X		SLC	MACREG.1	00006		0353	-5000010300	
24	X		STJ3	9991(0,1)	00013		0387	-3400450001	
25	X	FLOT3.A	B	0+X94			0388	+5000001002	
26	X	COMAREA.A	DA				0389	-3200019991	
				LITERALS			0390	+0194090000	
								+0003910392	
27	X			+0000000070					
28	X			+5000000000	00014	09	0393	+0000000070	0393
29	X			+9999999999		09	0394	+5000000000	0394
30	X			+0		09	0395	+9999999999	0395
31	X			+11		00	0396	+0	0396
32	X			+15		12	0396	+ 11	0396
						34	0396	+ 15	0396

ERROR MESSAGE LIST

ARITH Example 7

PG/LN MESSAGE  
AA 07 ARITH W 20 0404

Warning message W 20 has been produced because the division operation might develop more integer digits than can be accommodated by the mask 4.4, which has been included in the message in the form 0404.

PAGE	AA	PROGRAM	LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	101	*					ARITH EXAMPLE 8					
02	102				DA		1				+0003250334	
03	103	A					03,19A9*8	39		0325		0325
04	104	B					37,49A9*4	79		0328		0328
05	105	C					53,69A9*8	39		0330		0330
06	106	D					72,79A5*3	29		0332		0332
07	107	RESULT					80,99A12*8	09		0333		0333
08	108	*										
09	109	ANYLABEL			ARITH		RESULT=(A+B-ABS(C))*D/A					M
10	X	ANYLABEL			ZA1		C(0,6)	00001		0335	+1300390330	
11	X				ZA2		C(7,16)			0336	+2300090331	
12	X				MSP		9991			0337	-0300919991	
13	X				MSP		9992			0338	-0300919992	
14	X				ST1		COMAREA*A(0,9)+2			0339	+1200090511	
15	X				ST2		COMAREA*A(0,9)+3	00002		0340	+2200090512	
16	X				ZA1		+0			0341	+1300000516	
17	X				ZA2		B(0,2)			0342	+2300790328	
18	X				SL		1			0343	-5000000201	
19	X				A2		B(3,3)			0344	+2400000329	
20	X				BLX		94*LINK*A	00003		0345	+0200940372	
21	X				SL		8			0346	-5000000208	
22	X				A1		A(0,6)			0347	+1400390325	
23	X				A2		A(7,16)			0348	+2400090326	
24	X				BLX		94*LINK*A			0349	+0200940372	
25	X				S1		COMAREA*A(0,9)+2	00004		0350	-1400090511	
26	X				S2		COMAREA*A(10,19)+2			0351	-2400090512	
27	X				BLX		94*LINK*A			0352	+0200940372	
28	X				ST1		COMAREA*A(0,9)			0353	+1200090509	
29	X				ST2		COMAREA*A(10,19)			0354	+2200090510	
30	X				ZA1		+0	00005		0355	+1300000516	
31	X				ZA2		D(0,7)			0356	+2300290332	
32	X				BLX		94*LINK*A			0357	+0200940372	
33	X				BLX		93*DUBLMPY*A			0358	+0200930395	
34	X				XS		MACREG*1,0+17			0359	-4700010017	
35	X				B		DUBLCK*A	00006		0360	+0100090431	
36	X				ST1		COMAREA*A(0,9)			0361	+1200090509	
37	X				ST2		COMAREA*A(10,19)			0362	+2200090510	
38	X				ZA1		A(0,6)			0363	+1300390325	
39	X				ZA2		A(7,16)			0364	+2300090326	
40	X				BLX		94*LINK*A	00007		0365	+0200940372	
41	X				BLX		93*DUBLDIV*A			0366	+0200930453	
42	X				XA		MACREG*1,0+11			0367	+4700010011	
43	X				B		DUBLCK*A			0368	+0100090431	
44	X				ST1		RESULT(0,9)			0369	+1200090333	
45	X				ST2		RESULT(10,19)	00006		0370	+2200090334	
46	110	*										
47	111	*					THE FOLLOWING IS GENERATED OUT OF LINE					
48	X				B		LINK8*A+1			0371	+0100090394	

PAGE AB		PROGRAM					PAGE AB	
LN	CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF
01	X	LINK1.A	BV2	LINK1.A IS THERE A CARRY		0372	+2100090380	
02	X	LINK3.A	BZ2	LINK4.A		0373	+2000090377	
03	X		BZ1	LINK5.A		0374	+1000090378	
04	X		BM1	LINK6.A	00009	0375	-1000090385	
05	X		BM2	LINK7.A		0376	-2000090390	
06	X	LINK4.A	SR	0 USE SIGN OF 9991		0377	-5000000000	
07	X	LINK5.A	SL	0 USE SIGN OF 9992		0378	-5000000200	
08	X		B	0+X94		0379	+0194090000	
09	X	LINK1.A	BM2	LINK2.A	00010	0380	-2000090383	
10	X		A1	+1 POSITIVE CARRY		0381	+1400110516	
11	X		B	LINK3.A		0382	+0100090373	
12	X	LINK2.A	S1	+1 NEGATIVE CARRY		0383	-1400110516	
13	X		B	LINK3.A		0384	+0100090373	
14	X	LINK6.A	BM2	0+X94 ARE BOTH ACC MINUS	00011	0385	-2094090000	
15	X		A1	+1 REVERSE CARRY (NEG)		0386	+1400110516	
16	X		S2	+9999999999 COMPLEMENT + SIGN)		0387	-2400090515	
17	X		S2	+1 CHANGE 9992 NG)		0388	-2400110516	
18	X		B	0+X94		0389	+0194090000	
19	X	LINK7.A	S1	+1 REVERSE CARRY (POS)	00012	0390	-1400110516	
20	X		A2	+9999999999 COMPLEMENT + SIGNS)		0391	+2400090515	
21	X		A2	+1 CHANGE 9992		0392	+2400110516	
22	X	LINK8.A	B	0+X94		0393	+0194090000	
23	X		B	HOLD.A+3 TEMP BRANCH		0394	+0100090430	
24	X	DUBLMPY.A	SLC	MACREG.1	00013	0395	-5000010300	
25	X		ZST1	MULT.A		0396	-1100090425	
26	X		ZST2	MULT.A+1		0397	-2100090426	
27	X		ZA1	COMAREA.A		0398	+1300090509	
28	X		ZA2	COMAREA.A+1		0399	+2300090510	
29	X		SLC	MACREG.2	00014	0400	-5000020300	
30	X		XA	MACREG.1,0+MACREG.2		0401	+4702010000	
31	X		ZST1	HOLD.A		0402	-1100090427	
	X		ZA3	9992		0403	+3300099992	
<hr/>								
35	X		A3	+5000000000				
36	X		B	DUBLOVFL.A+2				
37	X		DA	1			+0005060508	
38	X	DSOR.A		00,19	09	0506		0506
39	X	QUOT.A		20,29	09	0508		0508
40	X	COMAREA.A	DA	LITERALS			+0005090512	
41	X			+0000010009	00035	09	0513	0513
42	X			+5000000000		09	0514	0514
43	X			+9999999999		09	0515	0515
44	X			+0		00	0516	0516
45	X			+1		11	0516	0516
46	X			+20		23	0516	0516

ERROR MESSAGE LIST

PG/LN MESSAGE  
AA 09 ARITH W 20 1208

ARITH Example 8

Warning message W 20 has been produced because a multiplication or division operation might develop more integer digits than can be accommodated by the mask 12, 8, which has been included in the message in the form 1208.

## COMP—Compare and Branch

COMP generates instructions to compare two fields and to branch according to the results of the comparison.

### Source Program Format

The basic format for the COMP statement in the source program is as follows:

Line	Label	Operation	OPERAND				Basic Au				
3	56	15	16	20	21	25	30	35	40	45	50
0.1	ANY LABEL	COMP	FIELD1	FIELD2	LOWBR	EQUBR	HIGHBR				
0.2											

ANYLABEL may be any symbolic label; it may be omitted. The entry COMP must be written exactly as shown. FIELD1 and FIELD2 are either the symbolic names of the fields to be compared, or alphameric or numerical literals. Adcons are not permitted. LOWBR, EQUBR, and HIGHBR are the symbolic labels of instructions to which the program will branch if FIELD1 has the following relation to FIELD2, depending on the mode:

MODE	LOWBR	EQUBR	HIGHBR
Numerical	is less than	equals	is greater than
Alphameric	precedes	is identical to	follows

These results are determined by comparison techniques described under "Processing Techniques," below.

The basic format may be varied in two ways:

1. If the comparison is between numerical fields, either one or both may be replaced by the expression ABS(FIELDX), in which case the absolute value of the field will be used for the comparison. The parentheses must be written as shown. An attempt to take the absolute value of an alphameric field will result in an error message during assembly.
2. One or two of the branch addresses may be omitted. Instead of the missing branch, the object program would then take the next instruction. If another branch is specified after an omitted branch, separating commas must be punched; e.g., FIELD1, FIELD2, LOWBR, , HIGHBR.

### Processing Techniques

All comparisons are made on the basis of the standard IBM 7070 collating sequence. The fields to be compared may be numerical (either automatic-decimal or floating-decimal) or alphameric.

### Limitations on Length

The number of parameters is limited by the format. Automatic-decimal fields to be compared may bridge words but may not exceed twenty digits in length. Floating-decimal numbers must be contained within one location. Numerical

literals may not exceed twenty digits in length and must be signed. There is no limit on the length of alphameric fields, except that alphameric literals may not exceed 120 characters.

**Address Modification**

Modification by indexing and address adjustment is permitted on all symbolic addresses.

**The Effect of COMP**

Comparison of numerical fields is accomplished by subtraction, the correct branch address being determined by a negative, zero, or positive difference.

When an automatic-decimal number is to be compared to a floating-decimal number, the automatic-decimal number is first converted to floating-decimal format.

For the comparison of two automatic-decimal fields, the alignment of decimal points is automatic. If, after alignment of decimal points, the total number of integer and decimal digits in both fields is larger than 20, the excess number of digits will be truncated on the right without rounding; no warning message will be issued. A difference between the numbers in the truncated digits would not register in the comparison. Thus, if the following fields were to be compared using the COMP macro-instruction, the transfer would be to EQUBR:

<i>Fields</i>	<i>Format Specifications</i>	<i>Object Program Contents</i>
FIELD1	A12.8	000000000001.23456789
FIELD2	A2.18	01.234567891000000000

⚠ an automatic-decimal field of more than twenty digits in length is compared to a floating-decimal field, an automatic-decimal field, or an alphameric field, a warning message will be issued. Instructions will be generated to compare the absolute values of the fields.

⚠ a numerical field (either automatic-decimal or floating-decimal), which is of proper length, is compared to an alphameric field, a warning message will be issued. Instructions will be generated to compare the absolute values of the fields.

Fields containing numbers that are in double-digit representation must be converted to single-digit form before the COMP statement is employed (e.g., by use of the EDMOV macro-instruction). But numerical fields of different modes may be compared to each other, as may fields of the various alphameric types.

It should be noted that the COMP macro-instruction treats all fields according to their defined characteristics (or the absence of them), and not according to their object-program contents. Thus, difficulties may arise if numerical data is stored in fields defined as alphameric, or vice versa.

**Error and Warning Messages**

The following error and warning messages will be issued during assembly under the conditions specified:

**ALL BRANCHES BLANK**

All branches have been omitted from a comparison in which one or both of the fields are absolute. The comparison is made; three branches will be generated to the next in-line instructions (\*, \*+1, and \*+2) for patching purposes.

**ALL BRANCHES EQUAL**

All three branches are identical; coding will be generated.

FIELD 1 BLANK

Field 1 has been omitted. A NOP will be generated.

FIELD 2 BLANK

Field 2 has been omitted. A NOP will be generated.

FIELD 1 NOT ACCEPTABLE

Field 1 is an adcon. A NOP will be generated.

FIELD 2 NOT ACCEPTABLE

Field 2 is an adcon. A NOP will be generated.

LESS THAN 3 INPUT PARAMETERS

The minimum input in the COMP statement is FIELD1, FIELD2, and a branch address. The above message is issued if this minimum requirement is not met. A NOP will be generated.

NUMERIC FIELD GREATER THAN 20 DIGITS

A numerical field is greater than 20 digits in length. Instructions will be generated to compare the absolute values of the fields.

W-BOTH FIELDS NOT ALPHA-NOFORM

Either one field is alphameric and the other is not or one field is unspecified ("noform") and the other is not. Instructions will be generated to compare the absolute values of the fields.

W-UNUSUAL BRANCH CONDITION

A branch address in the operand of the COMP statement has *not* been left blank or is an address *other* than the label of an imperative statement (symbolic machine instruction or macro-instruction), or an actual storage address.

## Examples

The following are examples of acceptable coding for the COMP macro-instruction. For each, the associated source-program entries are given, followed by the COMP statement, coding generated in-line, and (where applicable) coding generated out-of-line.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	115	*		COMP EXAMPLE 1.					
02	116		DA	1				+0003250328	
03	117	FIELD1		02,19A8.10		29	0325		0325
04	118	FIELD2		23,39A8.9		39	0327		0327
05	119	LOWBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0329	-0100090000	
06	120	EQUALBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0330	-0100090000	
07	121	HIGHBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0331	-0100090000	
08	122	*							
09	123	ANYLABEL	COMP	FIELD1,FIELD2,LOWBR,EQUALBR,HIGHBR					
10		X ANYLABEL	ZS1	FIELD2(0,6)			0332	-1300390327	
11		X	ZS2	FIELD2(7,16)			0333	-2300090328	
12		X	BLX	94,LINK.A	00002		0334	+0200940345	
13		X	SL	1			0335	-5000000201	
14		X	A1	FIELD1(0,7)			0336	+1400290325	
15		X	A2	FIELD1(8,17)			0337	+2400090326	
16		X	BLX	94,LINK.A			0338	+0200940345	
17		X	BZ2	*+2	00003		0339	+2000090341	
18		X	B	*+2			0340	+0100090342	
19		X	BZ1	EQUALBR			0341	+1000090330	
20		X	BM1	LOWBR			0342	-1000090329	
21		X	B	HIGHBR			0343	+0100090331	
22	124	*							
23	125	*		THE FOLLOWING IS GENERATED OUT OF LINE					
24	126	*							
25		X	B	LINK8.A+1	00004		0344	+0100090367	
26		X LINK.A	BV2	LINK1.A IS THERE A CARRY			0345	+2100090353	
27		X LINK3.A	bZ2	LINK4.A			0346	+2000090350	
<hr/>									
47		X LINK8.A	A2	+1 CHANGE 9992			0365	+2400000368	
			B	0+X94			0366	+0194090000	
			LITERALS						
48		X		+9999999999		09	0367	+9999999999	0367
<hr/>									
			OPERAND						
01		X		+1		00	0368	+1	0368

COMP Example 1

The program will branch to LOWBR if FIELD1 is less than FIELD2, to EQUALBR if they are equal, and to HIGHBR if FIELD1 is greater than FIELD2.



PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	130	*		COMP EXAMPLE 2					
02	131		DA	1				+0003250328	
03	132	FIELD1		02,19A8,10		29	0325		0325
04	133	FIELD2		23,39A8,9		39	0327		0327
05	134	EQUALBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0329	-0100090000	
06	135	*							
07	136	ANYLABEL	COMP	FIELD1,FIELD2,,EQUALBR					
08	X	ANYLABEL	ZS1	FIELD2(0,6)			0330	-1300390327	
09	X		ZS2	FIELD2(7,16)			0331	-2300090328	
10	X		BLX	94,LINK.A			0332	+0200940341	
11	X		SL	1			0333	-5000000201	
12	X		A1	FIELD1(0,7)	00002		0334	+1400290325	
13	X		A2	FIELD1(8,17)			0335	+2400090326	
14	X		BLX	94,LINK.A			0336	+0200940341	
15	X		BZ2	**2			0337	+2000090339	
16	X		B	**2			0338	+0100090340	
17	X		BZ1	EQUALBR	00003		0339	+1000090329	
18	137	*							
19	138	*		THE FOLLOWING IS GENERATED OUT OF LINE					
20	139	*							
21	X		B	LINK8.A+1			0340	+0100090363	
22	X	LINK.A	BV2	LINK1.A IS THERE A CARRY			0341	+2100090349	
23	X	LINK3.A	BZ2	LINK4.A			0342	+2000090346	
24	X		BZ1	LINK5.A			0343	+1000090347	
25	X		BM1	LINK6.A	00004		0344	-1000090354	
26	X		BM2	LINK7.A			0345	-2000090359	
27	X	LINK4.A	SR	0 USE SIGN OF 9991			0346	-5000000000	
28	X	LINK5.A	SL	0 USE SIGN OF 9992			0347	-5000000200	
29	X		B	0+X94			0348	+0194090000	
30	X	LINK1.A	BM2	LINK2.A	00005		0349	-2000090352	
31	X		A1	+1 POSITIVE CARRY			0350	+1400000364	
32	X		B	LINK3.A			0351	+0100090342	
33	X	LINK2.A	S1	+1 NEGATIVE CARRY			0352	-1400000364	
34	X		B	LINK3.A			0353	+0100090342	
35	X	LINK6.A	BM2	0+X94 ARE BOTH ACC MINUS	00006		0354	-2094090000	
36	X		A1	+1 REVERSE CARRY (NEG)			0355	+1400000364	
37	X		S2	+9999999999 COMPLEMENT + SIGNG)			0356	-2400090363	
38	X		S2	+1 CHANGE 9992 NG)			0357	-2400000364	
39	X		B	0+X94			0358	+0194090000	
40	X	LINK7.A	S1	+1 REVERSE CARRY (POS)	00007		0359	-1400000364	
41	X		A2	+9999999999 COMPLEMENT + SIGNS)			0360	+2400090363	
42	X		A2	+1 CHANGE 9992			0361	+2400000364	
43	X	LINK8.A	B	0+X94			0362	+0194090000	
				LITERALS					
44	X			+9999999999		09	0363	+9999999999	0363
45	X			+1	00008	00	0364	+1	0364

## COMP Example 2

This example is the same as Example 1 except that both LOWBR and HIGHBR have been omitted. The program will continue sequentially if FIELD1 is either greater than or less than FIELD2.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	142	*		COMP EXAMPLE 3.					
02	143		DA	1					
03	144	FIELD1		02,19A8,10				+0003250326	
04	145		DA	2,RDW,0+INDEXWORD			29	0325	0325
05		X						+0003270336	
06		X					00001	0327 +0003290332	0327
07	146	FIELD2		23,39A8,9				0328 -0003330336	0328
08	147	LOWBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH			39	0331	0002
09	148	HIGHBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			00002	0337 -0100090000	
10	149	*						0338 -0100090000	
11	150	ANYLABEL	COMP	ABS(FIELD1),FIELD2,LOWBR, HIGHBR					
12		X ANYLABEL	ZA1	FIELD1(0,7)					
13		X	ZA2	FIELD1(8,17)				0339 +1300290325	
14		X	MSP	9991				0340 +2300090326	
15		X	MSP	9992				0341 -0300919991	
16		X	ST1	COMAREA.A(0,9)+2			00003	0342 -0300919992	
17		X	ST2	COMAREA.A(0,9)+3				0343 +1200090382	
18		X	ZS1	FIELD2(0,6)+INDEXWORD				0344 +2200090383	
19		X	ZS2	FIELD2(7,16)+INDEXWORD				0345 -1301390002	
20		X	BLX	94,LINK.A				0346 -2301090003	
21		X	SL	1			00004	0347 +0200940358	
22		X	A1	COMAREA.A(0,9)+2				0348 -5000000201	
23		X	A2	COMAREA.A(10,19)+2				0349 +1400090382	
24		X	BLX	94,LINK.A				0350 +2400090383	
25		X	BZ2	**2				0351 +0200940358	
26		X	B	**2			00005	0352 +2000090354	
27		X	BZ1	**3				0353 +0100090355	
28		X	BM1	LOWBR				0354 +1000090357	
29		X	B	HIGHBR				0355 -1000090337	
30	151	*						0356 +0100090338	
31	152	*		THE FOLLOWING IS GENERATED OUT OF LINE					
32	153	*							
33		X	B	LINK8.A+1					
34		X LINK.A	BV2	LINK1.A IS THERE A CARRY			00006	0357 +0100090380	
35		X LINK3.A	BZ2	LINK4.A				0358 +2100090366	
36		X						0359 +2000090363	
								0360 +1000090364	
								0361 -1000090371	
03		X	B	0+X94				0374 -2400000385	
04		X LINK7.A	S1	+1 REVERSE CARRY (POS)				0375 +0194090000	
05		X	A2	+9999999999 COMPLEMENT + SIGNS)				0376 -1400000385	
06		X	A2	+1 CHANGE 9992			00010	0377 +2400090384	
07		X LINK8.A	B	0+X94				0378 +2400000385	
08		X COMAREA.A	DA					0379 +0194090000	
				LITERALS				+0003800383	
09		X		+9999999999					
10		X		+1			00011 09	0384 +9999999999	0384
							00	0385 +1	0385

COMP Example 3

The program will branch to LOWBR if the absolute value of FIELD1 is less than FIELD2 and to HIGHBR if it is greater. If the absolute value of FIELD1 equals FIELD2, the program will continue sequentially.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01 157	*		COMP EXAMPLE 4.					
02 158		DA	2,*0+INDEXWORD				+0003250328	
03 159	FIELD1		00,19'			09 0325		0000
04 160		DA	1				+0003290332	
05 161	FIELD2		20,39'			09 0331		0331
06 162	LOWBR	NOP		00001			0333 -0100090000	
07 163	EQUALBR	NOP					0334 -0100090000	
08 164	HIGHBR	NOP					0335 -0100090000	
09 165	*							
10 166	ANYLABEL	COMP	FIELD1,FIELD2,LOWBR,EQUALBR,HIGHBR					
11 X	ANYLABEL	ZAA	FIELD2(0,9)				0336 +1600090331	
12 X		CA	FIELD1(0,9)+INDEXWORD				0337 -1501090000	
13 X		BL	HIGHBR	00002			0338 +4000090335	
14 X		BH	LOWBR				0339 -4000090333	
15 X		ZAA	FIELD2(10,19)				0340 +1600090332	
16 X		CA	FIELD1(10,19)+INDEXWORD				0341 -1501090001	
17 X		BL	HIGHBR				0342 +4000090335	
18 X		BE	EQUALBR	00003			0343 -4100090334	
19 X		B	LOWBR				0344 +0100090333	
20 167	*							

## COMP Example 4

If FIELD1 precedes, is identical to, or follows FIELD2 in the standard collating sequence, the program will branch to LOWBR, EQUALBR or HIGHBR, respectively.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01 171	*		COMP EXAMPLE 5.					
02 172		DA	2,*0+X15				+0003250328	
03 173	FIELD1		00,19'			09 0325		0000
04 174		DA	2,*0+INDEXWORD				+0003290338	
05 175	FIELD2		20,45'			09 0331		0002
06 176	DIFFBR	NOP		00001			0339 -0100090000	
07 177	EQUBR	NOP					0340 -0100090000	
08 178	*							
09 179	ANYLABEL	COMP	FIELD1,FIELD2,DIFFBR,EQUBR,DIFFBR					
10 X	ANYLABEL	ZAA	FIELD1(0,9)+X15				0341 +1615090000	
11 X		CA	FIELD2(0,9)+INDEXWORD				0342 -1501090002	
12 X		BL	DIFFBR				0343 +4000090339	
13 X		BH	DIFFBR	00002			0344 -4000090339	
14 X		ZAA	FIELD1(10,19)+X15				0345 +1615090001	
15 X		CA	FIELD2(10,19)+INDEXWORD				0346 -1501090003	
16 X		BL	DIFFBR				0347 +4000090339	
17 X		BH	DIFFBR				0348 -4000090339	
18 X		ZA1	FIELD2(0,9)+2+INDEXWORD	00003			0349 +1301090004	
19 X		BZ1	EQUBR				0350 +1000090340	
20 X		B	DIFFBR				0351 +0100090339	
21 180	*							

## COMP Example 5

Fields which are identical will cause a branch to EQUBR; fields which differ to DIFFBR.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01 184	*		COMP EXAMPLE 6.						
02 185		DA	1				+0003250326		
03 186	FIELD1		00,09F		09	0325		0325	
04 187	FIELD2		10,19F		09	0326		0326	
05 188	EQUALBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0327	-0100090000		
06 189	HIGHBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0328	-0100090000		
07 190	*								
08 191	ANYLABEL	COMP	FIELD1,FIELD2,,EQUALBR,HIGHBR						
09	X ANYLABEL	ZA1	FIELD1(0,9)			0329	+1300090325		
10	X	S1	FIELD2(0,9)			0330	-1400090326		
11	X	BV1	**2			0331	+1100090333		
12	X	BZ1	EQUALBR	00002		0332	+1000090327		
13	X	BM1	**2			0333	-1000090335		
14	X	B	HIGHBR			0334	+0100090328		
15 192	*								

COMP Example 6

The program will go to the next sequential instruction if FIELD1 is less than FIELD2, to EQUALBR if FIELD1 is equal to FIELD2, and to HIGHBR if FIELD1 is greater than FIELD2.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01 201	*		COMP EXAMPLE 7.						
02 202		DA	1				+0003250332		
03 203	FIELD1		00,39'		09	0325		0325	
04 204	FIELD2		40,79'		09	0329		0329	
05 205	DIFFBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0333	-0100090000		
06 206	*								
07 207	ANYLABEL	COMP	FIELD1,FIELD2,DIFFBR,,DIFFBR						
08	X ANYLABEL	XL	MACREG.01,+0000000003			0334	+4500010340		
09	X M.4	ZAA	FIELD2(0,9)+MACREG.01			0335	+1601090329		
10	X	CA	FIELD1(0,9)+MACREG.01			0336	-1501090325		
11	X	BL	DIFFBR			0337	+4000090333		
12	X	BH	DIFFBR	00002		0338	-4000090333		
13	X	BIX	MACREG.01,M.4			0339	+4900010335		
14 208	*								
15	X	LITERALS	+0000000003		09	0340	+0000000003	0340	

COMP Example 7

If FIELD1 = FIELD2, the program will continue with the next sequential instruction. If not, it will branch to DIFFBR.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.								PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	214	*		COMP EXAMPLE 8.						
02	215		DA	1				+0003250327		
03	216	FIELD1		00,09F		09	0325		0325	
04	217	FIELD2		10,29A13.7		09	0326		0326	
05	218	LOWBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0328	-0100090000		
06	219	EQUALBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0329	-0100090000		
07	220	*								
08	221	ANYLABEL	COMP	FIELD1,FIELD2,LOWBR,EQUALBR						
09	X	ANYLABEL	ZA1	FIELD2(0,9)			0330	+1300090326		
10	X		ZA2	FIELD2(10,19)			0331	+2300090327		
11	X		ZA3	+0000000063			0332	+3300090347		
12	X		BLX	94,FLOT2.A	00002		0333	+0200940341		
13	X		S1	FIELD1(0,9)			0334	-1400090325		
14	X		BV1	**2			0335	+1100090337		
15	X		BZ1	EQUALBR			0336	+1000090329		
16	X		BM1	**2			0337	-1000090339		
17	X		B	LOWBR	00003		0338	+0100090328		
18	222	*								
19	223	*		THE FOLLOWING IS GENERATED OUT OF LINE						
20	224	*								
21	X	FLOT1.A	SLC1	MACREG.1			0339	+5000011300		
22	X		B	**2			0340	+0100090342		
23	X	FLOT2.A	SLC	MACREG.1			0341	-5000010300		
24	X		BZ1	0+X94			0342	+1094090000		
25	X		S3	MACREG.1(4,5)	00004		0343	-3400450001		
26	X		SR1	2			0344	+5000001002		
27	X		STD3	9991(0,1)			0345	-3200019991		
28	X	FLOT3.A	B	0+X94			0346	+0194090000		
			LITERALS							
29	X			+0000000063		09	0347	+0000000063	0347	

### COMP Example 8

The program will continue with the next instruction if FIELD1 is greater than FIELD2. It will branch to EQUALBR if FIELD1 equals FIELD2 and to LOWBR if FIELD1 is less than FIELD2.

PAGE AA	PROGRAM	LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01		227	*			COMP EXAMPLE 9						
02		228		AREANAME	DA	3,0+INDEXWORD			29	0325	+0003250336	0000
03		229		FIELD1		02,19A8.10			39	0327		0002
04		230		FIELD2		23,39A8.9	00001			0337	-0100090000	
05		231		LOWBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE				0338	-0100090000	
06		232		EQUALBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE				0339	-0100090000	
07		233		HIGHBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE						
08		234	*									
09		235		ANYLABEL	COMP	FIELD1,FIELD2,LOWBR, EQUALBR,HIGHBR						
10		236				REMARKS MAY BE USED.						
11			X	ANYLABEL	ZS1	FIELD2(0,6)+INDEXWORD				0340	-1301390002	
12			X		ZS2	FIELD2(7,16)+INDEXWORD				0341	-2301090003	
13			X		BLX	94,LINK.A	00002			0342	+0200940353	
14			X		SL	1				0343	-500000201	
15			X		A1	FIELD1(0,7)+INDEXWORD				0344	+1401290000	
16			X		A2	FIELD1(8,17)+INDEXWORD				0345	+2401090001	
17			X		BLX	94,LINK.A				0346	+0200940353	
18			X		BZ2	**2	00003			0347	+2000090349	
19			X		B	**2				0348	+0100090350	
20			X		BZ1	EQUALBR				0349	+1000090338	
21			X		BM1	LOWBR				0350	-1000090337	
22			X		B	HIGHBR				0351	+0100090339	
23		237	*									
24		238	*			THE FOLLOWING IS GENERATED OUT OF LINE						
25		239	*									
26			X		B	LINK8.A+1	00004			0352	+0100090375	
27			X	LINK.A	BV2	LINK1.A IS THERE A CARRY				0353	+2100090361	
28			X	LINK3.A	BZ2	LINK4.A				0354	+2000090358	
29			X		BZ1	LINK5.A				0355	+1000090359	
30			X		BM1	LINK6.A				0356	-1000090366	
31			X		BM2	LINK7.A	00005			0357	-2000090371	
32			X	LINK4.A	SR	0 USE SIGN OF 9991				0358	-5000000000	
33			X	LINK5.A	SL	0 USE SIGN OF 9992				0359	-5000000200	
34			X		B	0+X94				0360	+0194090000	
35			X	LINK1.A	BM2	LINK2.A				0361	-2000090364	
45			X	LINK7.A	S1	+1 REVERSE CARRY (POS)				0370	+0194090000	
46			X		A2	+9999999999 COMPLEMENT + SIGNS)	00008			0371	-1400000376	
47			X		A2	+1 CHANGE 9992				0372	+2400090375	
48			X	LINK8.A	B	0+X94				0373	+2400000376	
										0374	+0194090000	
						LITERALS						
01			X			+9999999999			09	0375	+9999999999	0375
02			X			+1			00	0376	+1	0376

**CYCLE—Cycle Branch**  
**RECYC—Reset Cycle**

CYCLE generates instructions to branch a specified number of times to each of a series of locations. RECYC generates instructions to reinitialize one or more CYCLE macro-instructions.

**Source Program Format**

The basic formats for the CYCLE and RECYC statements in the source program are as follows:

Line	Label	Operation	OPERAND					Basic Autocoder		
3	5 6	15 16	20 21	25	30	35	40	45	50	55
0 1	ANY LABEL	CYCLE	BRANCH1, COUNTER1, BRANCH2, COUNTER2, etc.							
0 2	ANY LABEL	RECYC	CYCLE1, CYCLE2, etc.							
0 3										

In these examples, ANYLABEL is any symbolic label; it may be omitted. The entries CYCLE and RECYC must be written exactly as shown. In the CYCLE statement, the various BRANCHX entries are the symbolic labels of instructions to which the program will transfer, taking each BRANCH from left-to-right and going to each as many times as is specified by the associated COUNTERX. The branching will occur every time the program reaches the position originally occupied by the CYCLE statement. After each BRANCH is taken, it is assumed that the program will return to a location preceding the location of the CYCLE statement or to that location itself. The COUNTERS may be symbolic or they may be unsigned actual integers (representing literal count numbers) up to a maximum of ten digits. If the reference is symbolic, the count number must be stored in a single location; i.e., it may not bridge words. Symbolic and literal counters may be freely mixed within the operand of a single CYCLE macro-instruction.

After all BRANCHES have been executed the number of times indicated, the next pass through the CYCLE statement will resume the entire process from the beginning. However, if there is only one BRANCH and one COUNTER, the macro-instruction becomes inactive (in effect, a NOP) after the branch has been made the requisite number of times, and the program continues sequentially.

The CYCLEX entries in the operand portion of the RECYC statement are the labels of CYCLE macro-instructions. The RECYC macro-instruction will restore the settings of the various counters in these CYCLE statements to their original values. Thus, after a CYCLE statement has been named in a RECYC macro-instruction, the next program pass through that CYCLE instruction will result in a "first-time" transfer to the BRANCH1 address.

The format of the CYCLE instruction may be modified in two ways:

1. A BRANCHX may be omitted, with separating commas entered as in the following example:

Line	Label	Operation	OPERAND				
3	5 6	15 16 20 21	25	30	35	40	45
0 1	ANY LABEL	CYCLE	BR1	, CTR1	, ,	CTR2	
0 2							

If this is done, the program will go to the next instruction of the source program for the specified number of times, instead of to the missing BRANCHX address. Thereafter, the next indicated BRANCH will be taken as usual. In particular, it is possible to omit the BRANCH1, beginning the operand portion of the CYCLE statement with a comma. In that case, the CYCLE statement will, in effect, result in a NOP on the first n program passes, where n is the number specified by the COUNTER1.

2. The last COUNTER may be omitted as follows:

Line	Label	Operation	OPERAND				
3	5 6	15 16 20 21	25	30	35	40	45
0 1	ANY LABEL	CYCLE	BR1	, CTR1	, BR2	, CTR2	, BR3
0 2							

This will cause the program to branch permanently to the last BRANCH after the other BRANCHES have been executed the required number of times. These two devices may be combined in a single CYCLE statement; for an example, see "Limitations on Length," below.

If the counters are symbolic, the number of times the program is to branch to each location must have been entered into the proper field before the set of instructions generated by CYCLE is entered for the first time. When the program reaches the CYCLE instruction, all COUNTERS in the operand are locked; they cannot be altered until all BRANCHES have been executed as required, except through a RECYC macro-instruction. Since the BRANCHES are taken from left-to-right, it is possible to alter the COUNTER for a *prior* BRANCH while a given BRANCH is being taken. This change will take effect when the entire CYCLE is restarted.

## Processing Techniques

### Limitations on Length

A maximum of five BRANCH and COUNTER pairs, including omitted parameters, may be entered. If more than five pairs are required, two successive CYCLE macro-instructions may be written as follows:

Line	Label	Operation	OPERAND					Basic Autocoder		
3	5 6	15 16 20 21	25	30	35	40	45	50	55	
0 1	ANY LABEL 1	CYCLE	BR1	, CTR1	, BR2	, CTR2	, BR3	, CTR3	, BR4	, CTR4
0 2	ANY LABEL 2	CYCLE	BR5	, CTR5	, BR6	, CTR6	, etc.			
0 3										



The first `CYCLE` statement will become inactive after the fourth `BRANCH` has been taken the required number of times; the program will then permanently “fall through” to the second statement.

`RECYC` allows for 94 entries in the operand portion.

### **Address Modification**

All symbolic addresses may be modified by indexing and address adjustment.

### **Error and Warning Messages**

The `CYCLE` macro generator will issue the following error and warning messages during assembly under the conditions specified:

#### **ASSUME COMMAS AFTER PAR. 2 — PERMANENT NOP**

This warning message is issued if the operand contains only one `BRANCH` and one `COUNTER`. After the `BRANCH` has been taken the specified number of times, the macro-instruction will not be reinitialized (which would make it the equivalent of an unconditional Branch) but will become permanently inactive (`NOP`). This is what would happen if the `COUNTER` were followed by two commas.

#### **BRIDGE CTR USING FIRST LOC ONLY PAR xx**

A symbol `COUNTER` has been used that bridges two storage locations. The message will give the parameter number of the faulty entry in place of the `xx`. Only the digits from the portion in the first location will be used as the `COUNTER`.

#### **ERROR — IMPROPER OPERAND**

This message is issued if there are less than two or more than ten parameters in the operand, or if an omitted `COUNTER` is followed by another `BRANCH`, which would never become effective.

#### **WARNING — SUCCESSIVE NOPS**

This message will be issued if two successive `BRANCHES` have been omitted from the operand. Nevertheless, the appropriate coding will be generated as if the double omission were intentional.

The `RECYC` macro generator will issue the following error message under the condition specified:

#### **PARAM NOT LABEL OF CYCLE MACRO xx**

The `xx` will be replaced by the number of the parameter which is not the label of a `CYCLE` statement.

### **Examples**

The following are examples of acceptable coding for the `CYCLE` and `RECYC` macro-instructions. For each, the associated source-program entries are given, followed by the `CYCLE` or `RECYC` statement, coding generated in-line, and (where applicable) coding generated out-of-line.

PAGE AA	PROGRAM				7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.				PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF	
01	243	*		CYCLE EXAMPLE 1.					
02	244	BRANCHA	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001	0325	-0100090000		
03	245	BRANCHB	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE		0326	-0100090000		
04	246	*							
05	247	ANYLABEL	CYCLE	BRANCHA,1,0BRANCHB,1					
06		X ANYLABEL	NOP	M.1		0327	-0100090330		
07		X	MSP	ANYLABEL		0328	-0300910327		
08		X	B	BRANCHA		0329	+0100090325		
09		X M.1	M5M	ANYLABEL	00002	0330	-0300610327		
10		X	B	BRANCHB		0331	+0100090326		
11		X M.6	NOP	0		0332	-0100090000		
12		X ORIGIN	CNTRL	*-1					
13	248	*							

#### CYCLE Example 1

CYCLE is used in this example to create a strictly alternating condition.

The first time the program arrives at this instruction, it will branch to

BRANCHA, the next time to BRANCHB, then to BRANCHA, then to

BRANCHB again, and so on for the duration of the program.

PAGE AA	PROGRAM				7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.				PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF	
01	252	*		CYCLE EXAMPLE 2.					
02	253	POINTX	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001	0325	-0100090000		
03	254	POINTY	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE		0326	-0100090000		
04	255	*							
05	256	ANYFIELD	CYCLE	POINTX,1,0POINTY					
06		X ANYFIELD	NOP	M.1		0327	-0100090330		
07		X	MSP	ANYFIELD		0328	-0300910327		
08		X	B	POINTX		0329	+0100090325		
09		X M.1	B	POINTY	00002	0330	+0100090326		
10		X M.6	NOP	0		0331	-0100090000		
11		X ORIGIN	CNTRL	*-1					
12	257	*							

#### CYCLE Example 2

When the program arrives at CYCLE the first time, it will branch to

POINTX. Each subsequent time it will branch to POINTY.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION QMY08, CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	261	*		CYCLE EXAMPLE 3.					
02	262	TOTALLINE	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0325	-0100090000	
03	263	REPORTLINE	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0326	-0100090000	
04	264	*							
05	265	ANYLABEL	CYCLE	REPORTLINE,35,TOTALLINE,1					
06	X	ANYLABEL	NOP	M.1			0327	-0100090331	
07	X		ZSA	+35			0328	-1600010338	
08	X		ST1	M.9(2,5)			0329	+1200250337	
09	X		MSP	ANYLABEL	00002		0330	-0300910327	
10	X	M.1	XL	MACREG.1,M.9			0331	+4500010337	
11	X		XA	MACREG.1,1			0332	+4700010001	
12	X		XU	MACREG.1,M.9			0333	-4500010337	
13	X		BXM	MACREG.1,REPORTLINE			0334	-4400010326	
14	X		MSM	ANYLABEL	00003		0335	-0300610327	
15	X		B	TOTALLINE			0336	+0100090325	
16	266	*							
17	267	*		THE FOLLOWING IS GENERATED OUT OF LINE					
18	268	*							
19	X	M.9	DA	1				+0003370337	
20	X			00,09		09	0337		0337
				LITERALS					
21	X			+35	00004	01	0338	+35	0338

## CYCLE Example 3

In this example, the program will branch to the REPORTLINE routine

35 times, to the TOTALLINE routine once, then to the REPORTLINE

routine 35 times, to the TOTALLINE routine once, and so on.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION QMY08, CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	272	*		CYCLE EXAMPLE 4.					
02	273	BRANCHONE	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0325	-0100090000	
03	274	BRANCHTWO	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0326	-0100090000	
04	275	*							
05	276	ANYLABEL	CYCLE	BRANCHONE,1,1,1,BRANCHTWO					
06	X	ANYLABEL	NOP	M.1			0327	-0100090331	
07	X		MSP	ANYLABEL			0328	-0300910327	
08	X		MSM	M.1			0329	-0300610331	
09	X		B	BRANCHONE	00002		0330	+0100090325	
10	X	M.1	NOP	M.2			0331	-0100090334	
11	X		MSP	M.1			0332	-0300910331	
12	X		B	M.6			0333	+0100090335	
13	X	M.2	B	BRANCHTWO			0334	+0100090326	
14	X	M.6	NOP	0	00003		0335	-0100090000	
15	X	ORIGIN	CNTRL	*-1					
16	277	*							

## CYCLE Example 4

The first time the program reaches CYCLE, a branch is made to BRANCHONE. On the second pass the next instruction following those generated by CYCLE is executed, since a branch address was omitted. Each subsequent time, a branch is made to BRANCHTWO since the final counter was omitted.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION QMY08; CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	301	*		CYCLE EXAMPLE 5.					
02	302	BRANCH1	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0325	-0100090000	
03	303	BRANCH3	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0326	-0100090000	
04	304	COUNTER1	NOP	REPRESENTS COUNTER LABEL			0327	-0100090000	
05	305	COUNTER2	NOP	REPRESENTS COUNTER LABEL			0328	-0100090000	
06	306	COUNTER3	NOP	REPRESENTS COUNTER LABEL			0329	-0100090000	
07	307	*							
08	308	ANYLABEL	CYCLE	BRANCH1,COUNTER1,,COUNTER2					
09	X	ANYLABEL	NOP	M.1	00002		0330	-0100090338	
10	X		ZSA	COUNTER1			0331	-1600090327	
11	X		ST1	M.9+1			0332	+1200090350	
12	X		ZSA	COUNTER2			0333	-1600090328	
13	X		ST1	M.9+2			0334	+1200090351	
14	X		ZAA	M.8	00003		0335	+1600090365	
15	X		ST1	M.9			0336	+1200090349	
16	X		MSP	ANYLABEL			0337	-0300910330	
17	X	M.1	XL	MACREG.1,M.9			0338	+4500010349	
18	X	M.2	ZA1	M.9+MACREG.1			0339	+1301090349	
19	X		A1	+1	00004		0340	+1400000367	
20	X		ZST1	M.9+MACREG.1			0341	-1101090349	
21	X		BM1	M.5-1+MACREG.1			0342	-1001090351	
22	X		XL	MACREG.1,M.9			0343	+4500010349	
23	X		XA	MACREG.1,1			0344	+4700010001	
24	X		XU	MACREG.1,M.9	00005		0345	-4500010349	
25	X		BCX	MACREG.1,M.2			0346	-4300010339	
26	X		MSM	ANYLABEL			0347	-0300610330	
27	X		B	ANYLABEL			0348	+0100090330	
28	X	M.9	DA	1				+0003490351	
29	X			00,29	09		0349		0349
30	X	M.5	B	BRANCH1	00006		0352	+0100090325	
31	X		B	M.6			0353	+0100090354	
32	X	M.6	NOP	0			0354	-0100090000	
33	X	ORIGIN	CNTRL	*-1					
34	309		CYCLE	BRANCH3,COUNTER3,,1					***
35	X	M.19	NOP	M.10			0354	-0100090358	
36	X		ZSA	COUNTER3			0355	-1600090329	
37	X		ST1	M.18(2,5)			0356	+1200250366	
38	X		MSP	M.19	00007		0357	-0300910354	
39	X	M.10	XL	MACREG.1,M.18			0358	+4500010366	
40	X		XA	MACREG.1,1			0359	+4700010001	
41	X		XU	MACREG.1,M.18			0360	-4500010366	
42	X		BXM	MACREG.1,BRANCH3			0361	-4400010326	
43	X		MSM	M.19	00008		0362	-0300610354	
44	X	M.15	NOP	0			0363	-0100090000	
45	X	ORIGIN	CNTRL	*-1					
46	310	SOMELABEL	RECYC	ANYLABEL					***
47	X	SOMELABEL	MSM	ANYLABEL			0363	-0300610330	
48	311		B	ANYLABEL			0364	+0100090330	

PAGE AB	PROGRAM								PAGE AB
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	312	*							
02	313	*		THE FOLLOWING IS GENERATED OUT OF LINE					
03	314	*							
04		X M.8	DC					+0003650365	
05		X		+0000010002			09 0365	+0000010002	0365
06		X M.18	DA	1				+0003660366	
07		X		00,09			09 0366		0366
				LITERALS					
08		X		+1	00009	00	0367	+1	0367

## CYCLE Example 5

This example illustrates a technique that will allow the modification of the contents of a counter associated with a branch when one of the preceding branches has already been entered. In this case, COUNTER3 may be freely changed while BRANCH1 is being taken by the program. (This would not be the case if COUNTER3 were written into the operand of the same CYCLE statement as BRANCH1.) COUNTER2 must be set to a value greater than the maximum possible content of COUNTER3. This must be done before the first CYCLE statement is entered. The program will BRANCH to BRANCH1 as many times as specified in COUNTER1. It will then "fall through" to the second CYCLE statement, taking BRANCH3 as many times as COUNTER3 indicates, and finally "fall through" to the RECYC statement. This will reinitialize the first CYCLE macro-instruction (the second one will have reinitialized itself, having completed all branches the required number of times), to which a transfer is then made by means of the unconditional Branch instruction.

NOTE: Lines marked thus \*\*\* in the example listing are original source statements; the intervening instructions are generated.

PAGE AA	PROGRAM								PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	318	*		RECYC EXAMPLE 1.					
02	319	SOMELABEL	RECYC	ANYLABEL,LABELZ,ANYSYMBOL,POINTX					
03		X SOMELABEL	MSM	ANYLABEL	00001		0325	-0300610329	
04		X	MSM	LABELZ			0326	-0300610353	
05		X	MSM	ANYSYMBOL			0327	-0300610377	
06		X	MSM	POINTX			0328	-0300610401	
07	320	*							
08	321	ANYLABEL	CYCLE	BRANCH1,COUNTER1,COUNTER2					
09		X ANYLABEL	NOP	M.1			0329	-0100090337	
10		X	ZSA	COUNTER1	00002		0330	-1600090426	
11		X	ST1	M.9+1			0331	+1200090349	
12		X	ZSA	COUNTER2			0332	-1600090427	
13		X	ST1	M.9+2			0333	+1200090350	
14		X	ZAA	M.8			0334	+1600090428	

## RECYC Example 1

The CYCLE macro-instructions whose labels are listed in the operand are reinitialized by this instruction.

## DECOD—Branch on Code Value

DECOD generates instructions to analyze a code field and to branch according to the value it contains. The code field must have been previously established through a CODE declarative statement.

### Source Program Format

The basic format for the DECOD statement in the source program is as follows:

Line	Label	Operation	OPERAND											
5	6	15/16	20/21	25	30	35	40	45	Basic Autocoder			Autocod		
									50	55	60	65	70	
0 1	ANY LABEL	DECOD	CODENAME	, CODEVALUE 1	, BRANCH 1	, CODEVALUE 2	, BRANCH 2	, etc.						
0 2														

ANYLABEL is any symbolic label; it may be omitted. The entry DECOD must be written exactly as shown. CODENAME must be the symbolic label of a CODE header line that appears elsewhere in the program. The CODEVALUEX entries are the symbolic names of the conditions for which tests are to be made; they must be subsequent entries under the CODE header line, CODENAME. The BRANCHX entries are the symbolic addresses of instructions to which the program is to transfer if the associated CODEVALUE is present in the CODE field to be analyzed.

Omission of BRANCH addresses is not permitted in writing the DECOD statement. If it is desired to have the program continue sequentially in case a given CODEVALUE is present in the CODENAME field, the particular CODEVALUE should be omitted from the operand altogether.

### Processing Techniques

#### Limitations on Length

A maximum of 45 CODEVALUE and BRANCH pairs may be entered.

#### Address Modification

The CODENAME and BRANCH entries may be modified by indexing and address adjustment.

#### Processing Sequence

The CODENAME field is analyzed by the generated instructions for the presence of the various CODEVALUES as they appear in the operand portion of the source statement from left to right. Thus, if two different labels of CODE subsequent entries refer to the same actual value in the CODE field, and if both are listed, with different BRANCHES, in the operand of a single DECOD statement, the presence of that value in the CODENAME field will cause a transfer to the first of these BRANCHES.

### Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified:

**CODENAME NOT DEFINED BY A CODE**

This error message is issued if the first parameter is not the label of a CODE statement.

**CODEVALUE NOT DEFINED UNDER A CODE**

This error message is issued if one of the CODEVALUE entries is not a subsequent entry under any CODE header line.

**CODEVALUES AND BRANCHES NOT PAIRED**

This message will be issued if the number of parameters following CODENAME is odd.

**WARNING — CODEVALUE NOT DEFINED UNDER CODENAME**

This warning message is issued if one of the CODEVALUE entries has, as its first parameter record, a header label which is different from the CODENAME written in the source statement. Coding will be generated nevertheless.

**Examples**

The following are examples of acceptable coding for the DECOD statement. For each, the associated source-program entries are given, followed by the DECOD statement, coding generated in-line, and coding generated out-of-line.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	332	*		DECOD EXAMPLE 1.					
02	333	AREANAME	DA	5,0+X30				+0003250329	
03	334	CODENAME	CODE	6,6			66	0325	0000
04	335	CODEVALUE		1					
05	336	BRANCH1	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0330	-0100090000	
06	337	*							
07	338	ANYLABEL	DECOD	CODENAME, CODEVALUE, BRANCH1					
08	X	ANYLABEL	ZA1	CODENAME+X30				0331 +1330660000	
09	X		CA	+1				0332 -1500000335	
10	X		BE	BRANCH1				0333 -4100090330	
11	339	*							
12	340	*		THE FOLLOWING IS GENERATED OUT OF LINE					
13	341	*							
				LITERALS					
14	X			1			00	0334 +1	0334
15	X			+1	00002	00	0335	+1	0335

DECOD Example 1

The program will branch to BRANCH1 if the CODENAME field contains the condition specified for CODEVALUE, or will go to the next sequential instruction if it does not.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	345	*		DECOD EXAMPLE 2.					
02	346		DA	1				+0003250325	
03	347	CODENAME	CODE	0,9'			09	0325	0325
04	348	CODE1		'RED '					
05	349	CODE2		'GREEN'					
06	350	CODE3		'BLUE '					
07	351	BR1	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0326	-0100090000	
08	352	BR2	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0327	-0100090000	
09	353	BR3	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0328	-0100090000	
10	354	*							
11	355	ANYLABEL	DECOD	CODENAME, CODE1, BR1, CODE2, BR2,					
12	356			CODE3, BR3					
13	X	ANYLABEL	ZA1	CODENAME				0329 +1300090325	
14	X		CA	RED				0330 -1500090341	
15	X		BE	BR1	00002		0331	-4100090326	
16	X		CA	GREEN				0332 -1500090340	
17	X		BE	BR2				0333 -4100090327	
18	X		CA	BLUE				0334 -1500090339	
19	X		BE	BR3				0335 -4100090328	
20	357	*							
21	358	*		THE FOLLOWING IS GENERATED OUT OF LINE					
22	359	*							
				LITERALS					
23	X			'RED '	00003	09	0336	'7965640000	0336
24	X			'GREEN'		09	0337	'6779656575	0337
25	X			'BLUE '		09	0338	'6273846500	0338

DECOD Example 2

If the CODENAME field contains the value specified for CODE1, CODE2, or CODE3, the program will branch to BR1, BR2, or BR3, respectively. Otherwise the program will go to the next sequential instruction.



## LOGIC — Logical Decisions

LOGIC generates instructions to test whether a given expression is true or false and, according to the result, to set a switch or to branch to a designated location, or both.

**Source Program Format** The basic formats for the LOGIC statement in the source program are as follows:

Line	Label	Operation	OPERAND	Basic Au
3	5 6	15 16	20 21 25 30 35 40 45	50
0 1	ANY LABEL	LOGIC	SWITCH = EXPRESSION	
0 2	ANY LABEL	LOGIC	EXPRESSION, TRUEBR, FALSEBR	
0 3	ANY LABEL	LOGIC	SWITCH = EXPRESSION, TRUEBR, FALSEBR	
0 4				

In these examples, ANYLABEL is any symbolic label; it may be omitted. The entry LOGIC must be written exactly as shown. SWITCH represents any electronic, program, or digit switch to be turned ON if the expression in question is true, OFF if false. If the label of an index word or a DA, DC, or DSW header line or subsequent entry is written, the digits in the first location of the current area referenced will be treated as a bank of digit switches; i.e., they will all be turned ON or OFF. TRUEBR and FALSEBR represent the symbolic labels of instructions to which the program will branch if the expression is, respectively, true or false.

The first format causes the object program to set the switch as indicated and to continue with the next instruction of the source program after the LOGIC statement. The second format causes the object program to branch only, not to set a switch. The third format does both.

In the second and third formats listed above, one of the branches may be omitted. The object program will then proceed sequentially instead of taking the missing branch. For example, if the operand reads SWITCH=EXPRESSION, TRUEBR and the expression is false, the designated switch will be turned OFF and the next instruction of the source program will be executed. Care must be taken to enter the separating commas if TRUEBR is omitted since the generator always interprets the branch following the first comma as TRUEBR. The operand should read SWITCH=EXPRESSION, FALSEBR.

The table on the following page is a list of the parameters which are valid in the LOGIC statement as well as the way they are treated by LOGIC when they are used as a SWITCH or as an EXPRESSION.

### Logical Expressions

Logical expressions, in their simplest form, may consist of a single *logical variable*. *Logical operators* are used to connect two or more logical variables in order to form more complex logical expressions. *Logical punctuation* is used when necessary to clarify an ambiguous logical expression.

Parameter	Switch	Expression
Index Word	Bank of Switches	Operand Field
Alteration Switch	INVALID	Alteration Switch
Electronic Switch Undefined Term	Electronic Switch	Electronic Switch
DLINE Header Line DLINE Subsequent Entry CODE Header Line	INVALID	Operand Field
CODE Subsequent Entry	INVALID	Code Field
DTF Header Line	Bank of Switches	Operand Field
DA Header Line DC Header Line DSW Header Line	Digit Switch or Bank of Switches	Operand Field
DTF Subsequent Entry DA Subsequent Entry DC Subsequent Entry	Digit Switch or Bank of Switches	Digit Switch or Operand Field
DSW Subsequent Entry	Digit Switch	Digit Switch
Program Switch	Program Switch	Program Switch

Also, literals are allowed in relational expressions. Any parameters not listed above are invalid.

Logical variables may be of seven types; they can assume either one of two values, which are interpreted as true and false, respectively.

Type	Address	True Condition	False Condition
Electronic Switch	Symbolic	ON	OFF
Alteration Switch	"	ON	OFF
Program Switch	"	Plus	Minus
Digit Switch	"	>0	=0
Code Value	"	Code value present	Code value not present
Operand Field	"	Not all zeros	All zeros
Relational	(See below)	True	False

*Electronic Switch, Alteration Switch, Program Switch, Digit Switch.* These four switches require no explanation. Examples of the use of these switches with the LOGIC macro-instruction are included under "Examples." The use of an electronic

switch is illustrated in examples 6, 9, and 10; of an alteration switch in examples 7 and 10; of a program switch in examples 7, 9, and 10; of a digit switch in examples 1, 8, and 10.

*Code Values.* If a CODE field has been defined in accordance with the instructions outlined on page 50, it can be interrogated for the presence of a specific code value by referencing the symbolic name of that code value. If the 50 states of the Union are assigned the integers 1 through 50 in alphabetical order as code values under a CODE header line labeled STATE, the integer 1 will be assigned to ALABAMA. The word ALABAMA then functions as a logical variable, and in any given expression it will be regarded as true if the CODE field STATE contains a 1, false otherwise. The use of a code value is illustrated in examples 5 and 10 under "Examples."

*Operand Fields.* An operand field is a contiguous field of any length, not necessarily confined to a single word. If an operand field is used as a SWITCH to be set, it is treated as a bank of switches. However, if the field is greater than one word, only the first word of the field will be turned on and a warning message will be issued.

If an operand field is used as an EXPRESSION in the LOGIC statement, it may be any length. The field will be regarded as false if it contains all zeros, and true if one digit contains a value other than zero. The sign of the zeros does not matter. Alphameric zeros (@9090909090) will register as digits different from zero, but alphameric blanks (@0000000000) will be treated as zeros in this context. The use of an operand field as an EXPRESSION is illustrated in example 3 under "Examples."

*Relationals.* Relational expressions are comparisons between two numerical or two alphameric fields; the two types should not be mixed in a single comparison. The six comparison operators available and their respective numerical and alphameric meanings are as follows:

<i>Operator</i>	<i>Numerical Meaning</i>	<i>Alphameric Meaning</i>
G	is greater than	follows
NOT G	is not greater than	does not follow
E	is equal to	is identical to
NOT E	is not equal to	is not identical to
L	is less than	precedes
NOT L	is not less than	does not precede

A relational expression is formed by placing one of these operators between two fields and enclosing the resulting expression in parentheses. Care must be taken that exactly one blank separates fields and the operator.

When comparing two numerical fields, either field may be symbolic or literal. Numerical literals may be signed or unsigned; if unsigned, they will be interpreted as positive. Both automatic-decimal and floating-decimal fields are acceptable, and the modes may be freely mixed within a single relational expression.

The conditions and restrictions which apply to the use of the LOGIC macro-instruction are identical to those which apply to the COMP macro-instruction. These are as follows:

1. No automatic-decimal field may have more than twenty digits.
2. If an automatic-decimal number is compared to a floating-decimal number, only its first eight significant digits will affect the comparison.
3. If two automatic-decimal numbers are compared, their combined length after alignment of the decimal point may not exceed twenty digits. Any excess decimals beyond this length are disregarded in the comparison.

Comparison of numerical fields is strictly algebraic. Of two positive fields, the one with the greater absolute value will be regarded as larger of two negative numbers, the one with the smaller absolute value will be regarded as larger; any negative number is treated as smaller than any positive; zeros, whether positive or negative, will be treated as smaller than any other positive and greater than any other negative number.

For example, (AGE G 17) will be true if the number in the field referenced by AGE is larger than +17, false otherwise. The expression (TEMP NOT L - 13) is true if TEMP is 2, 0, -5, or -13, false if it is -13.5. Examples 2, 8, and 10 under "Examples" illustrate additional uses of relationals for comparing numerical fields.

When comparing two *alphameric fields*, either field may be symbolic or literal. Literals must be enclosed in @ signs (e.g., @NYC@). There is no limitation on the length of alphameric fields to be compared, except that literals may not exceed 120 characters. The relational expression tests for dictionary ordering, and the comparison operators must be reinterpreted to the alphameric meaning indicated in the chart above. Special characters will be included in this dictionary ordering according to the standard collating sequence given in the IBM Reference Manual "7070 Data Processing System."

For example, (INITIAL L @K@) will be regarded as true whenever the field referenced by the word INITIAL contains a letter that precedes K in the alphabet, or a special character whose two-digit numerical representation is less than 72. (GRADE NOT G @X@) will be false if the GRADE field contains Y or Z, true if it contains V, W, or X. (@12@ G @AB@) is true. @ @ will be regarded as less than any other alpha field.

### **Logical Operators**

Logical operators permit the construction of more complex expressions from logical variables. The LOGIC macro generator interprets three operators: NOT, AND, and OR.

*Not.* If NOT precedes an expression, it has the effect of changing its value to the opposite; if it precedes a true expression, the resulting expression is false, and vice versa. The use of NOT is illustrated in examples 6, 9, and 10 under "Examples."

*And.* If AND is placed between two expressions, a new and more complex expression results. This expression is true if, and only if, both of the component expressions are true; otherwise it is false. The use of AND is illustrated in examples 7, 9, and 10 under "Examples."

*Or.* If OR is placed between two expressions, a new and more complex expression results. This expression is true if at least one of the component expressions is true, possibly both. The compound expression, therefore, is false only if both of the components are false. The use of OR is illustrated in example 8 and 10 under "Examples."

### **Logical Punctuation (Parentheses)**

Expressions resulting from logical operations upon variables may in turn serve as components for larger expressions. Let capital letters represent electronic

switches; then the expression `A OR NOT B` will be true if `A` is on, or if `NOT B` is true, i.e., if `B` is off, or both. It will be false if and only if `A` is off and `B` is on.

The expression `NOT A AND B`, however, is ambiguous as it stands. It might have been constructed from `A` and `B` by first operating upon them with `AND`, and then prefacing the result with `NOT`; in that case the expression will be true if `B` is off, regardless of the status of `A`. On the other hand, it might have been built by placing `NOT` before `A`, and then operating upon the resulting expression and `B` with the operator `AND`; in that case, if `B` is off the expression is false. To provide the necessary distinction between these meanings, any compound logical expression must be enclosed in parentheses before it is operated upon again. This would yield `NOT (A AND B)` and `(NOT A) AND B`, respectively, for the two cases above. The use of parentheses is illustrated in examples 9 and 10 under "Examples."

To avoid excessive parenthesization, three conventions are adopted:

1. The operator `NOT` applies only to the shortest complete logical expression immediately to its right. Thus `NOT A OR B` is taken to mean `(NOT A) OR B`, since `A` is a complete logical expression. If `NOT (A OR B)` was intended, the parentheses would have to be explicitly written; then, since the operator `NOT` is followed by a left parenthesis, the shortest complete expression to its right is the entire parenthesized expression.
2. If the operators `AND` and `OR` occur in the same expression without intervening parentheses, the terms connected by `AND` will be understood to be parenthesized. Thus `A AND B OR C` will be taken to mean `(A AND B) OR C`. If `A AND (B OR C)` is intended, the parentheses must be written.
3. Repeated use of either `AND` or `OR` in the same logical expression is not ambiguous and need not be parenthesized. The expression `A OR B OR C` always yields the same result, whether treated as `(A OR B) OR C` or as `A OR (B OR C)`. However, because of a saving in object program time, `LOGIC` will deal with it as though it had the latter form. See "Left-Orientation," below.

The `LOGIC` macro generator will automatically interpret punctuated logical expressions in the sense of these conventions. It will not, however, reject clear, correct expressions in which parentheses are explicit that might have been suppressed.

## Processing Techniques

### Limitations on Length

Not more than 24 parameters may occur in any one `LOGIC` operand, counting the switch to be set, each branch, each logical variable, and each logical operator; in the case of repetitions, each occurrence is counted separately. A relational expression without `NOT` has three parameters; with `NOT`, four. Punctuation and address modifiers are not counted as parameters.

### Spacing and Punctuation

A blank must both precede and follow each `AND` and `OR`. If `NOT` precedes a parenthesized expression, no blank need intervene, but if it precedes some other operand, a space should be left blank. In relational expressions, no blanks should occur between the enclosing parentheses and the fields to be compared. The operator should be separated by one blank from each of the fields, and, where `NOT` appears, from each other. No blanks should occur on either side of the equal sign or on either side of the separating commas preceding branches.

### Address Modification

Addresses occurring in the operand portion of a `LOGIC` statement may be freely modified by indexing and address adjustment. The same is true of symbolic

fields in relational expressions though not of literal fields in the same expressions; this will, of course, require parentheses within parentheses. Address modification is illustrated in example 10 under "Examples."

### Left-Orientation

In programming complex logical expressions, placing the simpler terms or conditions on the left side of AND and OR operators will often result in a substantial saving of object program time. Thus, the logical expression

$$A \text{ OR NOT } B \text{ AND } (\text{AGE } G \text{ } 26) \quad (1)$$

will yield coding which may allow faster resolution than the logically equivalent version

$$(\text{AGE } G \text{ } 26) \text{ AND NOT } B \text{ OR } A \quad (2)$$

In the object program, the truth of the individual logical variables is evaluated from left to right in the order in which they appeared in the source statement. Thus, if A is true in a given run of the object program, the above expression as a whole is true, and if it was originally coded as in version (1), the determination of the truth of B and of (AGE G 26) is bypassed. In version (2), no such quick resolution can be obtained. (AGE G 26) will be the first logical variable to be interrogated. The outcome of this test would not constitute sufficient information for an evaluation of the truth of the expression as a whole. If the relational expression is true, B would have to be tested next, and if false, A.

### Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified:

#### EQUAL SIGN BEGINS INPUT — WILL IGNORE

This warning will be issued if the operand portion of the source statement begins with an equal sign. The sign will be disregarded, and no instructions will be generated to set a switch *unless* a second equal sign occurs in the statement following the first operand entry.

#### ILLEGAL PUNCTUATION MARK USED

The only punctuation that may validly appear in a LOGIC statement consists of an equal sign and a maximum of two commas as indicated under "Source Program Format," as well as parentheses, and plus, minus, and alpha signs. Other special characters may occur only inside alphameric literals. In case of violation, a NOP will be generated.

#### ILLEGAL TERM ENDS LOGICAL EXPRESSION

Some term that is neither a logical variable nor a right parenthesis ends the source statement, or immediately precedes the first separating comma before the branch entries. This condition may be caused if an unintended double blank is interpreted by the generator to mean that the operand portion is complete. A NOP will be generated.

#### ILLEGAL TERM PRECEDES A BINARY OPERATOR

A term that is neither a right parenthesis nor a logical variable precedes an AND or OR. A NOP will be generated.

#### ILLEGAL TERM PRECEDES A LEFT PAREN

Some term that is neither another left parenthesis, nor a logical operator, nor an equal sign, precedes a left parenthesis. A NOP will be generated.

#### ILLEGAL TERM PRECEDES A NOT

A term that is neither a logical operator, nor a left parenthesis, nor an equal sign, precedes a NOT. A NOP will be generated.

#### ILLEGAL TERM PRECEDES A RIGHT PAREN

Some term that is neither another right parenthesis, nor a logical variable, nor an address modifier, precedes a right parenthesis. A NOP will be generated.

#### ILLEGAL TERM PRECEDES PARAMETER *xx*

*xx* will be replaced by the number of a parameter which has been preceded by a term that is neither a logical operator, nor a left parenthesis, nor an equal sign. A NOP will be generated.

#### INVALID PARAMETER *xx*

This message will be issued if a parameter is not one specifically listed as valid under "Source Program Format." *xx* will be replaced by the number of the parameter at fault. A NOP will be generated.

#### NO BRANCH OR SWITCH TO BE SET IN INPUT

The coding fails to indicate what implementation the generated instructions are to initiate. A NOP will be generated.

#### NOTHING TO TEST IN LOGIC STATEMENT

The source statement does not contain an expression whose truth or falsity is to be determined. A NOP is generated.

#### PAREN. MISSING AROUND ARITH-REL

Either the left or the right parenthesis has been omitted around a relational expression. This message will also be issued if an attempt is made to run together more than one comparison since the generator expects to find a right parenthesis after the second field. For example, this message will be issued if (MIN L AVERAGE L MAX) is encountered since the generator expects to find a right parenthesis after AVERAGE.

#### PARENTHESIS NOT CLOSED

There is an excess of left parentheses, leaving at least one left parenthesis unpaired. A NOP will be generated.

#### TOO MANY RIGHT PARENTHESSES

There is an excess of right parentheses, leaving at least one right parenthesis unpaired. A NOP will be generated.

#### WILL SET SWITCHES IN FIRST LOCATION ONLY

This warning is issued if the SWITCH to be set is the label of a DA or DC header line or subsequent entry, and if the referenced area bridges words. Only those digits that lie in the first location will be affected by the generated instructions.

## Examples

The following are examples of acceptable coding for the LOGIC macro-instruction. For each, the associated source-program entries are given, followed by the LOGIC statement, coding generated in-line, and (where applicable) coding generated out-of-line.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	363	*	LOGIC EXAMPLE 1.						
02	364		DSW DIGITSW	00001	00	0325	+1000000000	0325	
03	365	FALSEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE				0326	-0100090000
04	366	*							
05	367	ANYLABEL	LOGIC DIGITSW,,FALSEBR						
06		X ANYLABEL	CD DIGITSW,0			0327	+0300000325		
07		X	BE FALSEBR			0328	-4100090326		
08	368	*							

LOGIC Example 1

If the digit switch is ON, the program will continue with the next instruction.

If it is OFF, it will branch to FALSEBR.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	372	*	LOGIC EXAMPLE 2.						
02	373		DA 1				+0003250325		
03	374	INCOME	0,4A						
04	375	TRUEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE				0325	
05	376	*		00001	04	0326	-0100090000	0325	
06	377	ANYLABEL	LOGIC SWITCH=(INCOME G 4800),TRUEBR,						
07		X ANYLABEL	ZA2 INCOME(0,4)			0327	+2300040325		
08		X	S2 +4800			0328	-2400030334		
09		X	BZ2 M.2			0329	+2000090333		
10		X	BM2 M.2			0330	-2000090333		
11		X M.1	ESN SWITCH	00002		0331	+6100100000		
12		X	B TRUEBR			0332	+0100090326		
13		X M.2	ESF SWITCH			0333	+6100200000		
14	378	*							
15	379	*	THE FOLLOWING IS GENERATED OUT OF LINE						
16	380	*	LITERALS						
17		X	+4800			03	0334 +4800	0334	

LOGIC Example 2

If the numerical field referenced by INCOME is greater than +4800, SWITCH will be turned ON and the program will branch to TRUEBR. If INCOME is equal to or less than +4800, SWITCH will be turned OFF and the program will continue sequentially. In this example, since the type of switch is not specifically designated, SWITCH is assumed to be electronic switch 1.



PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF	
01 384	*		LOGIC EXAMPLE 3.					
02 385		DA	1			+0003250325		
03 386	OPFIELD		0,9A		09	0325	0325	
04 387	TRUEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0326 -0100090000		
05 388	*							
06 389	ANYLABEL	LOGIC	OPFIELD,TRUEBR					
07	X ANYLABEL	ZAI	OPFIELD(0,9)			0327 +1300090325		
08	X	BZ1	M.1			0328 +1000090330		
09	X	B	TRUEBR			0329 +0100090326		
10	X M.1	NOP				0330 -0100000000		
11	X ORIGIN	CNTRL	*-1					
12 390	*							

#### LOGIC Example 3

If OPFIELD contains any digits other than zeros, the program will continue with the instruction labeled TRUEBR. If every digit is zero, the program continues sequentially. Alphameric blanks (@0000000000) register as zeros; alphameric zeros (@9090909090) register as digits other than zero.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF	
01 401	*		LOGIC EXAMPLE 4.					
02 402		DA	1			+0003250325		
03 403	SUBSCRIPT		5,6'		56	0325	0325	
04 404	FALSEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0326 -0100090000		
05 405	*							
06 406	ANYLABEL	LOGIC	SWITCH=(SUBSCRIPT NOT E 'J'),,FALSEBR					
07	X ANYLABEL	ZAA	M.5(0,1)			0327 +1600010334		
08	X	CA	SUBSCRIPT(0,1)			0328 -1500560325		
09	X	BL	M.1			0329 +4000090333		
10	X	BH	M.1			0330 -4000090333		
11	X M.2	ESF	SWITCH	00002		0331 +6100200000		
12	X	B	FALSEBR			0332 +0100090326		
13	X M.1	ESN	SWITCH			0333 +6100100000		
14 407	*							
15 408	*		THE FOLLOWING IS GENERATED OUT OF LINE					
16 409	*							
17	X	DC				+0003340334		
18	X M.5		'J'	00003 01	0334	'71	0334	

#### LOGIC Example 4

If the field referenced by SUBSCRIPT, which must be alphameric, contains a J, SWITCH will be turned OFF and the program will branch to FALSEBR. If SUBSCRIPT contains a character different from J, SWITCH will be turned ON and the program will continue sequentially.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	413	*		LOGIC EXAMPLE 5.					
02	414		DA	1				+0003250325	
03	415	STATE	CODE	0		00	0325		0325
04	416	OHIO		5					
05	417	TRUEBR	NOP		00001		0326	-0100090000	
06	418	FALSEBR	NOP				0327	-0100090000	
07	419	*							
08	420	ANYLABEL	LOGIC	SWITCH=OHIO,TRUEBR,FALSEBR					
09	X	ANYLABEL	CD	STATE,5			0328	+0300500325	
10	X		BE	M.1			0329	-4100090332	
11	X	M.2	ESF	SWITCH			0330	+6100200000	
12	X		B	FALSEBR	00002		0331	+0100090327	
13	X	M.1	ESN	SWITCH			0332	+6100100000	
14	X		B	TRUEBR			0333	+0100090326	
15	421	*							
16	422	*		THE FOLLOWING IS GENERATED OUT OF LINE					
17	423	*		LITERALS					
18	X			5		00	0334	+5	0334

LOGIC Example 5

Assume that a CODE declarative statement has established a one-digit CODE field labeled STATE, and that the code value corresponding to OHIO is 5. If the STATE field contains a 5, SWITCH will be turned ON and the program will branch to TRUEBR; otherwise, SWITCH will be turned OFF and the program will branch to FALSEBR.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	427	*		LOGIC EXAMPLE 6.					
02	428	PROGSW	NOP	REPRESENTS PROGRAM SWITCH	00001		0325	-0100090000	
03	429	*							
04	430	ANYLABEL	LOGIC	PROGSW=NOT REGISTERED					
05	X	ANYLABEL	BES	REGISTERED,M.2			0326	+6100000329	
06	X	M.1	MSP	PROGSW			0327	-0300910325	
07	X		B	M.3			0328	+0100090330	
08	X	M.2	MSM	PROGSW			0329	-0300610325	
09	X	M.3	NOP		00002		0330	-0100000000	
10	X	ORIGIN	CNTRL	*-1					
11	431	*							

LOGIC Example 6

If the electronic switch REGISTERED is ON, the program switch PROGSW will be turned OFF (i. e., its sign will be made minus); if REGISTERED is OFF, PROGSW will be turned ON (plus). The program will continue sequentially in either case.

PAGE AA	PROGRAM		7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	435	*		LOGIC EXAMPLE 7.					
02	436	ALTSW	EQU	1,SN					
03	437	PROGSW	NOP	REPRESENTS PROGRAM SWITCH	00001		0325	-0100090000	
04	438	TRUEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH			0326	-0100090000	
05	439	FALSEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH			0327	-0100090000	
06	440	*							
07	441	ANYLABEL	LOGIC	PROGSW AND ALTSW,TRUEBR,FALSEBR					
08		X ANYLABEL	CSM	PROGSW			0328	-0300600325	
09		X	BE	FALSEBR			0329	-4100090327	
10		X M.5	BAS	ALTSW,TRUEBR	00002		0330	+5100100326	
11		X	B	FALSEBR			0331	+0100090327	
12	442	*							

LOGIC Example 7

If both the program switch PROGSW and the alteration switch ALTSW are ON, the program will branch to TRUEBR. If either one or both of these switches are OFF, the program will continue with FALSEBR.

PAGE AA	PROGRAM		7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	446	*		LOGIC EXAMPLE 8.					
02	447		DSW	-DIGITSW	00001	00	0325	+0000000000	0325
03	448		DA	1				+0003260326	
04	449	NET		0,9		09	0326		0326
05	450	TRUEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH	00002		0327	-0100090000	
06	451	*							
07	452	ANYLABEL	LOGIC	DIGITSW OR (NET G 1200),TRUEBR					
08		X ANYLABEL	CD	DIGITSW,0			0328	+0300000325	
09		X	BH	TRUEBR			0329	-4000090327	
10		X M.5	ZA2	NET(0,9)			0330	+2300090326	
11		X	S2	+1200			0331	-2400030336	
12		X	BV2	**2	00003		0332	+2100090334	
13		X	BZ2	**3			0333	+2000090336	
14		X	BM2	**2			0334	-2000090336	
15		X	B	TRUEBR			0335	+0100090327	
16	453	*							
17	454	*		THE FOLLOWING IS GENERATED OUT OF LINE					
18	455	*		LITERALS					
19		X		+1200	03	0336	+1200		0336

LOGIC Example 8

If the digit switch DIGITSW is ON, or if the numerical field NET contains a number greater than +1200, or both, the program will continue with the instruction labeled TRUEBR. Only if DIGITSW is OFF and the NET field contains a number equal to or less than +1200 will the program continue sequentially.

PAGE AA		PROGRAM		7070 COMPILER SYSTEM VERSION QMY08, CHANGE LEVEL 00001.				PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	459	*		LOGIC EXAMPLE 9.					
02	460		DSW	-DIGITSW	00001	00	0325	+0000000000	0325
03	461	PROGSW	NOP	REPRESENTS PROGRAM SWITCH			0326	-0100090000	
04	462	*							
05	463	ANYLABEL	LOGIC	DIGITSW=NOT(ELECSW AND PROGSW)					
06		X ANYLABEL	BES	ELECSW,**2			0327	+6100000329	
07		X	B	M.1			0328	+0100090331	
08		X M.5	CSM	PROGSW			0329	-0300600326	
09		X	BH	M.2	00002		0330	-4000090333	
10		X M.1	ZA1	+1111111111			0331	+1300090335	
11		X	B	M.4			0332	+0100090334	
12		X M.2	ZA1	+0			0333	+1300000336	
13		X M.4	ST1	DIGITSW			0334	+1200000325	
14	464	*							
15	465	*		THE FOLLOWING IS GENERATED OUT OF LINE					
16	466	*		LITERALS					
17		X		+1111111111	00003	09	0335	+1111111111	0335
18		X		+0		00	0336	+0	0336

### LOGIC Example 9

If both the electronic switch ELECSW and the program switch PROGSW are ON, digit switch DIGITSW will be turned OFF. If either ELECSW or PROGSW is OFF, or both, DIGITSW will be turned ON. In either case, the program will continue sequentially.

PAGE AA		PROGRAM		7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.					PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	470	*		LOGIC EXAMPLE 10.					
02	471		DA	1				+0003250325	
03	472	MATRIX		0,9			09	0325	0325
04	473	FALSEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001			0326	-0100090000
05	474	*							
06	475	ANYLABEL	LQIC	SWITCH(XWORD+125)=(MATRIX(ROW-1) NOT E 1),, REMARKS					
07	476			FALSEBR(7-3) MAY BE USED.					
08	X	ANYLABEL	ZA2	MATRIX(0,9)-1+ROW				0327	+2301090324
09	X		S2	+1				0328	-2400000335
10	X		BV2	**2				0329	+2100090331
11	X		BZ2	**2				0330	+2000090332
12	X		B	M.1				0331	+0100090334
13	X	M.2	ESF	SWITCH+125+XWORD	00002			0332	+6102200125
14	X		B	FALSEBR-3+X7				0333	+0107090323
15	X	M.1	ESN	SWITCH+125+XWORD				0334	+6102100125
16	477	*							
17	478	*		THE FOLLOWING IS GENERATED OUT OF LINE					
18	479	*							
				LITERALS					
19	X			+1				00 0335	+1 0335

#### LOGIC Example 10

If the field defined as MATRIX, indexed by the index word ROW and address-adjusted by -1, is not equal to +1, the electronic switch addressed by SWITCH, indexed by XWORD and address-adjusted by +125, will be turned ON. Otherwise the switch will be turned OFF and the program will continue with the instruction located at FALSEBR incremented by the contents of index word 7 and decremented by 3.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.										PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF			
01	483	*		LOGIC EXAMPLE 11								
02	484		DA	8,,0+X87								
03	485	STATE	CODE	0				+0003250332				
04	486	OHIO		5		00	0325				0000	
05	487	NEWYORK		6								
06	488	MAINE		7								
07	489	TRUEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0333	-0100090000				
08	490	FALSEBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0334	-0100090000				
09	491	*										
10	492	ANYLABEL	LOGIC	SWITCH=OHIO,TRUEBR,FALSEBR								
11	X	ANYLABEL	CD	STATE,5								
12	X		BE	M.1				0335 +0387500000				
13	X	M.2	ESF	SWITCH				0336 -4100090339				
14	X		B	FALSEBR				0337 +6100200000				
15	X	M.1	ESN	SWITCH	00002		0338	+0100090334				
16	X		B	TRUEBR				0339 +6100100000				
17	493	*						0340 +0100090333				
18	494	*		THE FOLLOWING IS GENERATED OUT OF LINE								
19	495	*										
				LITERALS								
20	X			5								
21	X			6		00	0341	+5			0341	
22	X			7		11	0341	+ 6			0341	
						22	0341	+ 7			0341	

LOGIC Example 11

This example is similar to Example 5. Additional CODE values have been established corresponding to NEWYORK and OHIO under a DA in which both relative addressing and implicit indexing are used.

## ZSIGN—Branch on Test for Zero and Sign

ZSIGN generates instructions that will analyze a field or area for the presence of zeros and, if it is not zero, for its sign, and then branch accordingly.

### Source Program Format

The two basic formats for the ZSIGN statement in the source program are as follows:

Line	Label	Operation	OPERAND							Basic Autocoder				
			15	16	20	21	25	30	35	40	45	50	55	
0 1	ANY LABEL	ZSIGN	FIELD	ZEROBR	PLUSBR	MINUSBR	ALPHABR							
0 2	ANY LABEL	ZSIGN	FIELD	NOZERO	PLUSBR	MINUSBR	ALPHABR							
0 3														

In these examples, ANYLABEL represents any symbolic label; it may be omitted. The entries ZSIGN and NOZERO must be written exactly as shown. FIELD is the symbolic name of the subsequent entry defining the field to be tested. It may also be the header label of a declarative statement in which case the record area defined will be tested. ZEROBR, PLUSBR, MINUSBR, and ALPHABR are the symbolic labels of instructions to which the program will branch if the contents of the field or area are found to be, respectively, zero, plus, minus, or alpha. In the second format, the entry NOZERO instead of ZEROBR will prevent testing for a zero condition; branching then is purely according to sign.

The source program formats may be modified by the omission of one, two, or in the case of the first format, even three branches. In that case, the object program would take the next instruction following the ZSIGN statement instead of the missing branch. Separating commas must be entered if any branch except the last is omitted.

In the first format, if the field contains zeros, transfer will be to ZEROBR whether the zeros are plus or minus. Alpha blanks (of the form @000000000) will also cause a branch to ZEROBR. Alpha zeros (@909090909), on the other hand, will cause the program to transfer to ALPHABR (except for the marginal case in which the field has only one digit and that contains the zero digit of an alpha zero, in which case transfer is to ZEROBR.) In this format, if the field does not contain zeros, transfer will be according to the sign of the word in which the left-most digit of the field is contained.

In the second format, branching will be determined by the sign of the word in which the left-most digit of the field is contained.

In the following examples, assume that the field consists of the underlined digits:

<i>Contents of Storage</i>	<i>Branch Taken Using the First Format</i>	<i>Branch Taken Using the Second Format</i>
+00 <u>01234567</u>	PLUSBR	PLUSBR
-12345 <u>67890</u>	MINUSBR	MINUSBR
@836177 <u>6979</u>	ALPHABR	ALPHABR

<i>Contents of Storage</i>	<i>Branch Taken Using the First Format</i>	<i>Branch Taken Using the Second Format</i>
+0000000000	ZEROBR	PLUSBR
-0000000000 +0000000005	ZEROBR	MINUSBR
@0000000000	ZEROBR	ALPHABR
@9090909090	ALPHABR	ALPHABR
@9090909090	ZEROBR	ALPHABR
+0001234567 -1234567890	PLUSBR	PLUSBR
@8272656274 +0000000000	ALPHABR	ALPHABR

When a **ZSIGN** statement references the label of a **DA** header line, coding will be generated to cause the following:

1. If the **DA** header line does not specify a relative address and implicit indexing, the *first* record area defined will be tested as specified.
2. If the **DA** header line specifies a relative address and implicit indexing, the *current* record area (as determined by the contents of the implicit index word) will be tested as specified.

When a **ZSIGN** statement references any other declarative statement header line, the entire area will be tested as specified.

## Processing Techniques

### Limitations on Length

The number of parameters is fixed by the format, subject only to the omission of one or more branches. There is no limit to the size of the field to be tested.

### Address Modification

Indexing and address adjustment are permitted on all symbolic addresses.

### Error and Warning Messages

The following error and warning messages are issued during assembly under the conditions indicated:

#### BRANCH TO NON-IMPERATIVE INSTRUCTION

One of the branch addresses to which the object program may transfer is not the label of an imperative instruction, but, for example, that of a **DA** subsequent entry.

#### FIELD MISSING

This message will be issued if **FIELD** is omitted from the operand, e.g., if the operand portion of the source statement begins with a comma.

#### FIELD UNACCEPTABLE

The **FIELD** entry contains a symbolic address of an entity that cannot meaningfully be tested for zero contents or sign, e.g., an alteration switch.

#### NO BRANCHES GIVEN

All four branches in the first format, or all three branches in the second format, have been omitted. A **NOP** will be generated to aid in patching if this condition was unintended. It should be noted, however, that this instruction will accomplish exactly what a literal interpretation of the source statement requires, i.e., the program will take the next instruction in any case.



#### UNLIKELY — ALL BRANCHES IDENTICAL

All four branches in the first format, or all three branches in the second format, are identical. In case the source statement is of the first format, coding for the zero test will have been generated. Then the procedure is the same as for the second format, i.e., a NOP will be generated, followed by an unconditional Branch instruction to the required branch. This again allows for patching while implementing a strict interpretation of the source statement.

#### Examples

The following are examples of acceptable coding for the ZSIGN macro-instruction. For each, the associated source-program entries are given, followed by the ZSIGN statement, and coding generated in-line.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001, PAGE AA							
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	501	*		ZSIGN EXAMPLE 1					
02	502		DA	1				+0003250325	
03	503	ANYFIELD		3,9A4.3		39	0325		0325
04	504	ZEROBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0326	-0100090000	
05	505	PLUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0327	-0100090000	
06	506	MINUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0328	-0100090000	
07	507	ALPHABR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0329	-0100090000	
08	508	*							
09	509	ANYLABEL	ZSIGN	ANYFIELD,ZEROBR,PLUSBR,MINUSBR,ALPHABR					
10	X	ANYLABEL	ZA1	ANYFIELD(0,6)			0330	+1300390325	
11	X		BZ1	ZEROBR	00002		0331	+1000090326	
12	X		CSM	ANYFIELD			0332	-0300600325	
13	X		BL	ALPHABR			0333	+4000090329	
14	X		BL	MINUSBR			0334	-4100090328	
15	X		B	PLUSBR			0335	+0100090327	
16	510	*							

ZSIGN Example 1

A field, less than one word in length, will be examined first for the presence of zeros and, if it does not contain all zeros, for a plus, minus, or alpha sign. A branch to the appropriate instruction will be made based on what is found.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001, PAGE AA							
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	514	*		ZSIGN EXAMPLE 2					
02	515		DA	1				+0003250325	
03	516	FIELDX		0,9A		09	0325		0325
04	517	PLUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0326	-0100090000	
05	518	MINUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0327	-0100090000	
06	519	*							
07	520	ANYLABEL	ZSIGN	FIELDX,PLUSBR,MINUSBR					
08	X	ANYLABEL	ZA1	FIELDX(0,9)			0328	+1300090325	
09	X		BZ1	M.1			0329	+1000090333	
10	X		CSM	FIELDX			0330	-0300600325	
11	X		BH	PLUSBR	00002		0331	-4000090326	
12	X		BL	MINUSBR			0332	-4100090327	
13	X	M.1	NOP				0333	-0100000000	
14	X	ORIGIN	CNTRL	*-1					
15	521	*							

ZSIGN Example 2

The program will examine the one-word field, FIELDX, for a plus or minus sign and branch accordingly. If FIELDX only contains zeros or has an alpha sign, the program will continue with the next sequential instruction.

PAGE	AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.				PAGE	AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	525	*		ZSIGN EXAMPLE 3					
02	526		DA	1				+0003250331	
03	527	AFIELD		03,69		39	0325		0325
04	528	BRANCHX	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0332	-0100090000	
05	529	BRANCHZ	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0333	-0100090000	
06	530	*							
07	531	ANYLABEL	ZSIGN	AFIELD, BRANCHX, BRANCHZ, BRANCHZ, BRANCHX					
08	X	ANYLABEL	XL	MACREG.01,+0000000005			0334	+4500010343	
09	X		ZAA	AFIELD			0335	+1600390325	
10	X	M.2	AA	AFIELD(7,16)+MACREG.01			0336	+1701090326	
11	X		BITX	MACREG.01,M.2	00002		0337	+4900010336	
12	X		BVI	**2			0338	+1100090340	
13	X		BZI	BRANCHX			0339	+1000090332	
14	X		CSM	AFIELD			0340	-0300600325	
15	X		BL	BRANCHX			0341	+4000090332	
16	X		B	BRANCHZ	00003		0342	+0100090333	
17	532	*							
18	533	*		THE FOLLOWING IS GENERATED OUT OF LINE					
19	534	*		LITERALS					
20	X			+0000000005		09	0343	+0000000005	0343

## ZSIGN Example 3

The program will examine a field greater than one word. If AFIELD contains zeros or if the sign of the word in which the left-most digit of AFIELD is contained is alpha, a branch to BRANCHX will be made.

If the sign is plus or minus, the branch will be to BRANCHZ.

PAGE	AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.				PAGE	AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	536	*		ZSIGN EXAMPLE 4					
02	537		DA	1				+0003250331	
03	538	AFIELD		03,69		39	0325		0325
04	539	NOZERO	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001		0332	-0100090000	
05	540	BRANCHZ	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0333	-0100090000	
06	541	BRANCHX	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0334	-0100090000	
07	542	*							
08	543	ANYLABEL	ZSIGN	AFIELD, NOZERO, BRANCHZ, BRANCHZ, BRANCHX					
09	X	ANYLABEL	CSM	AFIELD			0335	-0300600325	
10	X		BL	BRANCHX			0336	+4000090334	
11	X		B	BRANCHZ	00002		0337	+0100090333	
12	544	*							

## ZSIGN Example 4

A field greater than one word will be examined for sign only. The program will continue with the instruction located at BRANCHZ if the sign of the word which contains the left-most digit of AFIELD is plus or minus, or BRANCHX if the sign is alpha.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	548	*		ZSIGN EXAMPLE 5					
02	549		DA	4,RDW,0+X13				+0003250332	
03		X			00001		0325	+0003290329	0325
04		X					0326	+0003300330	0326
05		X					0327	+0003310331	0327
06		X					0328	-0003320332	0328
07	550	ANYFIELD		3,9A4.3			39	0329	0000
08	551	ZEROBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00002		0333	-0100090000	
09	552	PLUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0334	-0100090000	
10	553	MINUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0335	-0100090000	
11	554	ALPHABR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE			0336	-0100090000	
12	555	*							
13	556	ANYLABEL	ZSIGN	ANYFIELD,ZEROBR,PLUSBR,MINUSBR,ALPHABR					
14		X ANYLABEL	ZA1	ANYFIELD(0,6)+X13				0337 +1313390000	
15		X	BZ1	ZEROBR	00003			0338 +1000090333	
16		X	CSM	ANYFIELD+X13				0339 -0313600000	
17		X	BL	ALPHABR				0340 +4000090336	
18		X	BE	MINUSBR				0341 -4100090335	
19		X	B	PLUSBR				0342 +0100090334	
20	557	*							

### ZSIGN Example 5

The actual ZSIGN test is the same as in Example 1. In this case, however, four blocked records are in the defined area. The test will be made only on the field ANYFIELD in the current record area, as determined by the contents of the implicit index word.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	561	*		ZSIGN EXAMPLE 6					
02	562	ANYLABEL	DA	3,,0+INDEXWORD				+0003250333	
03	563	FIELDA		0,9A6.4			09	0325	0000
04	564	FIELDB		10,15'			05	0326	0001
05	565	FIELD C		16,28			69	0326	0001
06	566	ZEROBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE	00001			0334	-0100090000
07	567	PLUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE				0335	-0100090000
08	568	MINUSBR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE				0336	-0100090000
09	569	ALPHABR	NOP	REPRESENTS FIRST INSTRUCTION OF BRANCH ROUTINE				0337	-0100090000
10	570	*							
11	571	SOMELABEL	ZSIGN	ANYLABEL,ZEROBR,PLUSBR,					
12	572			MINUSBR,ALPHABR					REMARKS MAY BE USED.
13	X	SOMELABEL	ZAA	0(0,9)+INDEXWORD				0338	+1601090000
14	X		AA	0(10,19)+INDEXWORD	00002			0339	+1701090001
15	X		AA	0(20,29)+INDEXWORD				0340	+1701090002
16	X		BV1	*+2				0341	+1100090343
17	X		BZ1	ZEROBR				0342	+1000090334
18	X		CSM	0+INDEXWORD				0343	-0301600000
19	X		BL	ALPHABR	00003			0344	+4000090337
20	X		BE	MINUSBR				0345	-4100090336
21	X		B	PLUSBR				0346	+0100090335
22	573	*							

#### ZSIGN Example 6

Since the label of the DA header line is addressed in the ZSIGN statement, the entire area through digit position 29 will be tested for zeros. Since implicit indexing has been used, the current record will be tested. If the record does not contain all zeros, the sign of the word which contains the left-most digit of ANYLABEL (FIELD A) will be checked, and appropriate branches will be made.

## SETSW — Set Switch

SETSW generates instructions to set one or more digit, electronic, or program switches to an on or off condition.

### Source Program Format

The basic format for the SETSW statement in the source program is as follows:

Line	Label	Operation	OPERAND					
3	5 6	15 16	20 21	25	30	35	40	45
0 1	ANY LABEL	SETSW	±SWA	, ±SWB	, ±SWC	, etc.		
0 2								

In this example, ANYLABEL is any symbolic label; it may be omitted. The entry SETSW must be written exactly as shown. The entries SWA, SWB, etc., represent names of switches that are to be turned ON or OFF. If the switch is a program or a digit switch, the name must be symbolic; the label of a series of digit switches established by a DSW (Define Switch) statement may be used. If the switch is electronic, either its symbolic or its actual one- or two-digit name may be used. The various types may be freely intermingled within the operand portion of a single SETSW statement.

Switches preceded by a plus sign will be turned ON; those preceded by a minus sign will be turned OFF. The sign may be omitted, in which case it will be assumed to be plus. If a series of switches established by a DSW instruction is included, all of the component switches will be turned ON or OFF according to the sign (or its absence) preceding the label.

Commas must be entered between successive entries in the operand portion of the SETSW statement.

### Processing Techniques

#### Limitations on Length

The operand portion may include 94 parameters. Any excess will be ignored.

#### Address Modification

All symbolic addresses may be modified by indexing and address adjustment.

### Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified:

INSTRUCTION NOT PROGRAM SWITCH PAR. xx

An attempt has been made to turn a program switch ON or OFF (to make plus or minus an instruction with operation code 01) and the instruction addressed fails to have this operation code. xx will be replaced by the number of the parameter at fault. Coding will be generated to implement the sign adjustment nevertheless.

INVALID SWITCH, PARAMETER *xx*

An attempt has been made to set hardware switches, such as alteration switches, by use of the SETSW statement, or other unacceptable operand entries have been made. *xx* will be replaced by the number of the parameter at fault.

**Example**

The following is an example of acceptable coding for the SETSW macro-instruction. The associated source-program entries are given, followed by the SETSW statement, the coding generated in-line, and the coding out-of-line.

PAGE AA		PROGRAM				7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.				PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	577	*		SETSW EXAMPLE 1						
02	578		DSW	-DIGSWB,+DIGSWA	00001	01	0325	+0100000000	0325	
03	579	PROGSWA	NOP	REPRESENTS PROGRAM SWITCH			0326	-0100090000		
04	580	PROGSWB	NOP	REPRESENTS PROGRAM SWITCH			0327	-0100090000		
05	581	*								
06	582	ANYLABEL	SLTSW	ELECSWA,-21,PROGSWA,+DIGSWA,					REMARKS MAY	
07	583			-PROGSWB,-DIGSWB					BE USED.	
08	X	ANYLABEL	ESN	ELECSWA			0328	+6100100000		
09	X		ESF	21			0329	+6300200000		
10	X		MSP	PROGSWA	00002		0330	-0300910326		
11	X		ZA1	+1111111111			0331	+1300090336		
12	X		ST1	DIGSWA			0332	+1200110325		
13	X		MSM	PROGSWB			0333	-0300610327		
14	X		ZA2	+0			0334	+2300000337		
15	X		ST2	DIGSWB	00003		0335	+2200000325		
16	584	*								
17	585	*		THE FOLLOWING IS GENERATED OUT OF LINE						
18	586	*								
				LITERALS						
19	X			+1111111111			09 0336	+1111111111	0336	
20	X			+0			00 0337	+0	0337	

SETSW Example 1

Electronic switch A, program switch A, and digit switch(es) A will be turned ON. Electronic switch 21, program switch B, and digit switch(es) B will be turned OFF.



## ZERO—Zero Storage

ZERO generates instructions to replace the contents of fields or areas with zeros or blanks.

**Source Program Format** The basic format for the zero statement in the source program is as follows:

Line	Label	Operation	OPERAND
3	56	1516	2021 25 30 35 40 45
0 1	ANY LABEL	ZERO	FIELD A, AREA B, FIELD C, etc.
0 2			

ANY LABEL is any symbolic label; it may be omitted. The entry ZERO must be written exactly as shown. FIELD A, AREA B, etc., may be the symbolic names of any defined fields or areas. Areas, numerical fields and alphanumeric fields may be freely intermingled.

## Processing Techniques

### Limitations on Length

The operand portion of the ZERO macro-instruction may contain up to 94 entries. There is no limit on the size of the fields named.

### Address Modification

All symbolic addresses may be modified by indexing and address adjustment.

### The Effect of ZERO

*Fields and Areas.* *DLINE* areas and all fields defined as alphanumeric will be replaced by the double-digit representation of blanks. Whole words will be made alpha; the sign of partial words will not be altered.

In all other fields or areas, whole words will be replaced by plus zeros; partial words will be replaced by ZEROS but the sign will not be altered.

The following examples illustrate the effect of ZERO on various fields. The field addressed is underlined.

No.	Field Definition	Before ZERO	After ZERO
1.	0, 9A10.0	-8342168900	+0000000000
2.	0, 9@	@9192616263	@0000000000
3.	0, 13A	-9342168900 -1869123456	+0000000000 -0000123456
4.	0, 13A	-8342168900 +1869123456	+0000000000 +0000123456
5.	0, 15@	@9192616262 @6162636465	@0000000000 @0000006465
6.	0, 15@	-9192616263 -6162636465	@0000000000 -0000006465
7.	4, 8A5.0	-1234012345	-1234000005
8.	4, 8A5.0	@8283909195	@8283000005

In example 8, machine difficulties may arise when an attempt is made to print out the zeroed word, since the combination 05 has no meaning in double-digit, alphameric code.

*Declarative Statement Header Lines.* When ZERO references the label of a DA header line, coding will be generated to cause the following:

1. If the DA header line does not specify a relative address and implicit indexing, the *first* record area defined will be set to plus zeros.
2. If the DA header line specifies a relative address and implicit indexing, the *current* record area (as determined by the contents of the implicit index word) will be set to plus zeros.

If ZERO references the label of a DLINE header line, coding will be generated to cause the entire area, including constants (if any), to be set to blanks.

If ZERO references the label of a DRDW, a warning message will be issued, but coding will be generated to cause the first RDW generated (*not* the area it defines) to be set to plus zeros.

When ZERO references the label of any other declarative statement header line, coding will be generated to cause the entire area to be set to plus zeros.

*Instructions.* If ZERO references the label of an instruction, a warning message will be issued, but coding will be generated to cause the instruction to be set to plus zeros.

## Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified:

ALPHA BLANKS INTO UNDEFINED PAR. *xx*

The *xx* will be replaced by the parameter number of the operand without defined characteristics. This field will be filled with alphameric blanks.

ATTEMPTING TO ZERO HARDWARE. PAR. *xx*

The field to be zeroed has been defined by means of an EQU line as a hardware device. A NOP will be generated.

NO FIELD SIZE. PAR. *xx*

The parameter record of the operand entry whose number replaces the *xx* of the message does not indicate the size of the field to be zeroed out. A NOP will be generated.

ZEROING DC. PAR. *xx*

The label of a DC header line has been used as an operand. The parameter number will replace the *xx* of the warning message. Coding will be generated nevertheless.

ZEROING INSTRUCTION. PAR. *xx*

This warning message, with the parameter number of the faulty entry in place of the *xx*, will be issued whenever an attempt is made to zero out an instruction, whether symbolic machine or macro. Coding will be generated, however.

## Examples

The following are examples of acceptable coding for the ZERO macro-instruction. For each, the associated source-program entries are given, followed by the ZERO statement, coding generated in-line, and coding generated out-of-line.

PAGE AA		PROGRAM				7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.				PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	601	*		ZERO EXAMPLE 1						
02	602		DA	1				+0003250438		
03	603	FIELDA		00,99'	09		0325		0325	
04	604	FIELDDB		100,103A	03		0335		0335	
05	605	FIELDDC		110,129A	09		0336		0336	
06	606	FIELDDB		132,1131A	29		0338		0338	
07	607	*								
08	608	ANYLABEL	ZERO	FIELDA,FIELDDB,FIELDDC,				REMARKS MAY		
09	609			FIELDDB				BE USED.		
10		X ANYLABEL	ZA1	' '	00001		0439	+1300010455		
11		X	ST1	FIELDA(0,9)			0440	+1200090325		
12		X	XZA	MACREG.1,FIELDA			0441	+4600010325		
13		X	RS	MACREG.1,M.2			0442	+6500010452		
14		X	STD1	FIELDDB(0,3)			0443	-1200030335		
15		X	ZA2	+0	00002		0444	+2300000454		
16		X	ST2	FIELDDC(0,9)			0445	+2200090336		
17		X	ST2	FIELDDC(10,19)			0446	+2200090337		
18		X	STD1	FIELDDB(0,7)			0447	-1200290338		
19		X	ST2	FIELDDB(8,17)			0448	+2200090339		
20		X	XZA	MACREG.1,FIELDDB+1	00003		0449	+4600010339		
21		X	RS	MACREG.1,M.3			0450	+6500010453		
22		X	STD1	FIELDDB(998,999)			0451	-1200010438		
23	610	*								
24	611	*		THE FOLLOWING IS GENERATED OUT OF LINE						
25	612	*								
26		X M.2	DRDW	-FIELDDB+1,FIELDDB+9			0452	-0003260334		
27		X M.3	DRDW	-FIELDDB+2,FIELDDB+99			0453	-0003400437		
				LITERALS						
28		X		+0	00004	00	0454	+0	0454	
29		X		' '	00005	01	0455	'00	0455	

### ZERO Example 1

FIELDDB will be filled with alphameric blanks. FIELDDB will be filled with zeros; the sign of the word will not be changed. FIELDDB will be filled with plus zeros. The first and last words in which digits of FIELDDB occur will have the portion occupied by FIELDDB replaced by zeros; the sign of these words will not be changed. The rest of FIELDDB will be filled with plus zeros.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	616	*		ZERO EXAMPLE 2					
02	617		DA	2,RDW,0+INDEXWORD				+0003250352	
03		X			00001		0325	+0003270339	0325
04		X					0326	-0003400352	0326
05	618	FIELD A		00,95'			09	0327	0000
06	619	FIELD B		96,103A			69	0336	0009
07	620	FIELD C		110,129A			09	0338	0011
08	621	*							
09	622	ANYLABEL	ZERO	FIELD A, FIELD B, FIELD C					
10		X	ZA1	' '	00002		0353	+1300010366	
11		X	ST1	FIELD A(0,9)+INDEXWORD			0354	+1201090000	
12		X	XZS	MACREG.2, FIELD A+1+INDEXWORD			0355	-4601020001	
13		X	XSN	MACREG.2, FIELD A+8+INDEXWORD			0356	+4801020008	
14		X	XZA	MACREG.1, FIELD A+INDEXWORD			0357	+4601030000	
15		X	RS	MACREG.1, MACREG.2	00003		0358	+6500030002	
16		X	STD1	FIELD A(90,95)+INDEXWORD			0359	-1201050009	
17		X	STD1	FIELD B(0,3)+INDEXWORD			0360	-1201690009	
18		X	STD1	FIELD B(4,7)+INDEXWORD			0361	-1201030010	
19		X	ZA2	+0			0362	+2300000365	
20		X	ST2	FIELD C(0,9)+INDEXWORD	00004		0363	+2201090011	
21		X	ST2	FIELD C(10,19)+INDEXWORD			0364	+2201090012	
22	623	*							
23	624	*		THE FOLLOWING IS GENERATED OUT OF LINE					
24	625	*							
				LITERALS					
25		X		+0		00	0365	+0	0365
26		X		' '	00005	01	0366	'00	0366

ZERO Example 2

In the current record area, as determined by the contents of the implicit index word, FIELD A will be filled with alphameric blanks and FIELD B will be filled with zeros. Note that the last word that contains digits of FIELD A also contains digits of FIELD B. The sign of this word, therefore, will not be changed. The sign of the last word which contains digits of FIELD B will also remain unchanged since it is only a partial word. The contents of FIELD C will be replaced with plus zeros.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION QMY08, CHANGE LEVEL 00001, PAGE AA							
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	629	*		ZERO EXAMPLE 3					
02	630	AREANAME	DA	1				+0003250349	
03	631	FIELD A		00,99'	09		0325		0325
04	632	FIELD B		100,103A	03		0335		0335
05	633	FIELD C		104,129A	49		0335		0335
06	634	FIELD D		130,249A	09		0338		0338
07	635	*							
08	636	ANYLABEL	ZERO	AREANAME					
09	X	ANYLABEL		ZA2 +0	00001		0350	+2300000355	
10	X		ST2	AREANAME(0,9)			0351	+2200090325	
11	X		XZA	MACREG.1,AREANAME			0352	+4600010325	
12	X		RS	MACREG.1,M.2			0353	+6500010354	
13	637	*							
14	638	*		THE FOLLOWING IS GENERATED OUT OF LINE					
15	639	*							
16	X	M.2	DRDW	-AREANAME+1,AREANAME+24			0354	-0003260349	
				LITERALS					
17	X			+0	00002	00	0355	+0	0355

## ZERO Example 3

Since the ZERO statement references the label of the DA header line,  
the entire record area defined will be filled with plus zeros.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION QMY08, CHANGE LEVEL 00001, PAGE AA							
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	643	*		ZERO EXAMPLE 4					
02	644	AREANAME	DA	2,RDW.0+INDEXWORD				+0003250348	
03	X				00001		0325	+0003270337	0325
04	X						0326	-0003380348	0326
05	645	FIELD A		00,99'	09		0327		0000
06	646	FIELD B		100,103A	03		0337		0010
07	647	*							
08	648	ANYLABEL	ZERO	AREANAME					
09	X	ANYLABEL		ZA2 +0	00002		0349	+2300000355	
10	X		ST2	0(0,9)+INDEXWORD			0350	+2201090000	
11	X		XZS	MACREG.2,0+1+INDEXWORD			0351	-4601020001	
12	X		XSN	MACREG.2,0+10+INDEXWORD			0352	+4801020010	
13	X		XZA	MACREG.1,0+INDEXWORD			0353	+4601030000	
14	X		RS	MACREG.1,MACREG.2	00003		0354	+6500030002	
15	649	*							
16	650	*		THE FOLLOWING IS GENERATED OUT OF LINE					
17	651	*							
				LITERALS					
18	X			+0	00		0355	+0	0355

## ZERO Example 4

The ZERO statement in this example references the label of a DA header  
line which specifies a relative address and implicit indexing. The con-  
tents of the current record area will be filled with plus zeros.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	655	*		ZERO EXAMPLE 5					
02	656	DRDNAME	DRDW	AREANAME					
03	X	DRDNAME	DRDW	+AREANAME,AREANAME+24	00001		0325	+0003270351	
04	X		DRDW	-AREANAME+25,AREANAME+49			0326	-0003520376	
05	657	*							
06	658	AREANAME	DA	2,,0+INDEXWORD				+0003270376	
07	659	FIELDA		00,99'		09	0327		0000
08	660	FIELDB		100,103A		03	0337		0010
09	661	FIELD C		104,129A		49	0337		0010
10	662	FIELD D		130,249A		09	0340		0013
11	663	*							
12	664	ANYLABEL	ZERO	DRDNAME					
13	X	ANYLABEL	ZA2	+0	00002		0377	+2300000379	M
14	X		STD2	DRDNAME(0,9)			0378	-2200090325	
15	665	*							
16	666	*		THE FOLLOWING IS GENERATED OUT OF LINE					
17	667	*							
				LITERALS					
18	X			+0		00	0379	+0	0379

ERROR MESSAGE LIST

PG/LN MESSAGE

AA 12 ZEROING INSTRUCTION. PAR.000001000A

ZERO Example 5

The warning message shown will be issued since the ZERO statement references the label of a DRDW. Coding has been generated, however, to cause the first RDW generated to be filled with plus zeros.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	671	*		ZERO EXAMPLE 6					
02	672	LINENAME	DLINE		00005			+0003250348	
03	673			1'1'	00006	01	0325	'91	0325
04	674			10'TOTAL'		89	0326	' 83	0326
05		X				07	0327	'76836173	0327
06	675			18'S'		45	0328	' 25	0328
07	676	GROSSAMT		19X,XXX.ZZ)DR,CR		69	0328		0328
08	677	CHECKAMT		60\$X,XXZ.ZZ),C		89	0336		0336
09	678	ITEMNAME		80,94	00007	89	0340		0340
13	679	FLVAR		95F		89	0343		0343
14	680			120'R	00008	89	0348	' 80	0348
15	681	*							
16	682	ANYLABEL	ZERO	CHECKAMT					
17		X ANYLABEL	ZA2	+0	00009		0349	+2300000354	
18		X	STD2	CHECKAMT(0,1)			0350	-2200890336	
19		X	ZA1	' '			0351	+1300010355	
20		X	ST1	CHECKAMT(2,11)			0352	+1200090337	
21		X	STD1	CHECKAMT(12,19)			0353	-1200070338	
22	683	*							
23	684	*		THE FOLLOWING IS GENERATED OUT OF LINE					
24	685	*		LITERALS					
25		X		+0	00010	00	0354	+0	0354
26		X		' '	00011	01	0355	'00	0355

## ZERO Example 6

The contents of the DLINE field CHECKAMT is replaced with alphameric blanks. The sign of whole words in the field are set to alpha; the sign of partial words are not altered.

```

PAGE AA      PROGRAM
7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001. PAGE AA
LN CDREF LABEL OP OPERAND CDNO FD LOC INSTRUCTION REF
01 689 * ZERO EXAMPLE 7
02 690 LINENAME DLINE 00002 +0003250333
03 691 AMOUNTONE 195X,XXX,ZZ 69 0328 0328
04 692 AMOUNTTWO 305X,XXX,XXX,XX 89 0330 0330
05 693 *
06 694 ANYLABEL ZERO LINENAME
07 X ANYLABEL ZA1 ' ' 00003 0334 +1300010339
08 X ST1 LINENAME(0,9) 0335 +1200090325
09 X XZA MACREG.1,LINENAME 0336 +4600010325
10 X RS MACREG.1,M.2 0337 +6500010338
11 695 *
12 696 * THE FOLLOWING IS GENERATED OUT OF LINE
13 697 *
14 X M.2 DRDW -LINENAME+1,LINENAME+8 0338 -0003260333
LITERALS
15 X ' ' 00004 01 0339 '00 0339

```

ZERO Example 7

The entire area, including constants, will be filled with alphameric blanks since the label of the DLINE header line is referenced in the ZERO statement.

```

PAGE AA      PROGRAM
7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001. PAGE AA
LN CDREF LABEL OP OPERAND CDNO FD LOC INSTRUCTION REF
01 6901 * ZERO EXAMPLE 8
02 6902 AREANAME DA 1 00002 +0003250331
03 6903 FIELDA 00,23' 09 0325 0325
04 6904 FIELDB 30,35A 05 0328 0328
05 6905 FIELDC 45,55 59 0329 0329
06 6906 65 55 0331 0331
07 6907 ANYLABEL ZERO AREANAME
08 X ANYLABEL ZA2 +0 00001 0332 +2300000337
09 X S12 AREANAME(0,9) 0333 +2200090325
10 X XZA MACREG.1,AREANAME 0334 +4600010325
11 X RS MACREG.1,M.2 0335 +6500010336
12 6908 *
13 6909 * THE FOLLOWING IS GENERATED OUT OF LINE
14 6910 *
15 X M.2 DRDW -AREANAME+1,AREANAME+6 0336 -0003260331
LITERALS
16 X +0 00002 00 0337 +0 0337

```

ZERO Example 8

Since the label of the DA header line is referenced by the ZERO statement, the entire area through digit position 69 will be filled with plus zeros.



## FILL—Fill Storage

FILL generates instructions to replace the contents of fields or areas with a specified constant.

### Source Program Format

The basic format for the FILL statement in the source program is as follows:

Line	Label	Operation	OPERAND				Basic Autocoder				
3	5	15	20	25	30	35	40	45	50	55	60
01	ANY LABEL	FILL	FIELD A	AREA 1	WITH	ih	AND	FIELD B	WITH	@Z@	etc.
02											

ANY LABEL is any symbolic label; it may be omitted. The entry FILL and the AND and WITH separators must be written exactly as shown. FIELD A, AREA 1, etc., may be the symbolic names of any defined fields or areas. Areas, numerical fields, and alphanumeric fields may be freely intermingled.

### Processing Techniques

#### Limitations on Length

The operand portion of the FILL macro-instruction may contain up to 94 entries, including areas and fields to be filled, the WITH and AND separators, and the characters to be inserted. There is no limit on the size of the fields named.

#### Address Modification

All symbolic addresses may be modified by indexing and address adjustment.

#### The Effect of FILL

*Fields and Areas.* The sign of the words in the field or area to be filled will not affect the generated instructions. Each word in the entire field or area, and any word in which a segment of the field may appear, will be set to the sign of the filling constant. It is thus possible to introduce invalid alpha combinations in the following two cases:

1. The field specified is numerical and occupies part of a word(s) and is being filled with an alphanumeric character.
2. A field with an odd number of digits, or a field that begins in an odd-numbered position, is being filled with an alphanumeric character.

In case 2, a warning message will be issued.

If two fields are specified in the same word, and each field is to be filled with a different constant, the sign of the word is determined by the sign of the constant which fills a field last.

The following examples illustrate the effect of FILL on various fields. The field addressed is underlined>.

No.	Field Definition	Before FILL	Filling Constant	After FILL
1.	00,09A10.0	-1234567890	+9	+9999999999
2.	00,05A4.2	-1234560000	+1	+1111110000
3.	05,14	+6789523721 +4376411111	@Y@	@6789588888@8888811111
4.	08,08	+1234567890	-1	-1234567190
5.	00,15@	@0061626364 @6564630000	+9	+9999999999 +9999990000
6.	00,03	+1234567890	-1	-1111567890
	04,09	-1111567890	+2	+1111222222

In example 3, invalid, double-digit combinations (58, 81, and 11) are introduced. These may cause machine difficulties when an attempt is made to print out these words.

*Declarative Statement Header Lines.* When FILL references the label of a DA header line, coding will be generated to cause the following:

1. If the DA header line does not specify a relative address and implicit indexing, the *first* record area defined will be filled with the specified value.
2. If the DA header line specifies a relative address and implicit indexing, the *current* record area (as determined by the contents of the implicit index word) will be filled with the specified value.

If FILL references the label of a DLINE header line, coding will be generated to cause the entire area, including constants (if any), to be filled.

If FILL references the label of a DRDW, a warning message will be issued, but coding will be generated to cause the first RDW generated (*not* the area it defines) to be filled.

When FILL references the label of any other declarative statement header line, coding will be generated to cause the entire area to be filled.

*Instructions.* If FILL references the label of an instruction, a warning message will be issued, but coding will be generated to cause the instruction to be filled.

## Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified:

ATTEMPTING TO FILL HARDWARE. PAR. xx

The field to be filled has been defined by means of an EQU line as a hardware device. A NOP will be generated.

FILLING INSTRUCTION. PAR. xx

This warning message, with the parameter number of the faulty entry in place of the xx, will be issued whenever an attempt is made to fill an instruction, whether symbolic machine or macro. Coding will be generated, however.

NO FIELD SIZE. PAR. xx

The parameter record of the operand entry, whose number replaces the xx of the message, does not indicate the size of the field to be filled. A NOP will be generated.

WARNING. INVALID ALPHA MAY BE INTRODUCED.

An alphameric character is filling a field with an odd number of digits, or a field that begins in an odd-numbered position.

## Examples

The following are examples of acceptable coding for the `FILL` macro-instruction. For each, the associated source-program entries are given, followed by the `FILL` statement, coding generated in-line, and coding generated out-of-line.

```

PAGE AA      PROGRAM      7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001. PAGE AA
LN CDREF LABEL      OP      OPERAND      CDNO FD LOC INSTRUCTION      REF
01 902 *              FILL EXAMPLE 1
02 903              DA      1              +0003250328
03 904 FIELDA        06,09A          69 0325              0325
04 905 FIELDB        10,19A          09 0326              0326
05 906 FIELDC        28,31A          89 0327              0327
06 907 *
07 908 ANYLABEL     FILL  FIELDA WITH +1 AND FIELDB WITH -1 AND FIELDC WITH 'Y' M
08      X ANYLABEL   ZA1  +1111111111          00001      0329 +1300090336
09      X              ST1  FIELDA(0,3)          0330 +1200690325
10      X              ZA1  -1111111111          0331 +1300090337
11      X              ST1  FIELDB(0,9)          0332 +1200090326
12      X              ZA1  'YYYYY'          0333 +1300090338
13      X              ST1  FIELDC(0,1)          00002      0334 +1200890327
14      X              ST1  FIELDC(2,3)          0335 +1200010328
15 909 *
16 910 *              THE FOLLOWING IS GENERATED OUT OF LINE
17 911 *
                LITERALS
18      X              +1111111111          09 0336 +1111111111      0336
19      X              -1111111111          09 0337 -1111111111      0337
20      X              'YYYYY'          00003 09 0338 '8888888888      0338

```

ERROR MESSAGE LIST

PG/LN MESSAGE

AA 07 WARNING-INVALID ALPHA MAY BE INTRODUCED

FILL Example 1

FIELDA and FIELDB will be filled with 1s and the sign of the words in which each appears will be set to plus and minus respectively. FIELDC will be filled with Ys and, since it bridges words, the signs of the two words in which it appears will be set to alpha.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.										PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF			
01	914	*	FILL	EXAMPLE 2								
02	915		DA	2				+0003250332				
03	916	AFIELD		00,09'			09	0325		0325		
04	917	BFIELD		10,15'			05	0326		0326		
05	918	CFIELD		20,23'			03	0327		0327		
06	919	DFIELD		30,33'			03	0328		0328		
07	920	*										
08	921	ANYLABEL	FILL	AFIELD WITH ' ' AND BFIELD,CFIELD WITH +1 AND DFIELD								
09	922			WITH -1								
10		X ANYLABEL	ZA1	' '	00001		0333	+1300090342				
11		X	ST1	AFIELD(0,9)			0334	+1200090325				
12		X	ZA1	+1111111111			0335	+1300090340				
13		X	ST1	BFIELD(0,5)			0336	+1200050326				
14		X	ST1	CFIELD(0,3)			0337	+1200030327				
15		X	ZA1	-1111111111	00002		0338	+1300090341				
16		X	ST1	DFIELD(0,3)			0339	+1200030328				
17	923	*										
18	924	*		THE FOLLOWING IS GENERATED OUT OF LINE								
19	925	*										
				LITERALS								
20		X		+1111111111			09	0340	+1111111111	0340		
21		X		-1111111111			09	0341	-1111111111	0341		
22		X		' '	00003	09	0342	'0000000000		0342		

### FILL Example 2

In the first record area defined, AFIELD will be filled with alphameric blanks. BFIELD and CFIELD will be filled with 1s and the sign of the word in which each appears will be set to plus. DFIELD will also be filled with 1s, but the sign of the word will be set to minus.

PAGE AA	PROGRAM					7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.	PAGE AA	
LN	CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF
01	928	*		FILL EXAMPLE 3				
02	929		DA	2,RDW,0+INDEXWORD			+0003250332	
03		X			00001	0325	+0003270329	0325
04		X				0326	-0003300332	0326
05	930	A		00,01		01 0327		0000
06	931	B		10,11		01 0328		0001
07	932	C		20,21		01 0329		0002
08	933	*						
09	934	ANYLABEL	FILL	A WITH +9 AND B WITH -9 AND C WITH 'Z'				
10		X ANYLABEL	ZAL	+9999999999	00002	0333	+1300090339	
11		X	ST1	A(0,1)+INDEXWORD		0334	+1201010000	
12		X	ZAL	-9999999999		0335	+1300090340	
13		X	ST1	B(0,1)+INDEXWORD		0336	+1201010001	
14		X	ZAL	'ZZZZZ'		0337	+1300090341	
15		X	ST1	C(0,1)+INDEXWORD	00003	0338	+1201010002	
16	935	*						
17	936	*		THE FOLLOWING IS GENERATED OUT OF LINE				
18	937	*						
				LITERALS				
19		X		+9999999999		09 0339	+9999999999	0339
20		X		-9999999999		09 0340	-9999999999	0340
21		X		'ZZZZZ'	00004	09 0341	'8989898989	0341

FILL Example 3

In the current record area, as determined by the contents of the implicit index word, fields A and B will be filled with 9s and the sign of the word in which each appears will be set to plus and minus respectively. Field C will be filled with Zs and the sign of the word will be set to alpha.

```

PAGE AA      PROGRAM
LN CDREF LABEL      OP      OPERAND
7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001. PAGE AA
CDNO FD LOC INSTRUCTION REF
01 940 *          FILL EXAMPLE 4
02 941 AREANAME DA      1
03 942 FIELDA      00,03A      +0003250327
04 943 FIELDB      10,13'      03 0325      0325
05 944 FIELDDC     20,23      03 0326      0326
06 945 *          03 0327      0327
07 946 ANYLABEL FILL AREANAME WITH +0
08 X ANYLABEL ZAI +000000000
09 X ST1 AREANAME(0,9)      00001      0328 +1300090332
10 X ST1 AREANAME(10,19)   0329 +1200090325
11 X ST1 AREANAME(20,29)   0330 +1200090326
12 947 *          0331 +1200090327
13 948 *          THE FOLLOWING IS GENERATED OUT OF LINE
14 949 *          LITERALS
15 X              +0000000000      09 0332 +0000000000      0332

```

## FILL Example 4

Since the FILL statement references the label of the DA header line, the entire record area through digit position 29 will be filled with 0s and the signs of the three words affected will be set to plus.

```

PAGE AA      PROGRAM
LN CDREF LABEL      OP      OPERAND
7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001. PAGE AA
CDNO FD LOC INSTRUCTION REF
01 952 *          FILL EXAMPLE 5
02 953 AREANAME DA      2, RDW*0+INDEXWORD
03 X
04 X
05 954 FIELDA      00,03A      00001      0325 +0003250332
06 955 FIELDB      10,13'      0326 +0003270329      0325
07 956 FIELDDC     20,23      03 0327      0326 -0003300332      0326
08 957 *          03 0328      0000
09 958 ANYLABEL FILL AREANAME WITH -0
10 X ANYLABEL ZAI -000000000
11 X ST1 0(0,9)+INDEXWORD   00002      0333 +1300090337
12 X ST1 0(10,19)+INDEXWORD 0334 +1201090000
13 X ST1 0(20,29)+INDEXWORD 0335 +1201090001
14 959 *          0336 +1201090002
15 960 *          THE FOLLOWING IS GENERATED OUT OF LINE
16 961 *          LITERALS
17 X              -0000000000      09 0337 -0000000000      0337

```

## FILL Example 5

Since the label of the DA header line is addressed in the FILL statement, the entire area through digit position 29 will be affected. Since implicit indexing has been used, the current record area will be filled with zeros and the signs of the three words affected will be set to minus.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	964	*		FILL EXAMPLE 6					
02	965	AREANAME	DA	2,0+INDEXWORD				+0003250344	
03	966	AFIELD		00,19A		09	0325		0000
04	967	BFIELD		20,49I		09	0327		0002
05	968	CFIELD		50,99		09	0330		0005
06	969	*							
07	970	DRDNAME	DRDW	AREANAME					
08	X	DRDNAME	DRDW	+AREANAME,AREANAME+9	00001		0345	+0003250334	
09	X		DRDW	-AREANAME+10,AREANAME+19			0346	-0003350344	
10	971	*							
11	972	ANYLABEL	FILL	DRDNAME WITH +0					
12	X	ANYLABEL	ZAI	+0000000000			0347	+1300090349	
13	X		ST1	DRDNAME(0,9)			0348	+1200090345	
14	973	*							
15	974	*		THE FOLLOWING IS GENERATED OUT OF LINE					
16	975	*							
				LITERALS					
17	X			+0000000000		09	0349	+0000000000	0349

ERROR MESSAGE LIST

FILL Example 6

PG/LN MESSAGE  
AA 11 FILLING INSTRUCTION

The warning message shown will be issued since the FILL statement references the label of a DRDW. Coding has been generated, however, to cause the first RDW generated to be filled with 0s and the sign set to plus.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	978	*		FILL EXAMPLE 7					
02	979	LINENAME	DLINE		00002			+0003250334	
03	980	GROSSAMT		10\$XX,ZZZ.ZZ		89	0326		0326
04	981	TOTEXPENSE		25\$XX,XXX.ZZ		89	0329		0329
05	982	NETAMT		40\$XX,ZZZ.ZZ		89	0332		0332
06	983	*							
07	984	ANYLABEL	FILL	TOTEXPENSE WITH +0					
08	X	ANYLABEL	ZAI	+0000000000	00003		0335	+1300090339	
09	X		ST1	TOTEXPENSE(0,1)			0336	+1200890329	
10	X		ST1	TOTEXPENSE(2,11)			0337	+1200090330	
11	X		ST1	TOTEXPENSE(12,19)			0338	+1200070331	
12	985	*							
13	986	*		THE FOLLOWING IS GENERATED OUT OF LINE					
14	987	*							
				LITERALS					
15	X			+0000000000		09	0339	+0000000000	0339

FILL Example 7

The DLINE field TOTEXPENSE will be filled with 0s and the signs of the words affected will be set to plus.



7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001, PAGE AA

PAGE AA	PROGRAM	LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	990	*				FILL EXAMPLE 8					
02	991	LINENAME		DLINE			00002			+0003250333	
03	992					1' 1'	00003	09	0325	'0000000000	0325
04	993	CUSTOMER				10,29		89	0326		0326
09	994	AMOUNT				30\$XXX,ZZZ,ZZ)DR,CR		89	0330		0330
10	995	*									
11	996	ANYLABEL		FILL		LINENAME WITH +0					
12		X ANYLABEL		ZAI		+0000000000	00004		0334	+1300090339	
13		X		ST1		LINENAME(0,9)			0335	+1200090325	
14		X		XZA		MACREG.01,LINENAME			0336	+4600010325	
15		X		RS		MACREG.01,M.1			0337	+6500010338	
16	997	*									
17	998	*				THE FOLLOWING IS GENERATED OUT OF LINE					
18	999	*									
19		X M.1		DRDW		-LINENAME+1,LINENAME+8			0338	-0003260333	
						LITERALS					
20		X				+0000000000	00005	09	0339	+0000000000	0339

FILL Example 8

The entire area, including constants, will be filled with 0s and the signs

of the words affected will be made plus since the label of the DLINE

header line is referenced in the FILL statement.

## EDMOV—Edit and Move Data

EDMOV generates instructions to transfer data between specified fields in storage and to edit them to conform to the format of the field to which they are moved.

### Source Program Format

The basic format of the EDMOV statement in the source program is the following:

Line	Label	Operation	OPERAND													
			Basic						Autocoder							
3	56	15	16	20	21	25	30	35	40	45	50	55	60	65	Autoc	
01	ANYLABEL	EDMOV	FROMFIELD1	TO	TOFIELD1	,	FROMFIELD2	TO	TOFIELD2	,,	etc.,					
02																

ANYLABEL is any symbolic label; it may be omitted. The entry EDMOV must be written exactly as shown.

FROMFIELD entries may be literals or symbolic addresses of areas from which data is to be moved. The data of FROMFIELD may be alphameric, automatic-decimal, floating-decimal, mixed, or of unspecified characteristics. Automatic-decimal and floating-decimal data may be in single- or double-digit representation; data which is mixed or of unspecified characteristics is treated in the same way as alphameric data.

TOFIELD entries are symbolic addresses of areas to which the edited data is to be moved. A TOFIELD area may be any field defined as alphameric, automatic-decimal, or floating-decimal, or may be a print format defined in a DLINE subsequent entry. Automatic-decimal and floating-decimal data may be in a single- or double-digit representation.

The operator TO must be written exactly as shown, preceded and followed by a blank. Each entry must have both a FROMFIELD and a TOFIELD. If more than one entry is made in the operand of a statement, commas must be used to separate them.

### Processing Techniques

#### Limitations on Length

A maximum of four entries may be written in the operand. Field size is restricted as follows:

Field	Size
Automatic-decimal, single-digit	not more than 20 digits
Automatic-decimal, double-digit	not more than 20 characters (40 digits)
Floating-decimal, single-digit	exactly one word
Floating-decimal, double-digit	exactly two words
Alphameric, mixed, unspecified	unrestricted

#### Other Limitations

DA, DC, and DLINE header lines cannot be used as parameters in an EDMOV statement.

## Address Modification

### The Effect of EDMOV

All symbolic addresses may be modified by indexing and address adjustment.

All alphameric, mixed, and unspecified fields will be passed on to the MOVE macro generator for processing. No editing is done.

Each of the twenty-two possible types of numerical editing is accomplished by a sequence of steps selected from the following nine basic types of conversion:

- |                                    |                                    |
|------------------------------------|------------------------------------|
| 1. Automatic-decimal, single-digit | to Automatic-decimal, double-digit |
| 2. Floating-decimal, single-digit  | to Floating-decimal, double-digit  |
| 3. Automatic-decimal, double-digit | to Automatic-decimal, single-digit |
| 4. Floating-decimal, double-digit  | to Floating-decimal, single-digit  |
| 5. Automatic-decimal, old format   | to Automatic-decimal, new format   |
| 6. Automatic-decimal               | to Floating-decimal                |
| 7. Floating-decimal                | to Automatic-decimal               |
| 8. Automatic-decimal, double-digit | to Print format                    |
| 9. Floating-decimal, double-digit  | to Print format                    |

The chart on page 207 shows the sequence for each of the twenty-two types of editing.

The rules governing each of the nine types of conversion are as follows:

1. *Automatic-Decimal, Single-Digit to Automatic-Decimal, Double-Digit.* Conversion is accompanied by sign control; positive numbers will show a 6 in the next-to-last digit, negative numbers a 7. The sign position of the converted number will have an @. If the result is stored in part of a word, the sign of the entire word will be set to @. The following examples illustrate this conversion:

<i>Before Conversion</i>	<i>After Conversion</i>
+7627	@bb97969267
-4502	@bb94959072

2. *Floating-Decimal, Single-Digit to Floating-Decimal, Double-Digit.* FROMFIELD must occupy exactly one word; TOFIELD will occupy exactly two words, both of which will have an @ sign. Sign control will be indicated in digit 8 of the second word; a 6 indicates a positive number, 7 a negative number. The following examples illustrate this conversion:

<i>Before Conversion</i>	<i>After Conversion</i>
+5212345678	@9592919293@9495969768
-4587654321	@9495989796@9594939271

3. *Automatic-Decimal, Double-Digit to Automatic-Decimal, Single-Digit.* Conversion is accompanied by sign sensing. The sign of TOFIELD will be set to plus if the next-to-last digit of FROMFIELD is different from 7, minus if it is 7. If the result is stored in part of a word, the sign of the entire word will be set to the sign of the result. The converse of the examples included under 1, above, illustrates this conversion.
4. *Floating-Decimal, Double-Digit to Floating-Decimal, Single-Digit.* FROMFIELD must occupy exactly two words; TOFIELD will occupy one word. Conversion is accompanied by sign sensing. The sign of TOFIELD will be set to plus if digit 8 of the second word of FROMFIELD is different from 7, minus if it is 7. The converse of the examples included under 2, above, illustrates this conversion.
5. *Automatic-Decimal, Old Format to Automatic-Decimal, New Format.* Four cases are distinguished:

## S E Q U E N C E   O F   E D I T I N G

FROMFIELD	TOFIELD	T Y P E   O F   C O N V E R S I O N									
		Automatic-decimal, single-digit to Automatic-decimal double-digit	Floating-decimal, single-digit to Floating-decimal, double-digit	Automatic-decimal, double-digit to Automatic-decimal, single-digit	Floating-decimal, double-digit to Floating-decimal, single-digit	Automatic-decimal, old format to Automatic-decimal, new format	Automatic-decimal to Floating-decimal	Floating-decimal to Automatic-decimal	Automatic-decimal, double-digit to Print format	Floating-decimal, double-digit to Print format	
Automatic-decimal, single-digit	Automatic-decimal, single-digit					1					
	Automatic-decimal, double-digit	2				1					
	Floating-decimal, single-digit						1				
	Floating-decimal, double-digit		2				1				
	Automatic-decimal, print format	1							2		
	Floating-decimal, print format		2				1			3	
Automatic-decimal, double-digit	Automatic-decimal, single-digit			1				2			
	Automatic-decimal, double-digit	3		1				2			
	Floating-decimal, single-digit			1				2			
	Floating-decimal, double-digit		3	1				2			
	Automatic-decimal, print format								1		
	Floating-decimal, print format		3	1			2				4
Floating-decimal, single-digit	Automatic-decimal, single-digit							1			
	Automatic-decimal, double-digit	2						1			
	Floating-decimal, double-digit		1								
	Automatic-decimal, print format	2						1	3		
Floating-decimal, double-digit	Floating-decimal, print format		1								2
	Automatic-decimal, single-digit				1			2			
	Automatic-decimal, double-digit	3			1			2			
	Floating-decimal, single-digit				1						
	Automatic-decimal, print format	3			1			2	4		
Floating-decimal, print format	Floating-decimal, print format										1

- a. TOFIELD has more *decimal* places than FROMFIELD. These new decimal places will be filled with zeros. For example, a field whose automatic-decimal format is 3.2 is converted to a field whose format is 3.4. If the field contains 123.45 before conversion, it will contain 123.4500 after conversion.
- b. TOFIELD has fewer *decimal* places than FROMFIELD. Extraneous decimals will be truncated after rounding. For example, a field whose automatic-decimal format is 2.3 is converted to a field whose format is 2.2. If the field contains 55.467 before conversion, it will contain 55.47 after conversion; if it contains 55.464 before, it will contain 55.46 after.
- c. TOFIELD has more *integer* places than FROMFIELD. These new integer places will be filled with high-order zeros. For example, a field whose automatic-decimal format is 2.3 is converted to a field whose format is 4.3. If the field contains 56.125 before conversion, it will contain 0056.125 after conversion.
- d. TOFIELD has fewer *integer* places than FROMFIELD. A warning message will be issued during assembly that high-order digits may be lost. For example, a field whose automatic-decimal format is 4.1 is converted to a field whose format is 2.1. If the field contains 1545.7 before conversion, it will contain 45.7 after conversion.

Combinations of these conditions will cause all of the indicated actions to be taken.

6. *Automatic-Decimal to Floating-Decimal.* The first eight significant digits will be converted; others will be truncated without rounding. An automatic-decimal number is converted to a standard 7070 normalized, floating-decimal word. For example,  $-123.456789$  is converted to  $-5312345678$ .
7. *Floating-Decimal to Automatic-Decimal.* Four cases are distinguished:
  - a. TOFIELD can accommodate the entire converted field. Any excess decimal places or integer places are filled with zeros. For example, if a field which contains  $+5287654321$  is converted to a field whose automatic-decimal format is 4.7, the result will be  $+0087.6543210$ .
  - b. An integer of the converted number falls to the left of the high-order place of TOFIELD. This is an overflow condition, and the overflow latch on Accumulator 1 will be set ON. No warning message can be issued during assembly since this condition cannot be predicted on the basis of floating-decimal format alone. All digits that can be accommodated in their proper places will be stored. For example, if a field which contains  $-5412345678$  is converted to a field whose automatic-decimal format is 3.5, the result will be  $-234.56780$  and the overflow latch of Accumulator 1 will be turned ON.
  - c. The first digit of the converted number falls into one of the places of TOFIELD, but the decimals cannot be accommodated. Excess decimals will be truncated after rounding. For example, if a field which contains  $+4823456789$  is converted to a field whose automatic-decimal format is 1.7, the result will be  $+0.0023457$ ; if a field which contains  $-5398765432$  is converted to a field whose format is 3.3, the result will be  $-987.654$ .
  - d. The first digit of the converted number falls to the right of the low-order place of the result field. Since the decimal value of the number is too small to register in the format of the result field, the field will be set equal to zero. For example, if a field which contains  $+3575757575$  is converted to a field whose automatic-decimal format is 2.3, the result will be  $+00.000$ .

8. *Automatic-Decimal, Double-Digit to Print Format.* Editing is performed to fit data to a DLINE print image. (See page 65.) The necessary commas, decimal points, and other characters will be inserted.
9. *Floating-Decimal, Double-Digit to Print Format.* Editing is performed to convert a floating-decimal number to DLINE print format which is  $\pm nn \pm .xxxxxxx$ , where  $\pm nn$  is a two-digit exponent  $\pm .xxxxxxx$  is an eight-digit number. The value of the number is  $\pm .xxxxxxx$  multiplied by  $10^{\pm nn}$ . For example, @9591999897@9695949372 (Which is  $-5198765432$  in single-digit form) will be printed as +01-.98765432, representing the number  $-.98765432 \times 10^1$ .

## Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified.

### BLANK PARAMETER *xx*

A parameter has been omitted. Its number will replace the *xx* of the message. A NOP will be generated.

### HIGH-ORDER DIGITS LOST OF PARAMETER *xx*

The field in which the edited data is to be stored has a format that will cause loss of integer digits on the left. Such digits as can be accommodated in their proper places will be stored.

### PRINT SUPPRESSED IF ALL ZERO

A print-line format has been written in such a way that all numerical positions are marked by Xs. A zero value for this field will result in no print at all.

### TO-FIELD NOT ALPHA. PARAMETER *xx*

An attempt has been made to move an alphameric, mixed, or unspecified field to a field that is not alpha. The parameter number of the TOFIELD will replace the *xx* of the message. The field will be moved but not edited.

### UNACCEPTABLE PARAMETER *xx*

The *xx* will be replaced by the number of a parameter that is not one of the types listed as acceptable under "Source Program Format." A NOP will be generated.

## Examples

The following are examples of acceptable coding for the EDMOV macro-instruction. For each, the associated source-program entries are given, followed by the EDMOV statement coding generated in-line and coding generated out-of-line.

PAGE AA	PROGRAM	LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
		01	701	*		EDMOV EXAMPLE 1					
		02	702	SOMELABEL	DLINE		00001			+0003250328	
		03	703	FIELDDB		10X,XXX.ZZ)		89	0326		0326
		04	704		DA	1				+0003290329	
		05	705	FIELDA		00,05A4.2		05	0329		0329
		06	706	*							
		07	707	ANYLABEL	EDMOV	FIELDA TO FIELDDB					
		08	X	ANYLABEL	ZA2	FIELDDB(0,5)	00002		0330	+2300050329	
		09	X		ST2	COMAREA.A+2			0331	+2200090352	
		10	X		XZA	MACREG.01,COMAREA.A+2			0332	+4600010352	
		11	X		ENB	MACREG.01,EDMOV02.A			0333	+5700010349	
		12	X		ZA3	COMAREA.A(8,9)			0334	+3300890350	
		13	X		ST3	FIELDDB(0,1)	00003		0335	+3200890326	
		14	X		BZ3	*+2			0336	+3000090338	
		15	X		ZA3	'.'			0337	+3300230353	
		16	X		ST3	FIELDDB(2,3)			0338	+3200010327	
		17	X		ZA3	COMAREA.A(10,15)			0339	+3300050351	
		18	X		ST3	FIELDDB(4,9)	00004		0340	+3200270327	
		19	X		BZ3	*+2			0341	+3000090343	
		20	X		ZA3	'.'			0342	+3300010353	
		21	X		ST3	FIELDDB(10,11)			0343	+3200890327	
		22	X		ZA3	COMAREA.A(16,19)			0344	+3300690351	
		23	X		ST3	FIELDDB(12,15)	00005		0345	+3200030328	
		24	X		ZA3	'9'			0346	+3300450353	
		25	X		ST3	FIELDDB(12,12)			0347	+3200000328	
		26	X		ST3	FIELDDB(14,14)			0348	+3200220328	
		27	708	*							
		28	709	*		THE FOLLOWING IS GENERATED OUT OF LINE					
		29	710	*							
		30	X	EDMOV02.A	DRDW	-COMAREA.A,COMAREA.A+1			0349	-0003500351	
		31	X	COMAREA.A	DA					+0003500352	
						LITERALS					
		32	X			'.'	00006	01	0353	'15	0353
		33	X			'.'		23	0353	' 35	0353
		34	X			'9'		45	0353	' 99	0353

## EDMOV Example 1

The automatic-decimal field FIELDDB is edited to the print format specified in the DLINE entry.

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	714	*		EDMOV EXAMPLE 2					
02	715		DA	1				+0003250325	
03	716	AUTODECNO		00,09A8.2	09		0325		0325
04	717		DA	1				+0003260326	
05	718	FLTPTNO		00,09F	09		0326		0326
06	719	*							
07	720	ANYLABEL	EDMOV	AUTODECNO TO FLTPTNO					
08		X ANYLABEL	ZA1	+0	00001		0327	+1300000341	
09		X	ZA2	AUTODECNO(0,9)			0328	+2300090325	
10		X	ZA3	+0000000068			0329	+3300090340	
11		X	BLX	94,FLOT2.A			0330	+0200940334	
12		X	ZST1	FLTPTNO			0331	-1100090326	
13	721	*							
14	722	*		THE FOLLOWING IS GENERATED OUT OF LINE					
15	723	*							
16		X FLOT1.A	SLC1	MACREG.1	00002		0332	+5000011300	
17		X	B	**2			0333	+0100090335	
18		X FLOT2.A	SLC	MACREG.1			0334	-5000010300	
19		X	BZ1	0+X94			0335	+1094090000	
20		X	S3	MACREG.1(4,5)			0336	-3400450001	
21		X	SR1	2	00003		0337	+5000001002	
22		X	STD3	9991(0,1)			0338	-3200019991	
23		X FLOT3.A	B	0+X94			0339	+0194090000	
			LITERALS						
24		X		+0000000068		09	0340	+0000000068	0340
25		X		+0		00	0341	+0	0341

EDMOV Example 2

The automatic-decimal number is edited to the floating-decimal format.



LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	727	*		EDMOV EXAMPLE 3.					
02	728	SOMELABEL	DLINE		00001			+0003250328	
03	729	AMT		10\$X,XXX.ZZ)DR,CR		89	0326		0326
04	730		DA	2,0+INDEXWORD				+0003290330	
05	731	AMTFIELD		3,9A4.3		39	0329		0000
06	732	*							
07	733	ANYLABEL	EDMOV	AMTFIELD TO AMT					
08	X	ANYLABEL	ZA2	AMTFIELD(0,6)+INDEXWORD	00002		0331	+2301390000	
09	X		SRR2	1			0332	+5000002101	
10	X		BM2	M*6			0333	-2000090362	
11	X		ZA3	' DR'			0334	+3300090377	
12	X	M*3	ST3	AMT(18,21)			0335	+3200690328	
13	X		ST2	COMAREA.A+2	00003		0336	+2200090374	
14	X		XZA	MACREG*01,COMAREA.A+2			0337	+4600020374	
15	X		ENA	MACREG*01,EDMOV02.A			0338	+5600020364	
16	X		ZA3	' '			0339	+3300090375	
17	X		ST3	AMT(0,1)			0340	+3200890326	
18	X		ST3	AMT(2,9)	00004		0341	+3200070327	
19	X		SLC2	MACREG*02			0342	+5000032300	
20	X		ZA3	'S'			0343	+3300230378	
21	X		B	M*4-4+MACREG*02			0344	+0103090361	
22	X	M*5	ST3	AMT(0,1)			0345	+3200890326	
23	X		ZA3	COMAREA.A(8,9)	00005		0346	+3300890372	
24	X		ST3	AMT(2,3)			0347	+3200010327	
25	X		ZA3	' ,'			0348	+3300450378	
26	X		ST3	AMT(4,5)			0349	+3200230327	
27	X		ZA3	COMAREA.A(10,11)			0350	+3300010373	
28	X		ST3	AMT(6,7)	00006		0351	+3200450327	
29	X		ZA3	COMAREA.A(12,13)			0352	+3300230373	
30	X		ST3	AMT(8,9)			0353	+3200670327	
31	X		ZA3	COMAREA.A(14,15)			0354	+3300450373	
32	X		ST3	AMT(10,11)			0355	+3200890327	
33	X		ZA3	' ,'	00007		0356	+3300010378	
34	X		ST3	AMT(12,13)			0357	+3200010328	
35	X		ZA3	COMAREA.A(16,17)			0358	+3300670373	
36	X		ST3	AMT(14,15)			0359	+3200230328	
37	X		ZA3	COMAREA.A(18,19)			0360	+3300890373	
38	X		ST3	AMT(16,17)	00008		0361	+3200450328	
39	734	*							
40	735	*		THE FOLLOWING IS GENERATED OUT OF LINE					
41	736	*							
42	X	M*6	ZA3	' CR'			0362	+3300090376	
43	X		B	M*3			0363	+0100090335	
44	X	EDMOV02.A	DRDW	-COMAREA.A,COMAREA.A+1			0364	-0003720373	
45	X	M*4	B	M*5			0365	+0100090345	
46	X		B	M*5+4	00009		0366	+0100090349	
47	X		B	M*5+6			0367	+0100090351	
48	X		B	M*5+8			0368	+0100090353	

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	X		B	M*5+10			0369	+0100090355	
02	X		B	M*5+10			0370	+0100090355	
03	X		B	M*5+10	00010		0371	+0100090355	
04	X	COMAREA.A	DA					+0003720374	
				LITERALS					
05	X		'	'	00011	09	0375	'0000000000	0375
06	X		'	CR'		09	0376	'0000006379	0376
07	X		'	DR'		09	0377	'0000006479	0377
08	X		'	'		01	0378	'15	0378
09	X		'S'	'		23	0378	' 25	0378
10	X		' ,'	'		45	0378	' 35	0378

EDMOV Example 3

An automatic-decimal number is edited to the print format with a floating dollar sign and debit and credit indication.

## MOVE—Move Data

MOVE generates instructions that will transmit data from one specified field or area in storage to another.

### Source Program Format

The basic formats for the MOVE statement in the source program are as follows:

Line	Label	Operation	OPERAND			Basic Autocoder			
			25	30	35	40	45	50	55
01	ANY LABEL	MOVE	THISFIELD	TO	THATFIELD				
02	ANY LABEL	MOVE	FIELD A, FIELD B, etc.	TO	THATFIELD				
03	ANY LABEL	MOVE	THISFIELD	TO	FIELD D, FIELD E, etc.				
04	ANY LABEL	MOVE	S	TO	T AND U, V, etc.	TO	W AND X	TO	Y, Z, etc.
05									

In these examples, ANY LABEL is any symbolic label; it may be omitted. The entry MOVE must appear exactly as written. THISFIELD, FIELD A, FIELD B, S, U, V, and X are either the symbolic names of the fields or areas to be moved or alphameric or numerical literals. Numerical literals must be signed. TO and AND are operators that must be written exactly as shown, preceded and followed by a blank. THATFIELD, FIELD D, FIELD E, T, W, Y, and Z are storage locations to which the data is to be moved; the addresses must be symbolic names of fields or areas.

If there are several "from" fields (as in the second format) or several "to" fields (as in the third format), they must be separated by commas. It is not possible to move multiple "from" fields to multiple "to" fields; an attempt to do so will result in an error condition.

Data may bridge words and start at any position in a word, both in the "from" fields and in the "to" fields. Symbolically referenced fields may be any length; literals are restricted as indicated under "Limitations on Length." Data characteristics do not affect the transmission. Data will always be left-justified in the field(s) to which they are moved. The sign of the last item stored in any location determines the sign of the entire word.

In the first format, if THISFIELD is larger than THATFIELD, movement of data will be terminated when THATFIELD is filled. If THISFIELD is smaller than THATFIELD, the data from THISFIELD will be left-justified in THATFIELD and the remaining portion of THATFIELD filled with zeros. The sign of these zero words will be the same as that of the last word moved.

In the second format, the data in FIELD A will be moved to THATFIELD and left-justified. Data from FIELD B will be entered beginning with the digit position following the one in which data from FIELD A terminated. The movement of data continues in the same fashion until the contents of all the specified "from" fields have been moved or until THATFIELD is filled. If all the "from" fields have been moved before THATFIELD is filled, the remainder will be filled with zeros. The sign of the zero words will be the same as that of the last word moved.

In the third format, the data in THISFIELD will be moved to FIELD D and to each subsequent field until all such fields are filled or THISFIELD has been completely

transferred. If the data in THISFIELD is accommodated in the "to" fields without filling them, the remainder of the field(s) will be zeroed out; the sign of the zero words will be the same as that of the last word moved.

In the fourth format, several MOVE operations are performed. The MOVE operations are linked by the operator AND, as indicated. Any of the above three formats for the MOVE statement may be used.

## Processing Techniques

### Limitations on Length

The operand portion of the MOVE macro-instruction may contain 75 parameters. The operators TO and AND are counted as parameters. No limitation is placed on the size of the field if it is referenced symbolically. Literals are restricted as follows: alphameric literals, 120 characters; automatic-decimal literals, 20 digits.

### Address Modification

All symbolic addresses may be modified by indexing and address adjustment.

### The Effect of MOVE

The MOVE macro-instruction is non-destructive in that it does not alter the contents of the "from" field(s). A possible exception might be the case in which the "to" field(s) begin to overlay the "from" field(s), in which case a strict left-to-right procedure of data movement would be maintained.

The following examples illustrate the effect of MOVE on various fields.

No.	"From" field(s)	"To" field(s) before MOVE	"To" field(s) after MOVE
1	+1234567890+12345	-6857463590-7948375403	+1234567890+1234500000
2	+1111111111+22-3333	-9999999999+9999999999	+1111111111-2233330000
3	+1111111111	-99999-9999-9999	+11111+1111+1000
4	-9876543210	+55555	-98765
5	+8888888888+8888	+0000000000-00	+8888888888+88
6	@6162636465	+123456789	@616263646
7	@717273	+1234567	@7172730
8	@818283	+123+123	@818@283

Examples 1, 4, 6, and 7 illustrate a single "from" field and a single "to" field. Examples 2 and 5 illustrate multiple "from" fields. Examples 3 and 8 illustrate multiple "to" fields. In examples 6, 7, and 8, machine difficulty may arise when an attempt is made to print out the "to" fields since invalid double-digit combinations were created at the end of each field by the MOVE.

Although the MOVE macro-instruction may refer to the label(s) of any field or area, it is most frequently used in conjunction with input/output macro-instructions. For example, records (other than Form 3 or Form 4 records as described in the bulletin "IBM 7070 Input/Output Control System") that have undergone preliminary processing in the input area may be moved to a work area by means of a MOVE macro-instruction that references the labels of the DA header lines of the input area and the work area. As another example, a print line might be included in a tape output file by following a PUT macro-instruction by a MOVE macro-instruction that references the label of the DLINE header line and the label of the DA header line of the output area.

When MOVE references the label of a declarative statement other than DA or DRDW, coding will be generated to cause, as a maximum, the entire area or constant defined to be moved or filled.

When MOVE references the label of a DA header line which does not specify a relative address and implicit indexing, coding will be generated to cause, as a maximum, the first record area defined to be moved or filled.

When `MOVE` references the label of a `DA` header line which specifies a relative address and implicit indexing, coding will be generated to cause, as a maximum, the *current* record area (as determined by the contents of the implicit index word) to be moved or filled.

If `MOVE` references the label of a `DRDW`, coding will be generated to move the `RDW` only, *not* the area it defines.

## Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified:

### MULTIPLE FROM- AND TO-FIELDS

An attempt has been made to move data from several fields to several others, which is a format violation. A `NOP` will be generated.

### NO FIELD SIZE. PARAMETER *xx*

The *xx* of the message is replaced by the number of a parameter for which the record does not indicate its size. Thus no coding can be generated to move its contents or to store data in it. A `NOP` is generated.

### NO FROM-FIELD IN MOVE MACRO LINE

The first parameter in the operand portion is the operator `TO`. A `NOP` is generated.

### NO TO-FIELD IN MOVE MACRO LINE

No parameter follows the operator `TO`. A `NOP` is generated.

### TO-FIELD(S) SMALLER THAN FROM-FIELD(S)

The field(s) to which data is to be moved cannot accommodate all of the data of the "from" field(s). Transmission terminates when the "to" area is filled.

### UNACCEPTABLE PARAMETER *xx*

The *xx* of the message will be replaced by the number of a parameter that is not of the types listed under "Source Program Format" as valid. A `NOP` is generated.

## Examples

The following are examples of acceptable coding for the `MOVE` macro-instruction. For each, the associated source-program entries are given, followed by the `MOVE` statement, coding generated in-line, and (where applicable) coding generated out-of-line.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	740	*		MOVE EXAMPLE 1					
02	741		DA	1				+0003250325	
03	742	HERE		00,09A		09	0325		0325
04	743		DA	1				+0003260326	
05	744	THERE		00,09A		09	0326		0326
06	745	*							
07	746	ANYLABEL	MOVE	HERE TO THERE					
08		X ANYLABEL	ZA2	HERE(0,9)	00001		0327	+2300090325	
09		X	ST2	THERE(0,9)			0328	+2200090326	
10	747	*							

### MOVE Example 1

A ten-digit, automatic-decimal number is moved from one location to another location of the same size.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	751	*		MOVE EXAMPLE 2					
02	752		DA	1				+0003250339	
03	753	HERE		00,149'		09	0325		0325
04	754		DA	1				+0003400366	
05	755	AA		00,49'		09	0340		0340
06	756	BB		80,139'		09	0348		0348
07	757	CC		230,269'		09	0363		0363
08	758	*							
09	759	ANYLABEL	MOVE	HERE TO AA,BB,CC					
10		X ANYLABEL	XZA	MACREG.1,HERE	00001		0367	+4600010325	
11		X	RS	MACREG.1,M.1			0368	+6500010369	
12	760	*							
13	761	*		THE FOLLOWING IS GENERATED OUT OF LINE					
14	762	*							
15		X M.1	DRDW	+AA,AA+4			0369	+0003400344	
16		X M.3	DRDW	+BB,BB+5			0370	+0003480353	
17		X M.5	DRDW	-CC,CC+3			0371	-0003630366	

### MOVE Example 2

A 75-character (150-digit) alphameric field is moved to three fields, defined under a single DA, whose total storage area is also equal to 150 digits.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.										PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF			
01	766	*		MOVE EXAMPLE 3								
02	767	FIELDS	DA	1								
03	768	EE		00,07A				+0003250326				
04	769	FF		08,15A			07	0325		0325		
05	770	*					89	0325		0325		
06	771	ANYLABEL	MOVE	FF TO EE								
07	X	ANYLABEL	ZA2	FF(0,1)	00001							
08	X		SL	6			0327	+2300890325				
09	X		A2	FF(2,7)			0328	-5000000206				
10	X		ST2	EE(0,7)			0329	+2400050326				
11	772	*					0330	+2200070325				

MOVE Example 3

An eight-digit field that bridges two locations is moved to a field that does not bridge locations.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.										PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF			
01	776	*		MOVE EXAMPLE 4								
02	777	FIELDS	DA	1				+0003250347				
03	778	AA		00,19A			09	0325		0325		
04	779	BB		20,99A			09	0327		0327		
05	780	CC		100,229A			09	0335		0335		
06	781		DA	1				+0003480371				
07	782	THERE		00,239A			09	0348		0348		
08	783	*										
09	784	ANYLABEL	MOVE	AA,BB,CC TO THERE								
10	X	ANYLABEL	XZA	MACREG.1,THERE	00001		0372	+4600010348				
11	X		RG	MACREG.1,M.1			0373	-6500010377				
12	X		ZA2	THERE(229,229)			0374	+2300990370				
13	X		SL	20			0375	-5000000220				
14	X		ST2	THERE(230,239)			0376	+2200090371				
15	785	*										
16	786	*		THE FOLLOWING IS GENERATED OUT OF LINE								
17	787	*										
18	X	M.1	DRDW	+AA,AA+1	00002		0377	+0003250326				
19	X	M.3	DRDW	+BB,BB+7			0378	+0003270334				
20	X	M.5	DRDW	-CC,CC+12			0379	-0003350347				

MOVE Example 4

Three automatic-decimal fields are moved to a single, larger field.

The unoccupied portion of the "to" field is filled with zeros.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.							PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	801	*		MOVE EXAMPLE 5					
02	802	HERE	DA	1					
03	803			00,149'				+0003250339	
04	804		DA	1			09 0325		0325
05	805	THEREA		00,49'				+0003400344	
06	806		DA	1			09 0340		0340
07	807	THEREB		00,75'				+0003450352	
08	808		DA	1			09 0345		0345
09	809	THEREC		00,23'				+0003530355	
10	810	*					09 0353		0353
11	811	ANYLABEL	MOVE	HERE TO THEREA,THEREB,THEREC					
12		X ANYLABEL	XZA	MACREG.1,HERE	00001		0356	+4600010325	
13		X	RS	MACREG.1,M.1			0357	+6500010368	
14		X	ZA2	HERE(120,125)			0358	+2300050337	
15		X	ST2	THEREB(70,75)			0359	+2200050352	
16		X	ZA1	HERE(126,129)			0360	+1300690337	
17		X	ZA2	HERE(130,139)	00002		0361	+2300090338	
18		X	SL	6			0362	-5000000206	
19		X	A2	HERE(140,145)			0363	+2400050339	
20		X	ST2	THEREC(10,19)			0364	+2200090354	
21		X	ST1	THEREC(0,9)			0365	+1200090353	
22		X	ZA2	HERE(146,149)	00003		0366	+2300690339	
23		X	ST2	THEREC(20,23)			0367	+2200030355	
24	812	*							
25	813	*		THE FOLLOWING IS GENERATED OUT OF LINE					
26	814	*							
27		X M.1	DRDW	+THEREA,THEREA+4			0368	+0003400344	
28		X M.3	DRDW	-THEREB,THEREB+6			0369	-0003450351	

### MOVE Example 5

A 75-character (150-digit) alphameric field is moved into three fields.

The fields are each defined by subsequent entries of separate DAs and

their total storage area is equal to 150 digits. However, the RDWs

generated reserve an area of 160 digits. The possibility of invalid

alpha combinations exists in digit positions 76, 79 of THEREB and in

digit positions 24, 29 of THEREC since these segments are not affected

by MOVE.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF	
01 818	*		MOVE EXAMPLE 6					
02 819		DA	1			+0003250326		
03 820	HEREA		00,19A	09	0325		0325	
04 821		DA	1			+0003270329		
05 822	HEREB		00,29A	09	0327		0327	
06 823		DA	1			+0003300332		
07 824	HEREC		00,24A	09	0330		0330	
08 825	THERE	DA	1			+0003330340		
09 826			74	44	0340		0340	
10 827	*							
11 828	ANYLABEL	MOVE	HEREA,HEREB,HEREC TO THERE					
12	X ANYLABEL	XZA	MACREG.1,THERE	00001	0341	+4600010333		
13	X	RG	MACREG.1,M.1		0342	-6500010345		
14	X	ZA2	HEREC(20,24)		0343	+2300040332		
15	X	ZST2	THERE(70,74)		0344	-2100040340		
16 829	*							
17 830	*		THE FOLLOWING IS GENERATED OUT OF LINE					
18 831	*							
19	X M.1	DRDW	+HEREA,HEREA+1		0345	+0003250326		
20	X M.3	DRDW	+HEREB,HEREB+2	00002	0346	+0003270329		
21	X M.5	DRDW	-HEREC,HEREC+1		0347	-0003300331		

MOVE Example 6

Three automatic-decimal fields, each defined by a separate DA and with a total storage area of 75 digits, are moved to an area of equal size but with unspecified characteristics.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.						PAGE AA
LN CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF	
01 835	*		MOVE EXAMPLE 7					
02 836		DA	5,,0+INDEXWORD			+0003250329		
03 837	HERE		00,09A	09	0325		0000	
04 838		DA	1			+0003300330		
05 839	THERE		00,09A	09	0330		0330	
06 840	*							
07 841	ANYLABEL	MOVE	HERE TO THERE					
08	X ANYLABEL	ZA2	HERE(0,9)+INDEXWORD	00001	0331	+2301090000		
09	X	ST2	THERE(0,9)		0332	+2200090330		
10 842	*							

MOVE Example 7

The MOVE statement references the label of a subsequent entry under a DA that contains a relative address and implicit indexing. The contents of the HERE field of the current record (as determined by the contents of the implicit index word) will be moved to the THERE field.



PAGE	AA	PROGRAM	LN	LABEL	OP	OPERAND	LOC	INSTRUCTION	REF.ADR	PGLIN	CD	NO
			01	*		MOVE EXAMPLE 8						846
			02		DA	2,,0+INDEXWORD		+0003250326				847
			03	HERE		00,09A	0325		09 0000			848
			04		DA	2,,0+X15		+0003270330				849
			05	THERE		00,12A	0327		09 0000			850
			06	*								851
			07	ANYLABEL	MOVE	HERE TO THERE						852
			08	ANYLABEL	ZA2	HERE(0,9)+INDEXWORD	0331	+2301090000				GENRD 00001
			09		ST2	THERE(0,9)+X15	0332	+2215090000				GENRD
			10		SL	20	0333	-5000000220				GENRD
			11		ST2	THERE(10,12)+X15	0334	+2215020001				GENRD

## MOVE Example 8

The ten-digit HERE field of the current record is moved to the larger THERE field of the current record. The unoccupied portion of the THERE field will be filled with zeros and the sign will be made the same as the sign of the last word moved.

PAGE	AA	PROGRAM	LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
			01	8501	*		MOVE EXAMPLE 9					
			02	8502	HERE	DA	1				+0003250328	
			03	8503	AA		00,15A		09	0325		0325
			04	8504	BB		20,25'		05	0327		0327
			05	8505	CC		26,32		69	0327		0327
			06	8506	THERE	DA	1				+0003290333	
			07	8507			42		22	0333		0333
			08	8508	*							
			09	8509	ANYLABEL	MOVE	HERE TO THERE					
			10	X	ANYLABEL	XZA	MACREG.1;HERE	00001		0334	+4600010325	
			11	X		RS	MACREG.1;M.1			0335	+6500010339	
			12	X		ZA2	THERE(39,39)			0336	+2300990332	
			13	X		SL	20			0337	-5000000220	
			14	X		ST2	THERE(40,49)			0338	+2200090333	
			15	8510	*							
			16	8511	*		THE FOLLOWING IS GENERATED OUT OF LINE					
			17	8512	*							
			18	X	M.1	DRDW	-THERE;THERE+3	00002		0339	-0003290332	

## MOVE Example 9

The entire HERE field through digit position 39 is moved to the larger THERE field. The unoccupied portion of the THERE field (40,49) will be filled with zeros and the sign will be made the same as the sign of the last word moved. In this example, the possibility of creating invalid alpha combinations does not exist since the sign of the word to which field BB and part of field CC are to be moved is determined by the sign of CC. Since the format of field CC is unspecified, the processor assumes the field contains numerical data to be treated as a signed integer.

## SHIFT—Shift and Store

SHIFT generates instructions to place the contents of a field into one or more accumulators, to shift the data in a specified way, and to store the result.

### Source Program Format

The basic format for the SHIFT statement in the source program is as follows:

Line	Label	Operation	OPERAND					Basic Au			
3	56	15	16	20	21	25	30	35	40	45	50
0,1	ANY LABEL,	SHIFT	OPTION,	START,	COUNT,	FIELD A,	FIELD B				
0,2											

ANY LABEL is any symbolic label; it may be omitted. The entry SHIFT must be written exactly as shown. OPTION is one of the following one- or two-letter codes, specifying the type of shift to be executed:

OPTION	Type of Shift
L	Left
R	Right
LC	Left and Count
RR	Right and Round
LS	Left Split
RS	Right Split

START is an integer indicating where to begin shifting in case of a split-shift option; it is determined by counting from the left-most digit of the field, beginning with 1. (With other options, START will be blank, but the separating commas must be entered.) With one exception, COUNT is the actual number of positions to be shifted, restricted only by the size of the field. In the case of a shift-left-and-count option, however, COUNT must be an index word, referenced either by its actual, two-digit number, without a preceding x, or by its symbolic name, which will contain in digit positions 4-5 the number of high-order zeros found in the shifted field. FIELD A is the symbolic address of the field to be shifted; it must be less than twenty-one digits in length. FIELD B is the field in which the result is to be stored; it may also be referenced by its symbolic address. The sign of the field to be shifted is transmitted with the field and stored with the result in FIELD B. The basic format may be modified by the omission of FIELD B. In that case, the data will be restored in its original field after shifting.

No warning is issued either in case of possible digit loss when the result of a shift is larger than FIELD B or when invalid alpha combinations have been created.

### Processing Techniques

#### Limitations on Length

The number of parameters is fixed by the format. The field to be shifted is limited to twenty-digit length.

### **Address Modification**

All symbolic addresses may be modified by indexing and address adjustment.

### **Error and Warning Messages**

The following error and warning messages will be issued during assembly under the conditions specified:

#### **COUNT GREATER THAN FIELD-SIZE**

The **COUNT** field contains an integer larger than the number of digits in the field to be shifted.

#### **COUNT IS ZERO**

No shift will be carried out. The field will be stored without shifting.

#### **COUNT NOT AN INDEXWORD**

In a shift-left-and-count option, a label has been entered for **COUNT** that is defined elsewhere as other than an index word. A **NOP** will be generated.

#### **FIELD GREATER THAN 20 DIGITS**

An error condition, since fields to be shifted may not exceed this size. A **NOP** will be generated.

#### **FIELD UNACCEPTABLE**

An attempt has been made to shift an alteration switch or some other entity that is not a part of storage. A **NOP** will be generated.

#### **INCORRECT NUMBER OF PARAMETERS**

An attempt has been made to write more than five parameters into the operand portion of a **SHIFT** statement. Parameters in excess of the first five will be ignored.

#### **INCORRECT OPTION**

The option code is blank or is not one of those listed under "Source Program Format."

#### **SHIFTING INSTRUCTION**

The field to be shifted is an imperative instruction. Coding to accomplish this will be generated nevertheless.

#### **START GREATER THAN FIELD-SIZE**

The **START** counter in a split-shift option contains an integer larger than the number of digits in the field to be shifted.

### **Examples**

The following are examples of acceptable coding for the **SHIFT** macro-instruction. For each, the associated source-program entries are given, followed by the **SHIFT** statement, coding generated in-line, and (where applicable) coding generated out-of-line.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.								PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	856	*		SHIFT EXAMPLE 1						
02	857		DA	1				+0003250325		
03	858	FIELDA		00,09		09	0325		0325	
04	859	*								
05	860	ANYLABEL	SHIFT	L,,3,FIELDA						
06		X ANYLABEL	ZA2	FIELDA(0,9)	00001		0326	+2300090325		
07		X	SL2	3			0327	+5000002203		
08		X	ST2	FIELDA(0,9)			0328	+2200090325		
09	861	*								

SHIFT Example 1

The contents of FIELDA are shifted left three positions and the result is stored again in FIELDA.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.								PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	865	*		SHIFT EXAMPLE 2						
02	866	COUNT	EQU	90,X						
03	867		DA	1				+0003250325		
04	868	FIELDA		00,09		09	0325		0325	
05	869		DA	1				+0003260326		
06	870	FIELDDB		00,09		09	0326		0326	
07	871	*								
08	872	ANYLABEL	SHIFT	LC,,COUNT,FIELDA,FIELDDB						
09		X ANYLABEL	ZA2	FIELDA(0,9)	00001		0327	+2300090325		
10		X	SLC2	COUNT			0328	+5000902300		
11		X	ST2	FIELDDB(0,9)			0329	+2200090326		
12	873	*								

SHIFT Example 2

The digits in FIELDA are shifted to the left until a digit other than a zero is in the high-order position. If the high-order digit is non-zero to start with, no shift takes place. The number of positions shifted is recorded in the index word labeled COUNT.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.								PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	874	*		SHIFT EXAMPLE 3						
02	875		DA	1				+0003250328		
03	876	FIELD A		00,14						
04	877	FIELD B		20,34		09	0325		0325	
05	878	*				09	0327		0327	
06	879	ANYLABEL	SHIFT	RR,5,FIELD A,FIELD B						
07		X ANYLABEL	ZA1	+0	00001		0329	+1300000337		
08		X	ZA2	FIELD A(0,9)			0330	+2300090325		
09		X	SL	5			0331	-5000000205		
10		X	A2	FIELD A(10,14)			0332	+2400040326		
11		X	SRR	5			0333	-5000000105		
12		X	ST2	FIELD B(10,14)	00002		0334	+2200040328		
13		X	SR	5			0335	-5000000005		
14		X	ST2	FIELD B(0,9)			0336	+2200090327		
15	880	*								
16	881	*		THE FOLLOWING IS GENERATED OUT OF LINE						
17	882	*		LITERALS						
18		X		+0			00 0337	+0	0337	

## SHIFT Example 3

The contents of FIELD A are shifted right 5 positions and the amount shifted is rounded off. The result is stored in FIELD B.

PAGE AA	PROGRAM	7070 COMPILER SYSTEM VERSION OMY08, CHANGE LEVEL 00001.								PAGE AA
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF	
01	885	*		SHIFT EXAMPLE 4						
02	886		DA	1				+0003250328		
03	887	FIELD A		00,19						
04	888	FIELD B		20,39		09	0325		0325	
05	889	*				09	0327		0327	
06	890	ANYLABEL	SHIFT	RS,12,8,FIELD A,FIELD B						
07		X ANYLABEL	ZA1	FIELD A(0,9)	00001		0329	+1300090325		
08		X	ZA2	FIELD A(10,19)			0330	+2300090326		
09		X	SRS	8(11)			0331	-5000001608		
10		X	ST2	FIELD B(10,19)			0332	+2200090328		
11		X	ST1	FIELD B(0,9)			0333	+1200090327		
12	891	*								

## SHIFT Example 4

Starting with the 12th digit, the contents of FIELD A are shifted 8 positions to the right. The result is stored in FIELD B.

## SNAP—Memory Print-Out

SNAP generates instructions to provide a listing of a specified portion of storage.

### Source Program Format

The basic format for the SNAP statement in the source program is:

Line	Label	Operation	OPERAND												Basic Autocoder	Autocoder
3	5/6	15/16	20/21	25	30	35	40	45	50	55	60	65	70	75		
0.1	ANY LABEL	SNAP	FROMADDR	TOADDR	OUTPUTUNIT	SKIPSPRIOR	SNAPS	SKIPSAFTER								
0.2																

ANYLABEL is any symbolic label; it may be omitted. The entry SNAP must be written exactly as shown. FROMADDR and TOADDR represent the limits of storage to be listed and may be symbolic or actual addresses. OUTPUTUNIT is the symbolic address of any tape channel and unit, punch, printer, or typewriter that has been defined by an EQU statement elsewhere in the program. SKIPSPRIOR, SNAPS, and SKIPSAFTER are unsigned, one- or two-digit integer counters. SKIPSPRIOR controls the number of times the program will pass through the location of the SNAP statement without taking a print-out. SNAPS controls the number of times following this that print-outs will be taken. SKIPSAFTER controls the number of times that print-outs will again be omitted. Thereafter, control alternates between SNAPS and SKIPSAFTER for the duration of the program.

The basic format may be modified by the omission of one or more of the three counters. If omitted, SKIPSPRIOR will automatically be set equal to 0, SNAPS to 1, and SKIPSAFTER will cause permanent skipping once the specified number of SNAPS has been taken. If either SKIPSPRIOR or SNAPS is omitted, or both, separating commas must be punched.

### Processing Techniques

#### Limitations on Length

The number of parameters is fixed by the basic format, subject only to the omission of one or more counters. There is no limit to the amount of storage to be printed out, although approximately 150 locations of storage will be required by the calling sequence and the subroutine that constitute SNAP. Conceivably these may be positioned so as to fall between FROMADDR and TOADDR, in which case this area of storage will be shown as required by the SNAP instruction. The programmer has control over the placement of the subroutine, however, by means of the LITORIGIN statement.

#### Address Modification

All addresses may be modified by indexing; symbolic addresses may be modified by address adjustment.

#### The Effect of SNAP

The print-out provided by SNAP is non-destructive and will consist of the electronic switches, the accumulators, index words 10 and 11, and storage between the limits specified. (In the case of a type-out, only storage will be given.) The stor-

age area specified must be contiguous. Index words 10 and 11 are used by SNAP; if they fall between FROMADDR and TOADDR, their representation in that section of the print-out will be that required by SNAP, not that given by the source program. For this reason, the source-program contents of these words are printed out ahead of the storage area. Before the exit from the SNAP subroutine, both index words are restored to their original contents.

The output will be in the following format:

First line:                SNAP Electronic Switches, Instruction Counter at time of SNAP branch.

Second line:             Accumulators, Index Words 10 and 11.

Each succeeding line:    Five words of storage each, with the location of the first word printed on each line shown for identification purposes. Numerical words carry the proper sign. Alpha words are printed with an A in the sign position and contents in double-digit representation; this is done to avoid printer difficulty which might arise if an alpha location contained a double-digit combination that was not the code for any character.

When tape output is specified, tape density must have been set previously by the program or the operator. No provision for tape labels, end-of-file conditions, or tape mark and rewind routines has been made in order to conserve space, and because it will often be desirable to use the same output tape for SNAP print-outs as that used for utility program output.

When the SNAP routine is entered, the priority mask is set to "prohibit." On leaving the subroutine, the mask is set to "allow," regardless of its prior condition. If the programmer does not want the mask at "allow," he must restore it after each use of SNAP.

## Error and Warning Messages

The following error and warning messages will be issued during assembly under the conditions specified:

### BLANK PARAMETER *xx*

The *xx* of the message will be replaced by 01, 02, or 03, depending upon whether FROMADDR, TOADDR, or OUTPUTUNIT has been left blank. A NOP will be generated.

### INVALID PARAMETER *xx*

The *xx* will be replaced by 01 or 02 if FROMADDR or TOADDR, respectively, is not the label of a location of storage; by 03 if OUTPUTUNIT is not the label of one of the acceptable units listed under "Source Program Format." A NOP will be generated.

### SNAPS COUNTER ZERO

The SNAPS counter is zero, rendering the SNAP statement ineffective. A NOP will be generated; this will at once be an aid to patching and a literal implementation of the source statement as written.

### TOO MANY PARAMETERS. WILL IGNORE

More than six entries have been written into the operand portion of a SNAP statement. The *first* six will be accepted and coding generated accordingly.





LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	891	*		SNAP EXAMPLE 1					
02	892	SNAPOUTP	EQU	19, CU					
03	893	*							
04	894	ANYLABEL	SNAP	0,99,SNAPOUTP,1,2,3					
05	X	ANYLABEL	BLX	94,SNAP001.A	00001		0325	+0200940357	
06	X		DRDW	+0,99			0326	+0000000099	
07	X		PIW	SNAPOUTP,SNAP018.A			0327	+8100930356	
08	X		DC					+0003280328	
09	X			+0100020003			09	0328	+0100020003
10	895	*							0328
11	896	*		THE FOLLOWING IS GENERATED OUT OF LINE					
12	897	*							
13	X	SNAP002.A	DA	1				+0003290333	
14	X			49,49			99	0333	
15	X		DA	1				+0003340349	
16	X	SNAP015.A		00,159			09	0334	
17	X	SNAP016.A	DRDW	+SNAP015.A+2,SNAP015.A+3	00002		0350	+0003360337	0334
18	X		DRDW	+SNAP015.A+5,SNAP015.A+6			0351	+0003390340	
19	X		DRDW	+SNAP015.A+8,SNAP015.A+9			0352	+0003420343	
20	X		DRDW	+SNAP015.A+11,SNAP015.A+12			0353	+0003450346	
21	X		DRDW	-SNAP015.A+14,SNAP015.A+15			0354	-0003480349	
22	X	SNAP017.A	DRDW	-SNAP015.A+2,SNAP015.A+13	00003		0355	-0003360347	
23	X	SNAP018.A	DRDW	-SNAP015.A,SNAP015.A+15			0356	-0003340349	
24	X	SNAP001.A	PC	+1111111111			0357	+5500000470	
25	X		ZST1	SNAP002.A			0358	-1100090329	
<hr/>									
24	X	SNAP020.A	ZA1	2(8,9)+X94 LOAD SKIP CTR			0453	+1394890002	
25	X		STD1	2(0,1)+X94			0454	-1294010002	
26	X		ZA1	2(4,5)+X94 LOAD SNAP CTR	00023		0455	+1394450002	
27	X		BZ1	SNAP021.A+2 MAKE NOP IF SNAPCTR ZERO			0456	+1000090460	
28	X		B	SNAP001.A+4			0457	+0100090361	
29	X	SNAP021.A	ZA1	2(7,7)+X94 SW7 TEST FOR SKIP			0458	+1394770002	
30	X		BZ1	SNAP020.A SETUP SKIPSAFTER CTR			0459	+1000090453	
31	X		ZA1	-0100000000 SETUP SKIPSAFTER FOREVER	00024		0460	+1300090471	
32	X		XS	94,1			0461	-4700940001	
33	X		ZST1	0+X94			0462	-1194090000	
34	X		ZST1	1+X94			0463	-1194090001	
35	X		ZST1	2+X94			0464	-1194090002	
36	X		ZST1	3+X94			00025	0465	-1194090003
37	X	SNAP022.A	B	SNAP009.A			0466	+0100090440	
<hr/>									
38	X			LITERALS					
38	X			+0000000000			09	0467	+0000000000
39	X			+0000000006			09	0468	+0000000006
40	X			+0000000012			09	0469	+0000000012
41	X			+1111111111	00026		09	0470	+1111111111
42	X			-0100000000			09	0471	-0100000000
43	X			'SNAP '	00027		09	0472	'8275617700
44	X			'+'			01	0473	'20
45	X			'-'			23	0473	'30
46	X			'A'			45	0473	'61

SNAP Example 1

A listing of the index word area of storage is desired on tape for two consecutive passes of the program through this area following the first pass through. Tape listings are then omitted for the next three passes. Control then alternates between producing the two tape listings and causing the three skips for the duration of the program.

# The Processor

## The Organization of the Processor

The Autocoder processor performs functions such as assembly, compilation, and generation for the Autocoder portion of the IBM 7070/7074 Compiler Systems Tape. In addition to Autocoder, the Compiler Systems Tape contains the compilers for FORTRAN, the Report Program Generator, and Commercial Translator. The runs which perform these various functions for all of these compilers fall into two general categories: Compile Runs and Generator Runs. These runs are described in the 7070/7074 Data Processing System Bulletin "IBM 7070/7074 Compiler Systems: Operating Procedure," form J28-6105.

Compilation of a source-language program is performed by a *Compile Run*. The program is converted to machine language and shown on the output listing along with the original symbolic instructions, and then punched into a condensed program deck. A *Generator Run* compiles input statements which create a new macro generator and, if desired, enters the resultant program on a new systems tape. In addition, changes to the Compiler Systems Tape are processed during a *Systems Run* and an updated systems tape is produced.

The operation which is used most frequently in Autocoder processing is that of compilation; this operation is performed by a Compile Run, as well as the compilation portion of a Generator Run. The function of the processor during compilation is discussed on the following pages.

The Autocoder processor consists of the following major sections:

*Systems Control:* This section of the processor exercises overall control over the compilation process and is principally charged with "housekeeping" functions.

*Phase I:* This is the first phase in the conversion of the source program to an object program. Phase I scans the input statements of the source program and creates records which will constitute input to Phases II and III.

*Phase II:* This is the second phase in the conversion of the source program to an object program; Phase II will not become active until Phase I has completed its role in processing. Phase II, in conjunction with the macro generators, compiles the macro-instructions of a program.

*Phase III:* This is the third and last phase in the conversion of the source program to an object program; Phase III will not become active until Phase II has completed its role in processing. Phase III assembles the final program on the basis of records received from Phases I and II. The output from Phase III consists of a machine-language program and a program listing.

*Autosort:* Autosort is the portion of the processor which contains appropriate sort routines which are called when it is necessary for the processor to reorder its records for a pass.

*Macro Generators:* The macro generators and function subroutines are contained on the Compiler Systems Tape and are made available when required to compile macro-instructions in Phase II.

The functions of Systems Control and Phases I, II, and III during compilation and assembly will now be described in more detail.

## **Systems Control**

The chief responsibilities of Systems Control during compilation are as follows:

1. To determine the type of run to be made.
2. To check the validity of the operating options.
3. To supervise the assignment and readiness of the various tape units.
4. To locate and load the various coding blocks on the systems tape.
5. To turn control over to the coding blocks at appropriate times to effect actual processing.

In addition to these functions, Systems Control incorporates the desired modifications, additions, or deletions to the system on a new Compiler Systems Tape when required by a Systems Run or a Generator Run.

The Systems Control program consists of ten sections, called Communication Record, Systems Control 1-8 (SYCL1, SYCL2, etc.), and Update.

### **Communication Record**

The Communication Record consists entirely of data; it contains no instructions to be executed. The data includes the characteristics of the processor machine, the object machine, and the object program, as specified by the programmer through the use of Operating Option Control Cards. The Communication Record also includes various codes generated within the processor for communication between coding blocks.

### **Systems Control 1 (SYCL1)**

SYCL1 loads the Communication Record and turns control over to SYCL3. This section of Systems Control also contains the input/output routines and the Systems Tape Control; both of these remain in storage throughout the compilation process. The Systems Tape Control locates, loads, and transfers control to the various coding blocks.

### **Systems Control 2 (SYCL2)**

SYCL2 is generally called only at the end of a Compile Run to identify halts and output tapes by messages on the console typewriter. In Multifile Runs, it tests for the presence of another program to be compiled and turns control over to SYCL3 again if indicated. SYCL2 may also be used to discontinue a run when the processor determines that this is necessary.

### **Systems Control 3 (SYCL3)**

SYCL3, the largest and most complex portion of Systems Control, initiates the desired run. It reads the various options into the Communication Record, does some validity checking of options, and may type error messages and even cause a halt. It opens the required input and work tapes and does whatever else is necessary in the way of "housekeeping" functions to prepare the type of run specified by the RUN Control Card. If this is an Autocoder Compile Run, Phase I is called.

### **Systems Control 4-8 (SYCL4-SYCL8)**

These sections, which provide for the mounting and dismounting of tapes, are always called by, and in turn will call, a routine other than a Systems Control routine. No two of these routines are identical, but they are quite similar and are included on the systems tape as one program section at five different points. Of these programs, only SYCL6 is loaded during Compile Runs.

### **Update**

Update is a program controlling changes to the systems tape; this program is not loaded during a Compile Run.

## Phase I

Phase I reads the source program from cards or tape. Each input statement is assigned a serial number, for internal use by the processor, to govern program order during assembly. Source program page and line numbers are retained and will reappear in the program listing, but they are not utilized during the assembly process. A program identification record is created from columns 76-80 of the first program card. This identification applies to the entire program to be compiled, and columns 76-80 of all subsequent cards are ignored.

As the input statements are scanned by Phase I, a certain amount of error checking is done, and, where necessary, error messages are issued which will ultimately appear in the program listing. In general, Phase I confines its validity checking to such major format violations as omission of operation codes or operands, malformation of the operand portions of macro-instructions, use of illegal characters, errors in field definition, use of unallowable address types, exceeding machine capacity, or exceeding the maximum permissible number of LITORIGIN statements. If Phase I encounters an input statement which cannot be processed, it will usually retain as much of the statement that was scanned, or it will generate a NOP. An error message will be issued.

The chief function of Phase I is to furnish records to Phases II and III that will enable them to compile the finished program; this is done by analyzing the source program statements and by writing records for the label and operand entries. These "element" records contain pertinent information about the source program entry, such as its serial number, its operation code, and other attributes relevant to the processor. Output from Phase I consists of two files, File A and File B. File A, which contains all of the element records, goes directly to Phase III. File B, which contains all records that may be required for the compilation of macro-instructions, becomes input to Phase II. (Phase II will receive not only records derived from macro-instructions themselves, but also label records from all other instructions having symbolic labels, since these may be referenced by macro-instruction operands.)

Input statements are scanned one at a time. When comments cards are encountered, appropriate records are written onto File A for inclusion in the final program listing. Other source statements are classified according to their operation code, which will be that of a declarative statement, a control statement, an imperative statement (either a symbolic machine instruction or a macro-instruction), or blank. If the operation code does not fall into any of these categories, or if a blank operation code occurs on a card following one that does not permit continuations or subsequent entries, the code is converted to a NOP, and an error message will appear for that line of the program listing.

Phase I processes its input statements in three successive passes, each of which is processed by a separate coding block. These coding blocks contain the programs that will process statements with the following operation codes, respectively:

<i>Pass 1</i> <i>First Coding Block</i>	<i>Pass 2</i> <i>Second Coding Block</i>	<i>Pass 3</i> <i>Third Coding Block</i>
Symbolic machine instructions	DA CODE	XRESERVE Control SRESERVE Control
DRDW	DC	XRELEASE Control
EQU	DLINE	SRELEASE Control
BRANCH Control	DSW	DIF
END Control	ORIGIN Control LITORIGIN Control	DUF Macro-instructions

The first coding block is loaded into storage and the entire source program is scanned for the statements whose programs are contained on this block. These statements are processed at once; the others are written onto a temporary work tape. The second coding block is then loaded into storage. The work tape containing the statements which have not been processed is scanned. The statements whose programs are contained on the second coding block are processed; the remaining statements are again written onto a work tape. Finally, the third coding block is loaded into storage and the remaining statements are processed. Thus, each of the three coding blocks needs to be loaded into storage but once.

This procedure, of course, temporarily destroys source program (i.e., page and line) order, but this order will ultimately be restored by means of the internal serial number assigned to each statement. In addition, input statements are generally taken apart, with separate records being written for the label and the one or more operand entries. These individual components of a given source statement are also numbered and provided with indicators that will permit eventual recomposition of the source statements.

## Phase II

Phase II processes the records on File B in order to perform the compilation of macro-instructions. Its output, a series of records corresponding to symbolic machine instructions and declarative statements, is written on File A.

The first step of Phase II is a sort of File B, which contains all label records of the declarative or imperative statements having symbolic labels; the records are grouped by symbolic name.

The sort is followed by an Information Transfer pass. This pass provides additional necessary information concerning the characteristics of fields occurring in the operands of macro-instructions; this is performed by transferring the characteristics shown in the label records into the macro-instruction operand records.

At the conclusion of the Information Transfer pass, all records other than macro-instruction labels and operands are dropped from File B. The remaining records are sorted and grouped by macro-instruction operation code. Within this grouping, the records are arranged according to their order of appearance in the program. This sorting avoids frequent reloading of the same macro generator since the processor will be able to operate successively on all appearances of each macro-instruction.

The generating portion of Phase II is governed by a program called Phase II Control. This program determines which macro generator is required and calls the appropriate coding block from the library. While one generator is being executed, Systems Tape Control (see "Systems Control 1 (SYCL1)") positions the tape in anticipation of the next required generator. Phase II Control sets up certain counters required by the generators, such as parameter counters and counters to record the number of previously generated labels. In addition, it causes the macro generator required to be loaded into storage, furnishes appropriate "parameter" records to it as a basis for analysis, and then turns control over to the generator itself.

The parameter records are not quite identical to the element records written by Phase I and completed in the Information Transfer pass, but they are based on them. The parameter records need not contain the operation code or the serial number; this information is stripped from the element record and temporarily stored elsewhere, to be attached again to output records. In general, the rest of the element record goes into the parameter record area unchanged, but some items that cannot be conveniently accommodated there, such as long alpha-

meric literals, the input texts of ARITH and LOGIC statements, etc., are stored in a separate area. rdws defining their location are made available to the generator in the parameter record area. Records of address adjustment and indexing are not preserved as independent entities; the pertinent information is entered into the records of the parameters they modify. Since one macro-instruction is compiled at a time, the parameter record area will, at any one time, contain only those records derived from a single macro-instruction.

Generated output is not, strictly speaking, produced by the macro generators themselves, but by the GENER subroutine, which is in storage during the entire generating portion of Phase II. The macro generators contain “model statements” which, together with the parameter records and certain indicators set by the generator, serve as a guide for the construction of appropriate records by GENER. These records, which are written onto File D, correspond to a sequence of symbolic machine instructions or macro-instructions that will accomplish the operations indicated by the original macro statement.

Certain Autocoder macro generators pass on portions of their work to other “lower-level” macro-instructions. For example, the LOGIC macro-instruction often passes work on to the COMP macro-instruction. In such cases, the generated output on File D will include both macro-instructions which require further compilation and symbolic machine instructions which do not. Therefore, after the generating portion of Phase II, the temporary output file, File D, is edited into two further files. Symbolic machine instructions, declarative statements, and control statements are written onto File A as eventual input to Phase III; an editing program changes the format of these records into a format identical to that produced by Phase I. The generated macro-instructions which require further compilation are edited to a format identical to that produced by Information Transfer. Since these records have already been augmented to include data characteristics, the records are re-entered at a point *following* the Information Transfer pass and *beginning* with the sort by macro-instruction operation code.

When no further macro-instructions remain to be compiled, all Phase II output will have been edited and written onto File A; this file then goes to Phase III for assembly.

## Phase III

At the start of Phase III all macro input statements have been reduced to sets of elementary Autocoder statements consisting of machine instructions, declaratives, and several types of comments records. The input to Phase III was written on File A by Phases I and II. After sorting these statements back into the order of the source program, the statements are assigned machine locations, the operands are replaced by their equivalent machine location values, and the individual machine instructions are built. These are shown on the output listing, along with the original symbolic instructions, and then packed into a condensed program deck.

Phase III is organized into six passes and four sorts, as follows:

1. Record Construction
2. Sort Serial-Request File
3. Serial Transfer
4. Sort Statement File
5. Assignment
6. Sort Symbol File

7. Information Transfer
8. Sort Operand File
9. Output
10. Message

The Assignment pass is the central processing section of Phase III. Its primary function is to assign a machine location to each statement as it is encountered. The passes and sorts preceding the Assignment pass prepare the input for this pass. The succeeding passes and sorts handle the final preparation and editing of the output.

### **Record Construction**

The first pass, the Record Construction pass, writes two types of records. On one file, statement records are constructed from the element records written by Phases I and II. A statement record is written for each statement; it contains the label, operation code, one or two operands, field control, address adjustment, indexing, and remarks. On a second file, Record Construction creates the following request records:

1. A Serial Record for each labeled statement (including EQU statements).
2. An Equate Request Record for each symbolic operand of each EQU statement.
3. An Index Assignment Request Record for each symbolic index word appearing either in the operand of an index word operation such as XL or XZA, or as indexing.
4. A Switch Assignment Request Record for each symbolic electronic switch appearing in the operand of an electronic switch operation such as BES or ESN.
5. An Implicit Indexing Request Record for each labeled DA subsequent entry included under a DA header line showing implicit indexing.
6. An Origin Request Record for each symbolic NAMEONE operand of an ORIGIN or LITORIGIN statement.
7. A Literal Request Record for each literal or adcon appearing in the operand of an instruction.

### **Sort Serial-Request File**

The Serial Record and request records created by the Record Construction pass are sorted by symbolic name and, further, by the type code of each request record. This sort places the file in order so that all request records for each symbolic name are grouped together following the Serial Record, if one has been created. The literals and adcons will be sorted into two groups following the request records. These records will be in order by adcon name or by literal value.

### **Serial Transfer**

The Serial Transfer pass selects the request records to be added to the statement file created during the Record Construction pass; this statement file will be processed during the Assignment pass.

Each Index Assignment Request Record and each Switch Assignment Request Record is read; the request record for each index word and electronic switch with the lowest card serial number (i.e., the record of the *first* occurrence of the symbolic name in the program) is passed on to the statement file.

Serial Transfer retains a record of the card serial number of each Index Assignment Request Record it passes to the statement file. This number is transferred to the Implicit Indexing Request Record for each labeled DA subsequent entry. These records are passed on to the statement file.

Equate Request Records have been created in the Record Construction pass for the symbolic operands of EQU statements. In addition, Serial Records have been created for the label of each EQU statement. The Serial Transfer pass matches the Serial Records with their corresponding Equate Request Records. The card serial number from the Serial Record is added to the Equate Request Record; the Equate Request Records are then passed on to the statement file.

Finally, the Serial Transfer pass processes the Literal Request Records. This pass essentially creates a DC for the complete set of literals and adcons. A DC subsequent entry, with either the literal value or the adcon as its operand, is created for each unique literal or adcon; duplicate references are read and dropped.

### **Sort Statement File**

The complete statement file is now sorted by card serial number to give the Assignment pass a file in source program order. The request records have all been given card serial numbers and will be merged with their respective statements. Literals will have been numbered so that they fall last in each litorigin segment of the program.

### **Assignment**

The Assignment pass assigns a machine location to each statement as it is encountered. A statement might be a symbolic machine instruction which occupies one word of storage or it might be a declarative which occupies any amount of storage from part of one word to many words. The Assignment pass assigns a reference address to each statement regardless of the type; the reference address is composed of the address and field control of the part of the statement that occupies the first word of storage assigned to this statement.

In addition to its function of assigning a reference address to each statement, the Assignment pass creates two types of records, Definition Records and Operand Records, which are placed in the Information Transfer file. Definition Records are created for each symbolic label encountered; the record contains the reference address which has been assigned to the labeled statement. Operand Records are created for each symbolic operand encountered.

Symbolic index words and electronic switches create special cases; these symbols must be assigned machine addresses. The statement file sort has placed an Index Assignment Request Record following the instruction containing the first reference to the symbolic index word; likewise, it has placed a Switch Assignment Request Record following the instruction containing the first reference to the symbolic electronic switch. These request records initiate the assignment of a machine location and the creation of a Definition Record to be placed in the Information Transfer file.

Labels of DA subsequent entries whose DA header lines show implicit indexing also create special cases. Implicit Indexing Request Records, which were created by the Record Construction pass, have been placed following the appropriate Index Assignment Request Records. These request records initiate the creation of Definition Records which contain the implicit index value; the Definition Records are placed in the Information Transfer file.

Another special case is equating one symbol to another symbol. An Equate Request Record has been placed following the statement with the same label as the EQU statement operand. The request record initiates the creation of a Definition Record to be placed in the Information Transfer file.

When ORIGIN and LTORIGIN statements are encountered, the Assignment pass must be able to change the address assignment to continue from some previously



defined address. An Origin Request Record for the operand NAMEONE will follow the statement with the same label as the NAMEONE. It will cause an entry in a table containing values of addresses for reference by ORIGIN and LITORIGIN statements.

**Sort Symbol File**

The Information Transfer file is sorted by symbolic name. Within each set of records for one symbol, Definition Records precede Operand Records.

**Information Transfer**

The Information Transfer pass will transfer the assigned value in each Definition Record to each operand that refers to the defined symbol. Simultaneously, the pass will produce the Cross Reference Listing which will be written on the end of the statement file.

**Sort Operand File**

The Operand Records, with assigned values, will be sorted on the page and line number of the statement which contained the symbolic operand; the Operand Records will then be used as input to the Output pass.

**Output**

The Output pass will read each statement record and reconstruct the card image of the statement, i.e., label, operation code, and operand. The Operand Records that pertain to a statement will be read and the values substituted for symbols in the statement. The pass will construct the machine instructions which will be shown with the original symbolic instructions on the output listing; the machine instructions will be condensed into the output program deck. Any halt instructions that occur will be copied onto the Message file, which will also receive any Message Records that have been added to the statement file during the entire process. After the last statement has been processed, the Message pass will be called.

**Message**

This pass will first write the memory map, the index word and electronic switch availability tables, and the Cross Reference Listing from the end of the statement file onto the output listing. Next, a component listing from the end of the Operand Record file will be added to the output listing. The list of halt statements are then written, followed by the list of error messages.

**Output Listings****Program Listing**

The Program Listing is prepared during the Output pass of Phase III. This listing contains the original symbolic instructions with the assembled machine instructions. A sample Program Listing is shown on the opposite page. The heading (1) of the Program Listing is "7070 COMPILER SYSTEM VERSION XXXXX, CHANGE LEVEL YYYYY." Appropriate substitutions for the version number and change level number are determined by the identification on the Compiler Systems Tape. A page number (2) appears twice on each page. The page numbers are assigned to each page of the listing by Phase III; the pages are numbered sequentially by a two-letter symbol (AA, AB, . . . through ZZ). The PROGRAM ZZZZZ (3) entry is the program identification record; if no identification is used, this entry will be blank.

The listing on the left-hand portion of the page contains the original source-program entries as well as the symbolic machine instructions, DRDWS, literals, etc., that have been generated. The line number (4) is assigned by Phase III to each entry. The card reference number (5) is the page and line number



assigned by the programmer to each source-program entry. The label (6), operation (7), and operand (8) columns contain the source-program entries and the generated statements. Any remarks originally contained in the operand of the input statements are printed on the listing. An x (9) preceding the label column of an entry indicates that the instruction has been generated. An apostrophe (10) appears in the listings throughout this manual and indicates the use of an @ character; the ' is equivalent to the @ character on the H type wheel configuration.

The listing on the right-hand portion of the page contains the assembled machine instructions. The card number (11) is assigned by the Output pass of Phase III to each condensed card punched out. When the source statement specifies field definers for less than the full word, the field definition (12) is indicated. The location (13) is the machine location assigned to the instruction. The instruction (14) is the actual assembled machine instruction. The relocation indicators (15) are two-digit numbers assigned each assembled instruction to indicate which part(s) of the instruction must be adjusted if relocation of all or part of the program is desired. The reference address (16) is either (a) the actual address assigned to literals or to DA, DC, or DLINE subsequent entries, or (b) the relative address of DA subsequent entries when a relative address is specified in the DA header line. An M (17) following the reference address column indicates that an error or warning message is associated with that line.

### **Origin Counter Listing**

An Origin Counter Listing follows the Program Listing. Each counter is listed in alphameric order; the initial value, last value, highest value, and lowest value for each counter follow. An example of an Origin Counter Listing appears on page 94.

### **Availability Table**

An Availability Table follows the Origin Counter Listing. This table first lists the electronic switches which remain available for assignment; the index words which are available for assignment follow. If no index words or if no electronic switches remain, the word "NONE" appears in place of the list of switches or index words.

### **Cross Reference Listing**

The Cross Reference Listing is prepared by the Information Transfer pass of Phase III. All symbolic labels, symbolic index words, symbolic switches, CODE fields, literals, and adcons are listed in alphameric order. Each of these entries is followed first by page, line, and actual location of all operand usages of that label.

### **Component Assignment Listing**

The Component Assignment Listing presents all of the electronic switch and index words used in the compiled program. The symbolic name(s), if any, assigned to each of these components is included in this listing.

### **Halt Listing**

A Halt Listing follows the Component Assignment Listing. Each halt instruction occurring in the compiled program will be listed in the order in which it occurs in the program. The halt instructions in the listing will appear in the same format as they appear in the Program Listing.

### **Message Listing**

The Message Listing is provided by the Message pass of Phase III; this listing immediately follows the Halt Listing. The actual message is printed, as well as the page and line number (referred to by (2) and (4) on the Program Listing) of the entry concerned.

## Appendix A:

# Relationship of 7070/7074 Autocoder to Basic Autocoder and Four-Tape Autocoder

The advanced programming capabilities of the 7070/7074 Autocoder system, as compared to the 7070/7074 Four-Tape Autocoder, are due to more powerful macro-instructions, extensive control operations over processing, and increased input/output options. To effect these improvements, distinct statement types and language specifications peculiar to Autocoder have been developed. However, as in Four-Tape Autocoder, the new language characteristics are additions rather than changes to the 7070/7074 Basic Autocoder language. It is therefore possible for either of the more powerful systems to process any program that can be assembled with Basic Autocoder without modification. The converse is untrue since Basic Autocoder is not designed to process the advanced programming features provided in the larger systems. Similarly, since Autocoder and Four-Tape Autocoder each use a unique type of macro-instructions and other programming functions, neither is designed to process *all* programs which can be assembled by the other.

If desired, the user may advance from Four-Tape Autocoder to the Autocoder system by using macro generators to duplicate the substitution-type macros used in a program, or by recoding the program so that macro-instructions provided with or added to the Autocoder system may be utilized. In order for a program coded in Four-Tape Autocoder to be fully compatible with the programming requirements of the Autocoder processor, certain additional changes in the coding and treatment of some programming functions must be made. These differences are outlined in Appendix B.

## Appendix B: Differences Among 7070/7074 Autocoder Systems

The areas in which 7070/7074 Autocoder differs from 7070/7074 Basic Autocoder or 7070/7074 Four-Tape Autocoder are outlined below.

### Coding Sheet, Operand

*Basic Autocoder.* Columns 21-60.

*Four-Tape Autocoder and Autocoder.* Columns 21-75.

### Address Types

A blank address in a LITORIGIN statement has the following significance:

*Basic Autocoder.* Not to be used.

*Four-Tape Autocoder and Autocoder.* Assignment continues at one beyond the highest location previously assigned, except for the locations assigned by the special “s” counter.

The number of symbolic labels that can be used in a source program is as follows:

*Basic Autocoder.* The number is limited by size of storage area available for the symbol table.

*Four-Tape Autocoder.* The symbol table is written as a block on tape when the area in storage is filled. The number is limited only by the number of blocks that can be written on the symbol tape.

*Autocoder.* There is no practical limit to the number of symbolic labels that can be used in a source program.

### Index Words

*Basic Autocoder and Four-Tape Autocoder.* Once reserved, index words cannot be made available for reassignment later in the same program.

*Autocoder.* Index words may be reserved by means of an XRESERVE statement. Index words which have been previously assigned by *any* method may be made available for later assignment by means of an XRELEASE statement.

### Electronic Switches

*Basic Autocoder and Four-Tape Autocoder.* Once reserved, electronic switches cannot be made available for reassignment later in the same program.

*Autocoder.* Electronic switches may be reserved by means of an SRESERVE statement. Electronic switches which have been previously assigned by *any* method may be made available for later assignment by means of an SRELEASE statement.

### DA — Define Area

A DA header line has the following differences:

*Basic Autocoder and Four-Tape Autocoder.* N may be from 1 to 999. The number of storage words that may be reserved for an area is a maximum of 999.

*Autocoder.* N may be from 1 to 9999. The number of words that may be reserved is limited only by the size of the object machine. In addition, a symbolic or actual index word may be appended to the relative address to facilitate the writing of indexed instructions which reference fields in the defined area.

The subsequent entries have the following differences:

*Basic Autocoder and Four-Tape Autocoder.* These entries indicate the names of fields and the position of the field within the area only. A CODE entry is not permitted.

*Autocoder.* The format or characteristics of a field, i.e., numerical and an automatic-decimal number, numerical and a floating-decimal number, etc., may also be included following the indication of the position of the field within the area. A CODE entry may be used.

## DC — Define Constant

*Basic Autocoder and Four-Tape Autocoder.* Numerical constants may contain a maximum of ten digits. In alphameric constants, the @ symbol may appear only if it immediately precedes the terminal @ symbol. The @ symbol may appear in remarks which are on the same line as an alphameric constant. An adcon cannot be modified by address adjustment.

*Autocoder.* Numerical constants may be in standard floating-decimal format or in the form of automatic-decimal numbers of up to twenty digits in length. The @ symbol may appear *anywhere* within an alphameric constant. The @ symbol *cannot* be used in remarks which are on the same line as an alphameric constant.

## EQU — Equate

*Basic Autocoder and Four-Tape Autocoder.* The symbolic address to which a symbolic *name* is being equated must have appeared as a label *earlier* in the sequence of program entries. A symbolic name cannot be equated to a digit value.

*Autocoder.* A symbolic name may be made equivalent to a symbolic address which appears as the label for an entry anywhere in the source program. A symbolic name can be equated to a digit value.

## Origin Control

*Basic Autocoder and Four-Tape Autocoder.* The location assignment counter used by the processor may be set to the starting location specified in the operand of the ORIGIN statement. A blank operand signifies that the contents of the high assignment counter plus 1 is to be used by the processor. The letter “s” following an address in the operand orders the processor not to alter the high assignment counter during the processing of succeeding entries. If an ORIGIN statement does not appear before a location is assigned to the first source-language input entry, the processor will begin the assignment of storage locations at 0325.

*Autocoder.* In addition to the automatic location assignment counter, over 250 separate symbolic-location counters may be named by the programmer and used by the processor to control the placement of a program in storage. A blank operand in an ORIGIN statement indicates that the maximum value attained by *any* location counter, other than counter “s”, is to be used for the assignment of subsequent locations. The assignment of storage locations will begin at an address specified in the Compiler Systems Tape if an ORIGIN entry does not appear before a location is assigned to the first source-language input entry in a program.

## Litorigin Control

*Basic Autocoder.* Normally, literals will be stored in locations immediately following the highest location assigned to the source program. One LITORIGIN statement, placed at the end of the source program deck or immediately preceding the END Control Card (if used) may be utilized to start the assignment of literals at the location specified in the operand rather than following the highest location used in the program.

*Four-Tape Autocoder.* Library subroutines, as well as literals, may be stored in locations immediately following the highest location assigned to the source program. However, the LITORIGIN statement may be used to insert these generated subroutines into the program at the point where the LITORIGIN entry appears. More than one such entry may be used in a program but each must be

followed by an `ORIGIN` statement unless the `LITORIGIN` statement is either the last entry or is followed by an `END` Control operation.

*Autocoder.* Same format as Autocoder `ORIGIN` statement. May be used to regulate the placement of materials generated out-of-line, i.e., implicit adcons, area definitions, etc., as well as literals. Up to 25 `LITORIGIN` statements may be used in a program to cause the assignment of locations to all material generated up to the point at which another `LITORIGIN` entry is encountered. Each phase of a multi-phase program may thus be loaded with its own literals, generated constants, etc. `ORIGIN` statements are not required following a `LITORIGIN` entry.

## End Control

*Basic Autocoder and Four-Tape Autocoder.* If an `END` statement is not used, the processor will generate an unconditional branch to location 0325.

*Autocoder.* If an `END` statement is not used, or if used with a blank operand, the processor will generate an unconditional branch to the address specified in the Compiler Systems Tape.

## Macro-Instructions

*Basic Autocoder.* Cannot process macro-instructions or `DTFS` per se.

*Four-Tape Autocoder.* The 7070 Input/Output Control System macro-instructions may be used with the following restrictions:

1. When using `PUT` macro-instructions with Four-Tape Autocoder, the name preceding `IN` must be the name of either an `RDW` which defines one area or a tape input file. The use of a field name is not allowed unless the field name is the label of an `RDW` that defines the field.
2. The name of a card input file may not be used in the `PUT` macro-instruction if the output file is a tape file.
3. A record from a card input file may be included in a tape output file if the unit record area is defined by one `RDW`; the `PUT` would then be written using the name of the `RDW`; i.e., the name preceding the word `IN` would be the same as the fourth item in the `DUF` entry of the card input file.
4. Only nine tape files may be named in the operand of an `OPEN` or `CLOSE` macro.
5. The symbolic names `IOCSIXF`, `IOCSIXG`, and `IOCSIXH` may not appear in the operand of the `DIOCS` statement.
6. Any index words to be used in the `DIOCS` statement must be actual. These may later be equated to a symbolic name.
7. Comments cards cannot be inserted between `DTFS` or among `DTF` subsequent entries.

Other macro-instructions that can be used are substitution type-values in the operand of the macro in the source program are used to complete the labels, operation codes, etc. of a sequence of instructions contained in the Library portion of the Systems Tape.

*Autocoder.* The 7070 Input/Output Control System macro-instructions can be used without the restrictions noted above. The additional macros that can be used with Autocoder are processed by macro generators and therefore differ from the substitution type of macro-instructions that can be used with Four-Tape Autocoder.

## Appendix C: Reserved Index Words

The Compiler Systems Tape will not allow assignment of *symbolic* names to certain index words. The address and special function of each of these index words are listed below.

<i>Address</i>	<i>Function</i>
0093	Used by Autocoder and FORTRAN floating-decimal subroutines.
0094	Used by Autocoder and FORTRAN normal subroutines.
0095	Used by SPOOL operations on Channel 2.
0096	Used by SPOOL operations on Channel 1.
0097	Priority address word.
0098	Table lookup indexing value and found address.
0099	Address of priority final status word.



## Appendix D: 7070/7074 Operation Codes by Autocoder Mnemonics

In the following list, the symbols in the operand column indicate what is permissible in the operand and the order in which this information must be written on the coding form for each symbolic machine instruction.

In all cases where an "A" has been indicated, a literal may also be used. The list, however, indicates a literal, "L," and also field definition, "F," only where it would seem to be of practical value. Caution is advised when using literals with operation codes which do not specifically indicate them.

### Operand Symbol Key

<i>Symbol</i>	<i>Meaning</i>	<i>Type of Coding</i>	<i>Range of Actual</i>
A	Address	Symbolic or actual	Any storage location
AF	Arm and file	Symbolic or actual	00-03, 10-13, 20-23
B	Blank		
C	Channel number	Symbolic or actual	1-4
CU	Channel and unit	Symbolic or actual	10-49
D	Digit	Actual	0-9
F	Field Definition	Actual (Enclosed in parentheses)	0-9
I	Unit record latch	A or 1, B or 2	1-2
L	Literal		
N	Number	Actual	0-10 (for normal shifts) 0-20 (for coupled or split shifts) 0-9999 (for index word codes)
P	Digit position	Actual (Enclosed in parentheses)	0-9 (for CD) 0-19 (for split shifts)
Q	Inquiry synchronizer	Symbolic or actual	1-2
S	Unit record synchronizer	Symbolic or actual	1-4
SN	Alteration switch	Symbolic or actual	1-4
SW	Electronic switch	Symbolic or actual	1-30
X	Index Word	Symbolic or actual	1-99
,	Used as a separator and must be written on the coding sheet (unless the address which follows is blank).		
/	Used to indicate the word "or" (e.g., A/L means either an address or a literal).		
#	Used to indicate an accumulator number (1, 2, or 3, which must appear in place of the # symbol).		

## Alphabetic List of 7070/7074 Operation Codes by Autocoder Mnemonics

<i>Mnemonic</i>	<i>Operation</i>	<i>Operand</i>	
A#	Add to accumulator #	A/L	F
AA	Add absolute to accumulator 1	A/L	F
AAS#	Add to absolute storage from accumulator #	A	F
AS#	Add to storage from accumulator #	A	F
B	Branch	A	
BAL	Branch if any stacking latch is ON	A	
BAS	Branch if alteration switch is ON	SN,A	
BCB	Branch if channel is busy	C,A	
BCX	Branch compared index word	X,A	
BDL	Branch if disk storage latch is ON	AF,A	
BDX	Branch decremented index word	X,A	
BE	Branch if equal	A	
BES	Branch if electronic switch is ON	SW,A	
BFV	Branch if field overflow	A	
BH	Branch if high	A	
BIX	Branch incremented index word	X,A	
BL	Branch if low	A	
BLX	Branch and load location in index word	X,A	
BM#	Branch if minus in accumulator #	A	
BQL	Branch if inquiry latch is ON	Q,A	
BSC	Branch if sign change	A	
BSF	Branch if electronic switch is ON and set OFF if ON	SW,A	
BSN	Branch if electronic switch is ON and set ON if OFF	SW,A	
BTL	Branch if tape latch is ON	CU,A	
BUL	Branch if unit record latch is ON	I,A	
BV#	Branch if overflow in accumulator #	A	
BXM	Branch if index word is minus	X,A	
BXN	Branch if indexing portion in index word is nonzero	X,A	
BZ#	Branch if zero in accumulator #	A	
C#	Compare accumulator # to storage	A/L	F
CA	Compare absolute in accumulator 1 to absolute in storage	A/L	F
CD	Compare storage to digit	A(P),D	
CSA	Compare sign to alpha	A	
CSM	Compare sign to minus	A	
CSP	Compare sign to plus	A	
D	Divide	A/L	F
DAR	Disk storage arm release	AF	
DLF	Disk storage latch set OFF	AF,A/B	
DLN	Disk storage latch set ON	AF,A/B	
DR	Disk storage read	C,A/L	
DW	Disk storage write	C,A/L	
EAN	Edit alphameric to numerical	X,A/L	
ENA	Edit numerical to alphameric	X,A/L	
ENB	Edit numerical to alphameric with blank insertion	X,A/L	
ENS	Edit numerical to alphameric with sign control	X,A/L	
ESF	Electronic switch set OFF	SW,A/B	
ESN	Electronic switch set ON	SW,A/B	
FA	Floating add	A/L	
FAA	Floating add absolute	A/L	
FAD	Floating add double precision	A/L	

<i>Mnemonic</i>	<i>Operation</i>	<i>Operand</i>
FADS	Floating add double precision and suppress normalization	A/L
FBU	Floating branch underflow	A
FBV	Floating branch overflow	A
FD	Floating divide	A/L
FDD	Floating divide double precision	A/L
FM	Floating multiply	A/L
FR	Floating round	A/B
FS	Floating subtract	A/L
FSA	Floating subtract absolute	A/L
FZA	Floating zero and add	A/L
HB	Halt and branch	A
HMFV	Halt mode for field overflow	A/B
HMSC	Halt mode for sign change	A/B
HP	Halt and proceed	A/B
LE	Lookup equal only	A/L F
LEH	Lookup equal or high	A/L F
LL	Lookup lowest	A/L F
M	Multiply	A/L F
MSA	Make sign alpha	A
MSM	Make sign minus	A
MSP	Make sign plus	A
NOP	No operation	A/B
PC	Priority control	A/L
PDR	Priority disk storage read	C,A/L
PDS	Priority disk storage seek	A
PDW	Priority disk storage write	C,A/L
PR	Priority release	A/B
PTM	Priority tape mark write	CU
PTR	Priority tape read	CU,A/L
PTRR	Priority tape read per record mark control	CU,A/L
PTSB	Priority tape segment backspace	CU,A/L
PTSF	Priority tape segment forward space	CU,A/L
PTSM	Priority tape segment mark write	CU
PTW	Priority tape write	CU,A/L
PTWC	Priority tape write with zero elimination and per record mark control combined	CU,A/L
PTWR	Priority tape write per record mark control	CU,A/L
PTWZ	Priority tape write with zero elimination	CU,A/L
QLF	Inquiry latch set OFF	Q,A/B
QLN	Inquiry latch set ON	Q,A/B
QR	Inquiry read	Q,A/L
QW	Inquiry write	Q,A/L
RG	Record gather	X,A/L
RS	Record scatter	X,A/L
S#	Subtract from accumulator #	A/L F
SA	Subtract absolute from accumulator 1	A/L F
SL	Shift left coupled	N
SL#	Shift left accumulator #	N
SLC	Shift left and count coupled	X
SLC#	Shift left and count accumulator #	X
SLS	Shift left split	N(P)
SMFV	Sense mode for field overflow	A/B

<i>Mnemonic</i>	<i>Operation</i>	<i>Operand</i>
SMSC	Sense mode for sign change	A/B
SR	Shift right coupled	N
SR#	Shift right accumulator #	N
SRR	Shift right and round coupled	N
SRR#	Shift right and round accumulator #	N
SRS	Shift right split	N(P)
SS#	Subtract accumulator # from storage	A F
ST#	Store accumulator #	A F
STD#	Store digits from accumulator # and ignore sign	A F
TEF	Tape end of file turn OFF	CU
TLF	Tape latch set OFF	CU,A/B
TLN	Tape latch set ON	CU,A/B
TM	Tape mark write	CU
TR	Tape read	CU,A/L
TRB	Tape record backspace	CU
TRR	Tape read per record mark control	CU,A/L
TRU	Tape rewind and unload	CU
TRW	Tape rewind	CU
TSB	Tape segment backspace	CU,A/L
TSEL	Tape select	CU
TSF	Tape segment forward space	CU,A/L
TSHD	Tape set high density	CU
TSK	Tape skip	CU
TSLD	Tape set low density	CU
TSM	Tape segment mark write	CU
TW	Tape write	CU,A/L
TWC	Tape write with zero elimination and per record mark control combined	CU,A/L
TWR	Tape write per record mark control	CU,A/L
TWZ	Tape write with zero elimination	CU,A/L
TYP	Type	A/L
ULF	Unit record latch set OFF	I,A/B
ULN	Unit record latch set ON	I,A/B
UP	Unit record punch	S,A/L
UPIV	Unit record punch invalid	S,A/L
UR	Unit record read	S,A/L
US	Unit record signal	S,A/B
UW	Unit record write	S,A/L
UWIV	Unit record write invalid	S,A/L
XA	Index word add to indexing portion	X,A/N
XL	Index word load	X,A/L
XLIN	Index word load and interchange	X,A/L
XS	Index word subtract from indexing portion	X,A/N
XSN	Index word set nonindexing portion	X,A/N
XU	Index word unload	X,A
XZA	Index word zero and add to indexing portion	X,A/N
XZS	Index word zero and subtract from indexing portion	X,A/N
ZA#	Zero accumulator # and add	A/L F
ZAA	Zero accumulator 1 and add absolute	A/L F
ZS#	Zero accumulator # and subtract	A/L F
ZSA	Zero accumulator 1 and subtract absolute	A/L F
ZST#	Zero storage and store accumulator #	A F

## List of 7070/7074 Operation Codes by Function

<i>Mnemonic</i>	<i>Operation</i>	<i>Operand</i>
TAPE OPERATIONS		
TR	Tape read	CU,A/L
PTR	Priority tape read	CU,A/L
TRR	Tape read per record mark control	CU,A/L
PTRR	Priority tape read per record mark control	CU,A/L
TW	Tape write	CU,A/L
PTW	Priority tape write	CU,A/L
TWR	Tape write per record mark control	CU,A/L
PTWR	Priority tape write per record mark control	CU,A/L
TWZ	Tape write with zero elimination	CU,A/L
PTWZ	Priority tape write with zero elimination	CU,A/L
TWC	Tape write with zero elimination and per record mark control combined	CU,A/L
PTWC	Priority tape write with zero elimination and per record mark control combined	CU,A/L
TSB	Tape segment backspace	CU,A/L
PTSB	Priority tape segment backspace	CU,A/L
TSF	Tape segment forward space	CU,A/L
PTSF	Priority tape segment forward space	CU,A/L
TSEL	Tape select	CU
TRB	Tape record backspace	CU
TRW	Tape rewind	CU
TRU	Tape rewind and unload	CU
PTM	Priority tape mark write	CU
TM	Tape mark write	CU
PTSM	Priority tape segment mark write	CU
TSM	Tape segment mark write	CU
TEF	Tape end of file turn OFF	CU
TSHD	Tape set high density	CU
TSLD	Tape set low density	CU
TSK	Tape skip	CU
TLN	Tape latch set ON	CU,A/B
TLF	Tape latch set OFF	CU,A/B
BTL	Branch if tape latch is ON	CU,A
DISK STORAGE OPERATIONS		
PDS	Priority disk storage seek	A
DR	Disk storage read	C,A/L
PDR	Priority disk storage read	C,A/L
DW	Disk storage write	C,A/L
PDW	Priority disk storage write	C,A/L
DAR	Disk storage arm release	AF
DLN	Disk storage latch set ON	AF,A/B
DLF	Disk storage latch set OFF	AF,A/B
BDL	Branch if disk storage latch is ON	AF,A
UNIT RECORD OPERATIONS		
UR	Unit record read	S,A/L
UP	Unit record punch	S,A/L
UW	Unit record write	S,A/L
UPIV	Unit record punch invalid	S,A/L
UWIV	Unit record write invalid	S,A/L

<i>Mnemonic</i>	<i>Operation</i>	<i>Operand</i>
US	Unit record signal	S,A/B
ULN	Unit record latch set ON	I,A/B
ULF	Unit record latch set OFF	I,A/B
BUL	Branch if unit record latch is ON	I,A
INQUIRY OPERATIONS		
QR	Inquiry read	Q,A/L
QW	Inquiry write	Q,A/L
QLN	Inquiry latch set ON	Q,A/B
QLF	Inquiry latch set OFF	Q,A/B
BQL	Branch if inquiry latch is ON	Q,A
CONSOLE TYPEWRITER OPERATIONS		
TYP	Type	A/L
ARITHMETIC OPERATIONS		
A#	Add to accumulator #	A/L F
ZA#	Zero accumulator # and add	A/L F
AA	Add absolute to accumulator 1	A/L F
ZAA	Zero accumulator 1 and add absolute	A/L F
AS#	Add to storage from accumulator #	A F
AAS#	Add to absolute storage from accumulator #	A F
S#	Subtract from accumulator #	A/L F
ZS#	Zero accumulator # and subtract	A/L F
SA	Subtract absolute from accumulator 1	A/L F
ZSA	Zero accumulator 1 and subtract absolute	A/L F
SS#	Subtract accumulator # from storage	A F
M	Multiply	A/L F
D	Divide	A/L F
SHIFTING OPERATIONS		
SR#	Shift right accumulator #	N
SRR#	Shift right and round accumulator #	N
SL#	Shift left accumulator #	N
SLC#	Shift left and count accumulator #	X
SR	Shift right coupled	N
SRR	Shift right and round coupled	N
SL	Shift left coupled	N
SLC	Shift left and count coupled	X
SRS	Shift right split	N(P)
SLS	Shift left split	N(P)
INDEX WORD OPERATIONS		
XL	Index word load	X,A/L
XLIN	Index word load and interchange	X,A/L
XU	Index word unload	X,A
XA	Index word add to indexing portion	X,A/N
XZA	Index word zero and add to indexing portion	X,A/N
XS	Index word subtract from indexing portion	X,A/N
XZS	Index word zero and subtract from indexing portion	X,A/N
XSN	Index word set nonindexing portion	X,A/N
BLX	Branch and load location in index word	X,A

<i>Mnemonic</i>	<i>Operation</i>	<i>Operand</i>
BXM	Branch if index word is minus	X,A
BXN	Branch if indexing portion in index word is nonzero	X,A
BCX	Branch compared index word	X,A
BIX	Branch incremented index word	X,A
BDX	Branch decremented index word	X,A
SWITCH OPERATIONS		
BAS	Branch if alteration switch is ON	SN,A
BES	Branch if electronic switch is ON	SW,A
BSF	Branch if electronic switch is ON and set OFF if ON	SW,A
BSN	Branch if electronic switch is ON and set ON if OFF	SW,A
HMFV	Halt mode for field overflow	A/B
HMSC	Halt mode for sign change	A/B
SMFV	Sense mode for field overflow	A/B
SMSC	Sense mode for sign change	A/B
ESN	Electronic switch set ON	SW,A/B
ESF	Electronic switch set OFF	SW,A/B
DATA MOVEMENT OPERATIONS		
RG	Record gather	X,A/L
RS	Record scatter	X,A/L
EAN	Edit alphameric to numerical	X,A/L
ENA	Edit numerical to alphameric	X,A/L
ENB	Edit numerical to alphameric with blank insertion	X,A/L
ENS	Edit numerical to alphameric with sign control	X,A/L
LOGIC OPERATIONS		
BM#	Branch if minus in accumulator #	A
BZ#	Branch if zero in accumulator #	A
BV#	Branch if overflow in accumulator #	A
BFV	Branch if field overflow	A
BSC	Branch if sign change	A
BH	Branch if high	A
BE	Branch if equal	A
BL	Branch if low	A
BCB	Branch if channel is busy	C,A
BAL	Branch if any stacking latch is ON	A
C#	Compare accumulator # to storage	A/L F
CA	Compare absolute in accumulator 1 to absolute in storage	A/L F
CD	Compare storage to digit	A P ,D
CSA	Compare sign to alpha	A
CSP	Compare sign to plus	A
CSM	Compare sign to minus	A
MISCELLANEOUS OPERATIONS		
B	Branch	A
HB	Halt and branch	A
HP	Halt and proceed	A/B
LE	Lookup equal only	A/L F
LL	Lookup lowest	A/L F
LEH	Lookup equal or high	A/L F
MSA	Make sign alpha	A

<i>Mnemonic</i>	<i>Operation</i>	<i>Operand</i>	
MSP	Make sign plus	A	
MSM	Make sign minus	A	
NOP	No operation	A/B	
PC	Priority control	A/L	
PR	Priority release	A/B	
ST#	Store accumulator #	A	F
STD#	Store digits from accumulator # and ignore sign	A	F
ZST#	Zero storage and store accumulator #	A	F
FLOATING POINT OPERATIONS			
FA	Floating add	A/L	
FZA	Floating zero and add	A/L	
FAA	Floating add absolute	A/L	
FAD	Floating add double precision	A/L	
FADS	Floating add double precision and suppress normalization	A/L	
FS	Floating subtract	A/L	
FSA	Floating subtract absolute	A/L	
FM	Floating multiply	A/L	
FD	Floating divide	A/L	
FDD	Floating divide double precision	A/L	
FR	Floating round	A/B	
FBV	Floating branch overflow	A	
FBU	Floating branch underflow	A	



## Appendix E:

## Note on Optional Characters

In certain cases, special characters used on IBM printers and other equipment have optional equivalents. In each case the character must be punched according to the card code, regardless of which option has been chosen for printing on the printer in a given installation.

The special characters which have been used in this manual and their optional equivalents for each type wheel configuration available are given in the following table.

Character Used in This Manual	IBM Card Code	Type Wheel Configuration									
		A	B	C	D	E	F	G	H	K	
(	0-8-4	%	%	%	%	%	(	%	(	(	
) and □	12-8-4	□	□	□	□	<	)	□	)	)	
@	8-4	@	@	@	@	>	-	-	'	@	
+	12	&	/	&	-	-	+	+	+	+	
=	8-3	#	#	#	#	#	=	+	=	=	
/	0-1	1	&	0	/	&	/	/	/	/	

## Appendix F:

## Glossary

### ACTUAL ADDRESS

The word "actual" usually refers to machine language. An actual address is the same as an absolute or machine address.

### ADDRESS ADJUSTMENT

Address adjustment refers to the procedure of changing an address at process time by means of an increment or decrement placed after the named address, i.e., NAME + 26.

### ADDRESS CONSTANT (ADCON)

An adcon is a numerical literal created by entering  $\pm$  ANYLABEL (a symbolic address written elsewhere in the program) in the operand of an instruction. The actual address assigned to ANYLABEL and the sign of the adcon entry become the address constant.

### ALPHAMERIC

A term which refers to symbols that are numerical digits, alphabetic characters, or special characters.

### ASSEMBLY PROGRAM

A translation program which substitutes machine-language instructions for symbolic instructions, assigns storage locations, and performs other activities necessary to produce the final object program.

### 7070/7074 AUTOCODER

Autocoder is a symbolic programming system consisting of a symbolic language and a processor; this system is designed for use in installations which have a minimum of six 729 Model IV (or Model II) tape units and a machine with 5,000 words of core storage. The language consists of symbolic machine instructions and generator-type macro-instructions; the processor converts a program written in this language to machine language.

### 7070/7074 BASIC AUTOCODER

Basic Autocoder is a symbolic programming system consisting of a symbolic language and a processor; this system is designed for use in installations which have a minimum of one IBM 7500 Card Reader, one IBM 7550 Card Punch, and a machine with 5,000 words of core storage. The language consists of symbolic machine instructions; the processor converts a program written in this language to machine language.

### BRANCH CONSTANT

A branch constant is an instruction which is used to provide the address of a field to a subroutine; it is interrogated by the subroutine but it is never executed.

#### COLLATING SEQUENCE

The relative order of precedence which a computer assigns to the numbers, letters, and special characters for compare operations.

#### COMPILE

To produce a machine-language routine from a routine written in non-machine language. See also: **COMPILER**.

#### COMPILER

A compiler is a complex program in which several different functions are performed. It typically includes the following:

1. Extensive program analysis during which information is collected or tabulated for later recombination.
2. Generation of instructions by synthesis of tabulated information and use of skeleton or model routines.
3. Translation of symbolic instructions into machine language.

A compiler is itself a routine, not a machine—although a machine could be built to do compiling.

#### 7070/7074 COMPILER SYSTEMS TAPE

This tape combines and links 7070/7074 Autocoder, 7070 FORTRAN, and the 7070/7074 Report Program Generator, permitting the use of these compilers with a minimum of effort and a maximum of efficiency.

#### CONDENSED CARD FORMAT

This is a format for placing several machine language instructions on a single card along with information sufficient to load the instructions into their proper storage locations.

#### CONTROL CARD

A card which contains the parameters required to set up a generalized program for one particular application.

#### DEBUGGING

The process of locating errors in a computer routine and correcting them.

#### EXPRESSION

An element of the source language where a combination of several names and operators may be used, as well as a single name or address.

#### FIELD DEFINER(s)

A number placed after an address to indicate the particular digit(s) in a word which are occupied by a field.

#### FIELD DEFINITION

Indication of the starting and ending positions of a field within a word, or the starting and ending positions of a part of a field relative to the field itself; for example, SYMBOL(6,9).

#### 7070 FORTRAN

FORTRAN is a symbolic programming system consisting of a symbolic language and a processor. The language, which closely resembles the language of mathematics, is essentially problem-oriented rather than machine-oriented. The processor converts a program written in this language to 7070/7074 Autocoder language. The Autocoder program is converted to a machine-language object program for a 7070 or a 7074 by the Autocoder processor. Other IBM Data Processing Systems are equipped with processors which convert programs coded in the FORTRAN language to their respective machine-oriented languages.

#### 7070/7074 FOUR-TAPE AUTOCODER

Four-Tape Autocoder is a symbolic programming system consisting of a symbolic language and a processor; this system is designed for use in installations which have a minimum of four 729 Model IV (or Model II) tape units and a machine with 5,000 words of core storage. The language consists of symbolic machine instructions and substitution-type macro-instructions; the processor converts a program written in this language to machine language.

#### GENERATOR

A program which accepts input parameters and uses them to modify skeleton instructions or skeleton routines to produce the desired output routine. A small number of parameters are capable of producing a large number of output instructions.

#### INDEXING

Refers to the procedure of changing an address, at object time, according to the contents of the indexing portion of a specified index word.

#### INSTRUCTION

An instruction is generally a single entry in symbolic machine language or in machine language, as opposed to a statement, or macro-instruction, which usually means a language entry that can produce many machine-language entries.

#### LITERAL

A literal is the actual data to be operated on by an instruction, as distinguished from the location or address of the data.

#### MACHINE LANGUAGE

A language in which the instructions are in a form which may be executed by the computer without translation.

#### MACRO GENERATOR

A macro generator, an abbreviation of macro-instruction generator, is that part of Autocoder which processes macro statement entries. It generates a sequence of symbolic machine instructions which, collectively, perform the operation specified by the macro-instruction entry. Each macro generator is associated with a given macro statement and processes all occurrences of that macro statement in a program.

#### MACRO-INSTRUCTION

A symbolically coded instruction resulting in a group of machine-language instructions which will perform a desired operation.

#### MNEMONIC

“Mnemonic” means “aiding memory.” The term is used to describe operation codes which are written in a symbolic notation to make them easier to remember than the actual operation codes.

#### NUMERICAL

In the 7070 or the 7074, numerical refers to either a field with a plus or minus sign rather than an alpha sign or to data which is to be treated as numbers, whether in single-digit or double-digit form.

#### OBJECT MACHINE

The machine on which an object program is to be run. See: **PROCESSOR MACHINE**.

#### OBJECT PROGRAM

The output from a processor. In this case, a 7070/7074 machine-language program assembled from a source program coded in symbolic language.

#### OBJECT TIME

The time at which the object program is being run. This is opposed to process time, the time at which a compiler, such as Autocoder, is being run.

#### OFF-LINE

Operation of input/output and other devices not under direct computer control. Most commonly used to designate the transfer of information between magnetic tapes and other input/output media.

#### ON-LINE

Operation of an input/output device as a component of the computer, under programmed control.

#### OPERAND

The operand is the factor or data acted upon during an operation; it may be a result, a parameter, or an indication of the location of the next instruction. Also, the entire field beginning in column 21 of the Autocoder source card and coding sheet is considered to be the operand.

#### OPERATION CODE

The operation code designates the machine function to be performed. Distinction is sometimes made between a "mnemonic operation code," such as ZAL, and its equivalent "machine operation code," +13.

#### OPERATOR

The term operator usually refers to such characters as +, -, =, which are said to "operate" on quantities.

#### PARAMETER

A factor which is left unspecified and to which the user may assign a value.

#### PROCESS TIME

The time at which the source program is being changed into an object program by a compiler, such as Autocoder. This is opposed to object time, the time at which the object program is being run.

#### PROCESSOR

A program which performs the functions of assembly, compilation, generation, or any similar functions to convert a source program into the desired object program.

#### PROCESSOR MACHINE

The machine on which a processor is to be run.

#### PROGRAMMING SYSTEM

A programming system is any method of programming problems other than machine language, such as Autocoder. A system consists of a language and its associated processor(s).

#### REPORT

A printed document which presents data arranged in an orderly manner for ease of reference.

#### REPORT GENERATION

A technique for producing complete reports given the content and format of the input file and the desired content and format of the output reports.

#### 7070/7074 REPORT PROGRAM GENERATOR

The 7070/7074 Report Program Generator consists of coding forms and a compiler. The format of the input file and the specifications for the output report are placed on the coding sheets which are then punched into cards and used as input for the compiler. The system produces a program in 7070/7074 Autocoder language which is converted to a machine-language object program for a 7070 or a 7074 by the Autocoder processor.

#### **SORT**

To place a file of records in order according to a specified sequence.

#### **SOURCE LANGUAGE**

The language in which a program is coded, e.g., the 7070/7074 Autocoder language.

#### **SOURCE PROGRAM**

The original coding of a program, usually coding in a language other than machine language.

#### **SPECIAL CHARACTER**

One of a set of special symbols. Some common special characters are:

# \$ + - \* ( ) / , . □ =

#### **STATEMENT**

Usually a source-language entry on the coding sheet, especially a line which might eventually produce several machine-language instructions, such as the ARITH statement or the GET statement.

#### **STORAGE**

Any medium into which data may be transferred and where it may be retained for later use.

#### **SUBROUTINE**

A subroutine is usually a series of instructions to perform some specific mathematical or logical operation. Subroutines are entered by a Branch instruction, as opposed to macro-instructions which are normally entered sequentially.

#### **SYMBOL**

In Autocoder, symbol is used to refer to a name used instead of a machine address. Thus "symbolic address," "symbolic name," or "symbolic label," conveys that one is not specifying machine addresses.

#### **SYMBOLIC ADDRESS**

An alphameric name used in place of an actual machine address.

#### **SYMBOLIC LANGUAGE**

A symbolic language is a collection of mnemonic symbols with rules of usage, such as Autocoder; the symbols are used in programming to represent operation codes, functions, and/or addresses. Symbolic-language coding must be translated into machine language to be used by the computer.

#### **SYMBOLIC MACHINE LANGUAGE**

Symbolic machine language is a language which is similar to machine language except for symbolic addresses and mnemonic operation codes. Symbolic machine instructions are sometimes referred to as "one-for-one" instructions since each instruction encountered in a source program will cause the corresponding machine language instruction to be inserted in the object program.

## Appendix G: Illustration of Autocoder Programming

### Input Statements

The listing on the following page contains the input statements of a sample payroll problem. The purpose of this problem is to illustrate coding with 7070/7074 Autocoder, including the use of input/output statements and macro-instructions.

DA statements define the formats of the input, output, and detail files, the file from which checks will be printed, and temporary storage fields. A DLINE statement describes the check format. An error message and a space for the insertion of an identification number are provided under a DC statement, for use if the desired record fails to appear in the master file. The processing routine follows the declarative statements in the listing.

### Generated Coding

The listings which follow contain coding generated from the previously described input statements. This listing is incomplete; the DIF entries, DIOS entries, and the major portion of the generated input/output instructions have been omitted.



```

AF01      *      AN ILLUSTRATION OF AUTOCODER PROGRAMMING
AF02      *
AF03      *      DESCRIPTION OF THE INPUT MASTER FILE
AF04      IMASTER DA      10,RDW,0+IMASTERX
AF05      IMANNNUMBER      00,09A10
AF06      INAME            10,39'
AF07      IDEPENDNTS      40,41A2
AF07A     IPAYRATE        42,46A2.3
AF08      ICOMISRATE      47,49A.3
AF09      IYTDPAY         50,56A5.2
AF10      IYTDFICA        65,69A3.2
AF11      *
AF12      *      DESCRIPTION OF THE OUTPUT MASTER FILE
AF13      MASTEROUT DA     10,RDW,0+OMASTERX
AF14      *              00,69
AF15      *
AF16      *      DESCRIPTION OF THE DETAIL FILE
AF17      DETAIL DA       20,RDW,0+DETAILX
AF18      DMANNUMBER      00,09A10
AF19      DHOOURS         17,19A2.1
AF20      DSALES          22,29A6.2
AF21      *
AF22      *      DESCRIPTION OF THE CHECK FORMAT
AF23      CHECKLINE DLIN  22,ZZZ,ZZZ,ZZZ
AF24      CMANNUMBER      18,32
AF25      CNAME           40$XX,XXZ+ZZ
AF26      CNETPAY         40$XX,XXZ+ZZ
AG01      *      DESCRIPTION OF THE FILE FROM WHICH CHECKS WILL BE PRINTED
AG02      CHECKTAPE DA    10,RDW,0+CHECKX
AG03      *              00,99
AG04      *
AG05      *      DESCRIPTION OF TEMPORARY STORAGE FIELDS
AG06      WORKAREA DA     1
AG07      GROSSPAY        03,09A5.2
AG08      TAX             13,19A5.2
AG09      FICA            25,29A3.2
AG10      TFICA          35,39A3.2
AG11      NETPAY          43,49A5.2
AG12      *
AG13      ERMESAGE DC     -RDW
AG14      *              'MASTER IS MISSING FOR MANNUMBER'
AG15      ERRORNO        '
AG16      *
AG17      *      PROGRAM
AG18      IOPEN          OPEN  IMASTER,DETAIL,MASTEROUT,CHECKTAPE
AG19      START          GET   DETAIL
AG20      NEXTMASTER    GET   IMASTER
AG21      COMP           DMANNUMBER,IMANNUMBER,NOMASTER,,NODETAIL
AG22      ARITH          GROSSPAY=IPAYRATE*DHOOURS+USALES*ICOMISRATE
AG23      ARITH          TAX=.18*(GROSSPAY-IDEPENDNTS*13.00)
AG24      ZSIGN          TAX,,ZEROTAX
AG25      FICATEST      ARITH FICA=GROSSPAY*.03
AH01      ARITH          TFICA=IYTDFICA+FICA-144.00
AH02      ZSIGN          TFICA,FICALC,,FICALC
AH03      ARITH          FICA=FICA-IFICA
AH04      FICALC        ARITH IYTDFICA=IYTDIFICA+FICA
AH05      ARITH          IYTDPAY=IYTDPAY+GROSSPAY
AH06      ARITH          NETPAY=GROSSPAY-TAX-FICA
AH07      EDMOV         IMANNUMBER TO CMANNUMBER,INAME TO CNAME,NETPAY TO
AH08      *              CNETPAY
AH08A     *      PREPARE TAPE RECORD FOR PRINTING CHECKS OFFLINE
AH09      PUT           CHECKLINE IN CHECKTAPE
AH10      PUTX          IMASTER IN MASTEROUT
AH11      B             START
AH12      ZEROTAX       ZERO  TAX
AH13      B             FICATEST
AH14      NOMASTER     MOVE  DMANNUMBER TO ERRORNO
AH15      TYP           ERMESAGE
AH16      NOP
AH17      B             START
AH18      NODETAIL     PUTX  IMASTER IN MASTEROUT
AH19      B             NEXTMASTER
AH20      EOFDETAIL    BSN   1,IEND
AH21      RNOUSTR      PUTX  IMASTER IN MASTEROUT
AH22      GET           IMASTER
AH23      B             RNOUSTR
AH24      EOFMASTER    BSN   1,IEND
AH25      RNOUSTR      MOVE  DMANNUMBER TO ERRORNO
AI01      TYP           ERMESAGE
AI02      NOP
AI03      GET           DETAIL
AI04      B             RNOUSTR
AI05      IEND         END
AI06      END          CNTRL IOPEN

```

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PAGE AY		PROGRAM							PAGE AY
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
34	AF01	*		AN ILLUSTRATION OF AUTOCODER PROGRAMMING					
35	AF02	*							
36	AF03	*		DESCRIPTION OF THE INPUT MASTER FILE					
37	AF04	IMASTER	DA	10:RDW*0+IMASTERX				+0013641443	
38		X			00219		1364	+0013741380	1364
39		X					1365	+0013811387	1365
40		X					1366	+0013881394	1366
41		X					1367	+0013951401	1367
42		X					1368	+0014021408	1368
43		X			00220		1369	+0014091415	1369
44		X					1370	+0014161422	1370
45		X					1371	+0014231429	1371
46		X					1372	+0014301436	1372
47		X					1373	-0014371443	1373
48	AF05	IMANNUMBER		00:09A10	09		1374		0000

PAGE AZ		PROGRAM							PAGE AZ
LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	AF06	INAME		10:391			09	1375	0001
02	AF07	IDPENDNTS		40:41A2			01	1378	0004
03	AF07A	IPAYRATE		42:46A2.3			26	1378	0004
04	AF08	ICOMISRATE		47:49A.3			79	1378	0004
05	AF09	IYDPAY		50:56A5.2			06	1379	0005
06	AF10	IYDFICA		65:69A3.2			59	1380	0006
07	AF11	*							
08	AF12	*		DESCRIPTION OF THE OUTPUT MASTER FILE					
09	AF13	MASTEROUT	DA	10:RDW*0+OMASTERX				+0014441523	
10		X			00221		1444	+0014541460	1444
11		X					1445	+0014611467	1445
12		X					1446	+0014681474	1446
13		X					1447	+0014751481	1447
14		X					1448	+0014821488	1448
15		X			00222		1449	+0014891495	1449
16		X					1450	+0014961502	1450
17		X					1451	+0015031509	1451
18		X					1452	+0015101516	1452
19		X					1453	-0015171523	1453
20	AF14			00:69	09		1454		0000
21	AF15	*							
22	AF16	*		DESCRIPTION OF THE DETAIL FILE					
23	AF17	DETAIL	DA	20:RDW*0+DETAILX				+0015241603	
24		X			00223		1524	+0015441546	1524
25		X					1525	+0015471549	1525
26		X					1526	+0015501552	1526
27		X					1527	+0015531555	1527
28		X					1528	+0015561558	1528
29		X			00224		1529	+0015591561	1529
30		X					1530	+0015621564	1530
31		X					1531	+0015651567	1531
32		X					1532	+0015681570	1532
33		X					1533	+0015711573	1533
34		X			00225		1534	+0015741576	1534
35		X					1535	+0015771579	1535
36		X					1536	+0015801582	1536
37		X					1537	+0015831585	1537
38		X					1538	+0015861588	1538
39		X			00226		1539	+0015891591	1539
40		X					1540	+0015921594	1540
41		X					1541	+0015951597	1541
42		X					1542	+0015981600	1542
43		X					1543	-0016011603	1543
44	AF18	DMANNUMBER		00:09A10	09		1544		0000
45	AF19	DHOURS		17:19A2.1	79		1545		0001
46	AF20	DSALES		22:29A6.2	29		1546		0002
47	AF21	*							
48	AF22	*		DESCRIPTION OF THE CHECK FORMAT					

PAGE	BA	PROGRAM				CDNO	FD	LOC	INSTRUCTION	PAGE	BA
LN	CDREF	LABEL	OP	OPERAND						REF	
01	AF23	CHECKLINE	DLINE			00228			+0016041613		
02	AF24	CMANNUBER		2Z,ZZZ,ZZZ,ZZZ			29	1604		1604	
03	AF25	CNAME		18,32		00229	49	1607		1607	
07	AF26	CNETPAY		40\$XX,XXZ,ZZ			89	1611		1611	
08	AG01	*		DESCRIPTION OF THE FILE FROM WHICH CHECKS WILL BE PRINTED							
09	AG02	CHECKTAPE	DA	10,RDW,0+CHECKX					+0016141723		
10		X				00230		1614	+0016241633	1614	
11		X						1615	+0016341643	1615	
12		X						1616	+0016441653	1616	
13		X						1617	+0016541663	1617	
14		X						1618	+0016641673	1618	
15		X				00231		1619	+0016741683	1619	
16		X						1620	+0016841693	1620	
17		X						1621	+0016941703	1621	
18		X						1622	+0017041713	1622	
19		X						1623	-0017141723	1623	
20	AG03			00,99			09	1624		0000	
21	AG04	*									
22	AG05	*		DESCRIPTION OF TEMPORARY STORAGE FIELDS							
23	AG06	WORKAREA	DA	1					+0017241728		
24	AG07	GROSSPAY		03,09A5.2			39	1724		1724	
25	AG08	TAX		13,19A5.2			39	1725		1725	
26	AG09	FICA		25,29A3.2			59	1726		1726	
27	AG10	TFICA		35,39A3.2			59	1727		1727	
28	AG11	NETPAY		43,49A5.2			39	1728		1728	
29	AG12	*									
30	AG13	ERMESAGE	DC	-RDW					+0017291738		
31		X				00232		1729	-0017301738	1729	
32	AG14			'MASTER IS MISSING FOR MANNUMBER'		00233	09	1730	'7461828365	1730	
33		X					09	1731	'7900698200	1731	
34		X					09	1732	'7469828269	1732	
35		X					09	1733	'7567006676	1733	
36		X					09	1734	'7900746175	1734	
37		X				00234	09	1735	'7584746265	1735	
38		X					01	1736	'79	1736	
39	AG15	ERRORNO					29	1736	' 00000000	1736	
40		X					09	1737	'0000000000	1737	
41		X					01	1738	'00	1738	
42	AG16	*									
43	AG17	*	PROGRAM								
44	AG18	IOPEN	OPEN	IMASTER,DETAIL,MASTEROUT,CHECKTAPE							
45		X	IOPEN	BLX	IOCSIXG,IOC,IOPEN	00235		1739	+0200040578		
46		X		B	TAPEFILEIM			1740	+0100091327		
47		X		B	TAPEFILEDI			1741	+0100091345		
48		X		B	TAPEFILEMO			1742	+0100091336		

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	X		B	TAPEFILECO					
02	X		NOP						
03	AG19	START	GET	DETAIL	00236		1743	+0100091354	
04	X	START	BIX	DETAILB,**2			1744	-0100000000	
05	X		BLX	DETAILX,IOC.NSE03A			1745	+4900051747	
06	X		XL	DETAILX,0+DETAILB			1746	+0200062039	
07	AG20	NEXTMASTER	GET	IMASTER			1747	+4505060000	
08	X	NEXTMASTER	BIX	IMASTERB,**2					
09	X		BLX	IMASTERX,IOC.NSE01A	00237		1748	+4900071750	
10	X		XL	IMASTERX,0+IMASTERB			1749	+0200081945	
11	AG21		COMP	DMANNUMBER,IMANNUMBER,NOMASTER,,NODetail			1750	+4507080000	
12	X		ZA2	DMANNUMBER(0,9)+DETAILX			1751	+2306090000	
13	X		S2	IMANNUMBER(0,9)+IMASTERX			1752	-2408090000	
14	X		BV2	**2			1753	+2100091755	
15	X		BZ2	**3			1754	+2000091757	
16	X		BM2	NOMASTER	00238		1755	-2000091881	
17	X		B	NODetail			1756	+0100091891	
18	AG22		ARITH	GROSSPAY=IPAYRATE*DHOURLS+DSALES*ICOMISRATE					
19	X		ZA3	IPAYRATE(0,4)+IMASTERX			1757	+3308260004	M
20	X		M	DHOURLS(0,2)+DETAILX			1758	+5306790001	
21	X		SR	1	00239		1759	-5000000001	
22	X		ST2	COMAREA.A(0,9)+1			1760	+2200092122	
23	X		ZA3	DSALES(0,3)+DETAILX			1761	+3306250002	
24	X		M	ICOMISRATE(0,2)+IMASTERX			1762	+5308790004	
25	X		A2	COMAREA.A(0,9)+1			1763	+2400092122	
26	X		SRR2	1	00240		1764	+5000002101	
27	X		ST2	GROSSPAY(0,6)			1765	+2200391724	
28	AG23		ARITH	TAX=.18*(GROSSPAY-IDEPENDNTS*13.00)					
29	X		ZS3	IDEPENDNTS(0,3)+IMASTERX			1766	-3308030004	M
30	X		M	+1300			1767	+5300142124	
31	X		A2	GROSSPAY(0,6)			1768	+2400391724	
32	X		ZA3	+18	00241		1769	+3300232125	
33	X		M	9992			1770	+5300099992	
34	X		SRR	2			1771	-5000000102	
35	X		ST2	TAX(0,6)			1772	+2200391725	
36	AG24		ZSIGN	TAX,,,ZEROTAX					
37	X		ZA1	TAX(0,6)			1773	+1300391725	
38	X		BZ1	M.24	00242		1774	+1000091777	
39	X		CSM	TAX			1775	-0300601725	
40	X		BE	ZEROTAX			1776	-4100091878	
41	X	M.24	NOP				1777	-0100000000	
42	X	ORIGIN	CNTRL	*-1					
43	AG25	FICATEST	ARITH	FICA=GROSSPAY*.03					
44	X	FICATEST	ZA3	GROSSPAY(0,6)			1777	+3300391724	
45	X		M	+03			1778	+5300012125	
46	X		SRR	2	00243		1779	-5000000102	
47	X		ST2	FICA(0,4)			1780	+2200591726	
48	AH01		ARITH	TFICA=1YDFICA+FICA-144.00					

PAGE	BC	PROGRAM	CDNO	FD	LOC	INSTRUCTION	PAGE	BC
LN	CDREF	LABEL	OP	OPERAND			REF	
01	X		ZA2	FICA(0,3)		1781 +2300581726		
02	X		SL2	2		1782 +5000002202		
03	X		A2	IYDFICA(0,4)+IMASTERX		1783 +2408590006		
04	X		S2	+14400	00244	1784 -2400592124		
05	X		ST2	TFICA(0,4)		1785 +2200591727		
06	AH02		ZSIGN	TFICA,FICALC,,FICALC				
07	X		ZA1	TFICA(0,4)		1786 +1300591727		
08	X		BZ1	FICALC		1787 +1000091794		
09	X		CSM	TFICA		1788 -0300601727		
10	X		BE	FICALC	00245	1789 -4100091794		
11	AH03		ARITH	FICA=FICA-TFICA				M
12	X		ZS2	TFICA(0,3)		1790 -2300581727		
13	X		SL2	2		1791 +5000002202		
14	X		A2	FICA(0,4)		1792 +2400591726		
15	X		ST2	FICA(0,4)		1793 +2200591726		
16	AH04	FICALC	ARITH	IYDFICA=IYDFICA+FICA				M
17	X	FICALC	ZA2	FICA(0,3)	00246	1794 +2300581726		
18	X		SL2	2		1795 +5000002202		
19	X		A2	IYDFICA(0,4)+IMASTERX		1796 +2408590006		
20	X		ST2	IYDFICA(0,4)+IMASTERX		1797 +2208590006		
21	AH05		ARITH	IYDTPAY=IYDTPAY+GROSSPAY				M
22	X		ZA2	GROSSPAY(0,3)		1798 +2300361724		
23	X		SL2	2	00247	1799 +5000002202		
24	X		A2	IYDTPAY(0,6)+IMASTERX		1800 +2408060005		
25	X		ST2	IYDTPAY(0,6)+IMASTERX		1801 +2208060005		
26	AH06		ARITH	NETPAY=GROSSPAY-TAX-FICA				M
27	X		ZS2	TAX(0,3)		1802 -2300361725		
28	X		S2	FICA(0,3)		1803 -2400581726		
29	X		SL2	2	00248	1804 +5000002202		
30	X		A2	GROSSPAY(0,6)		1805 +2400391724		
31	X		ST2	NETPAY(0,6)		1806 +2200391728		
32	AH07		EDMOV	IMANNUMBER TO CMANNUMBER, INAME TO CNAME, NETPAY TO				
33	AH08			CNETPAY				
34	X		ZA2	IMANNUMBER(0,9)+IMASTERX		1807 +2308090000		
35	X		ST2	COMAREA.A+2		1808 +2200092123		
36	X		XZA	MACREG.01,COMAREA.A+2	00249	1809 +4600092123		
37	X		ENA	MACREG.01,EDMOV02.A		1810 +5600092111		
38	X		ZA3	COMAREA.A(0,1)		1811 +3300012121		
39	X		ST3	CMANNUMBER(0,1)		1812 +3200231604		
40	X		ZA3	' , '		1813 +3300452127		
41	X		ST3	CMANNUMBER(2,3)	00250	1814 +3200451604		
42	X		ZA3	COMAREA.A(2,5)		1815 +3300252121		
43	X		ST3	CMANNUMBER(4,7)		1816 +3200691604		
44	X		ZA3	COMAREA.A(6,7)		1817 +3300672121		
45	X		ST3	CMANNUMBER(8,9)		1818 +3200011605		
46	X		ZA3	' , '	00251	1819 +3300452127		
47	X		ST3	CMANNUMBER(10,11)		1820 +3200231605		
48	X		ZA3	COMAREA.A(8,9)		1821 +3300892121		

LN	CDREF	LABEL	OP	OPERAND	CDNO	FD	LOC	INSTRUCTION	REF
01	X		ST3	CMANNUBER(12,13)			1822	+3200451605	
02	X		ZA3	COMAREA.A(10,13)			1823	+3300032122	
03	X		ST3	CMANNUBER(14,17)	00252		1824	+3200691605	
04	X		ZA3	' , '			1825	+3300452127	
05	X		ST3	CMANNUBER(18,19)			1826	+3200011606	
06	X		ZA3	COMAREA.A(14,19)			1827	+3300492122	
07	X		ST3	CMANNUBER(20,25)			1828	+3200271606	
08	X		ZA1	INAME(0,9)+IMASTERX	00253		1829	+1308090001	
09	X		ZA2	INAME(10,19)+IMASTERX			1830	+2308090002	
10	X		ST2	CNAME(16,19)			1831	+2200031609	
11	X		SR	4			1832	-5000000004	
12	X		ST2	CNAME(6,15)			1833	+2200091608	
13	X		ST1	CNAME(0,5)	00254		1834	+1200491607	
14	X		ZA2	INAME(20,29)+IMASTERX			1835	+2308090003	
15	X		ST2	CNAME(26,29)			1836	+2200031610	
16	X		SR2	4			1837	+5000002004	
17	X		ST2	CNAME(20,25)			1838	+2200491609	
18	X		ZA2	NETPAY(0,6)	00255		1839	+2300391728	
19	X		ST2	COMAREA.A+2			1840	+2200092123	
20	X		XZA	MACREG.01,COMAREA.A+2			1841	+4600092123	
21	X		ENA	MACREG.01,EDMOV02.A			1842	+5600092111	
22	X		ZA3	' '			1843	+3300092126	
23	X		ST3	CNETPAY(0,1)	00256		1844	+3200891611	
24	X		ST3	CNETPAY(2,9)			1845	+3200071612	
25	X		SLC2	MACREG.02			1846	+5000102300	
26	X		ZA3	' 5 '			1847	+3300232127	
27	X		B	M.22-3+MACREG.02			1848	+0110092109	
28	X	M.23	ST3	CNETPAY(0,1)	00257		1849	+3200891611	
29	X		ZA3	COMAREA.A(6,7)			1850	+3300672121	
30	X		ST3	CNETPAY(2,3)			1851	+3200011612	
31	X		ZA3	COMAREA.A(8,9)			1852	+3300892121	
32	X		ST3	CNETPAY(4,5)			1853	+3200231612	
33	X		ZA3	' , '	00258		1854	+3300452127	
34	X		ST3	CNETPAY(6,7)			1855	+3200451612	
35	X		ZA3	COMAREA.A(10,11)			1856	+3300012122	
36	X		ST3	CNETPAY(8,9)			1857	+3200671612	
37	X		ZA3	COMAREA.A(12,13)			1858	+3300232122	
38	X		ST3	CNETPAY(10,11)	00259		1859	+3200891612	
39	X		ZA3	COMAREA.A(14,15)			1860	+3300452122	
40	X		ST3	CNETPAY(12,13)			1861	+3200011613	
41	X		ZA3	' . '			1862	+3300012127	
42	X		ST3	CNETPAY(14,15)			1863	+3200231613	
43	X		ZA3	COMAREA.A(16,17)	00260		1864	+3300672122	
44	X		ST3	CNETPAY(16,17)			1865	+3200451613	
45	X		ZA3	COMAREA.A(18,19)			1866	+3300892122	
46	X		ST3	CNETPAY(18,19)			1867	+3200671613	
47	AH08A	*		PREPARE TAPE RECORD FOR PRINTING CHECKS OFFLINE					
48	AH09		PUT	CHECKLINE IN CHECKTAPE					

PAGE BE		PROGRAM				PAGE BE		
LN	CDREF	LABEL	OP	OPERAND	CDNO FD	LOC	INSTRUCTION	REF
01		X	BIX	CHECKB,**2		1868	+4900111870	
02		X	BLX	CHECKX,IOC.NSE04A	00261	1869	+0200122086	
03		X	XL	CHECKX,0+CHECKB		1870	+4511120000	
04		X	RG	CHECKX,IOC.PUT001		1871	-6500122120	
05	AH10		PUTX	IMASTER IN MASTEROUT				
06		X	BIX	OMASTERB,**2		1872	+4900131874	
07		X	BLX	OMASTERX,IOC.NSE02A		1873	+0200141992	
08		X	XL	OMASTERX,0+OMASTERB	00262	1874	+4513140000	
09		X	XU	IMASTERX,0+OMASTERB		1875	-4513080000	
10		X	XU	OMASTERX,0+IMASTERB		1876	-4507140000	
11	AH11		B	START		1877	+0100091745	
12	AH12	ZEROTAX	ZERO	TAX				
13		X	ZEROTAX	ZA2		1878	+2300002124	
14		X	STD2	TAX(0,6)	00263	1879	-2200391725	
15	AH13		B	FICATEST		1880	+0100091777	
16	AH14	NOMASTER	MOVE	DMANNUBER TO ERRORNO				
17		X	NOMASTER	ZA2		1881	+2306090000	
18		X	ST2	ERRORNO(8,9)		1882	+2200011737	
19		X	SR2	2		1883	+5000002002	
20		X	ST2	ERRORNO(0,7)	00264	1884	+2200291736	
21		X	SL	20		1885	-5000000220	
22		X	ST2	ERRORNO(10,17)		1886	+2200291737	
23		X	ST2	ERRORNO(18,19)		1887	+2200011738	
24	AH15		TYP	ERMESAGE		1888	+6900041729	
25	AH16		NOP		00265	1889	-0100090000	
26	AH17		B	START		1890	+0100091745	
27	AH18	NODETAIL	PUTX	IMASTER IN MASTEROUT				
28		X	NODETAIL	BIX		1891	+4900131893	
29		X	BLX	OMASTERX,IOC.NSE02A		1892	+0200141992	
30		X	XL	OMASTERX,0+OMASTERB		1893	+4513140000	
31		X	XU	IMASTERX,0+OMASTERB	00266	1894	-4513080000	
32		X	XU	OMASTERX,0+IMASTERB		1895	-4507140000	
33	AH19		B	NEXTMASTER		1896	+0100091748	
34	AH20	EOFDETAIL	BSN	1, IEND		1897	+6100301921	
35	AH21	RNOUTMASTR	PUTX	IMASTER IN MASTEROUT				
36		X	RNOUTMASTR	BIX		1898	+4900131900	
37		X	BLX	OMASTERX,IOC.NSE02A	00267	1899	+0200141992	
38		X	XL	OMASTERX,0+OMASTERB		1900	+4513140000	
39		X	XU	IMASTERX,0+OMASTERB		1901	-4513080000	
40		X	XU	OMASTERX,0+IMASTERB		1902	-4507140000	
41	AH22		GET	IMASTER				
42		X	BIX	IMASTERB,**2		1903	+4900071905	
43		X	BLX	IMASTERX,IOC.NSE01A	00268	1904	+0200081945	
44		X	XL	IMASTERX,0+IMASTERB		1905	+4507080000	
45	AH23		B	RNOUTMASTR		1906	+0100091898	
46	AH24	EOFMASTER	BSN	1, IEND		1907	+6100301921	
47	AH25	RNOUTDTAIL	MOVE	DMANNUBER TO ERRORNO				
48		X	RNOUTDTAIL	ZA2		1908	+2306090000	

PAGE BF		PROGRAM						PAGE BF	
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01	X		ST2	ERRORNO(8,9)	00269	1909	+2200011737		
02	X		SR2	2		1910	+5000002002		
03	X		ST2	ERRORNO(0,7)		1911	+2200291736		
04	X		SL	20		1912	-5000000220		
05	X		ST2	ERRORNO(10,17)		1913	+2200291737		
06	X		ST2	ERRORNO(18,19)	00270	1914	+2200011738		
07	A101		TYP	ERMESAGE		1915	+6900041729		
08	A102		NOP			1916	-0100090000		
09	A103		GET	DETAIL					
10	X		BIX	DETAILB,**2		1917	+4900051919		
11	X		BLX	DETAILX,IOC.NSE03A		1918	+0200062039		
12	X		XL	DETAILX,G+DETAILB	00271	1919	+4505060000		
13	A104		B	RNOUTDTAIL		1920	+0100091908		
14	A105	IEND	END						
15	X	IEND	BLX	IOCSIXG,IOC.IEND		1921	+0200041227		
16	X		NCP	0					



# Appendix H: Index of Messages

<i>Message</i>	<i>Macro</i>	<i>Page</i>	<i>Message</i>	<i>Macro</i>	<i>Page</i>
ACTUAL ADDRESS NOT ALLOWED.....	DUF	30	INCORRECT NUMBER OF PARAMETERS.....	SHIFT	222
ALL BRANCHES BLANK.....	COMP	145	INCORRECT OPTION.....	SHIFT	222
ALL BRANCHES EQUAL.....	COMP	145	INSTRUCTION NOT PROGRAM SWITCH. PAR. <i>xx</i> ....	SETSW	185
ALPHA BLANKS INTO UNDEFINED PAR. <i>xx</i> .....	ZERO	189	INVALID PARAMETER <i>xx</i> .....	LOGIC	170
AN ELEMENT OF THE OPERAND STARTS ILLEGALLY	DUF	30		SNAP	226
ASSUME COMMAS AFTER PAR.2 — PERMANENT NOP	CYCLE	156	INVALID SWITCH. PARAMETER <i>xx</i> .....	SETSW	186
ATTEMPTING TO FILL HARDWARE. PAR. <i>xx</i> .....	FILL	197	IOCSIXF(C, H) ENTRY OUT OF RANGE, IGNORED...	DIOCS	25
ATTEMPTING TO ZERO HARDWARE. PAR. <i>xx</i> .....	ZERO	189	LABEL SHOULD BE BLANK.....	DUF	30
BLANK PARAMETER <i>xx</i> .....	EDMOV	209	LESS THAN 3 INPUT PARAMETERS.....	COMP	146
	SNAP	226	MULTIPLE FROM- AND TO-FIELDS.....	MOVE	215
BRANCH TO NON-IMPERATIVE INSTRUCTION.....	ZSIGN	179	N # # (SEE NOTE).....		
BRIDGE CTR USING FIRST LOC ONLY PAR. <i>xx</i> .....	CYCLE	156	NO BRANCHES GIVEN.....	ZSIGN	179
CHAN ENTRY INVALID. TWO CHANNELS ASSUMED..	DIOCS	24	NO BRANCH OR SWITCH TO BE SET IN INPUT....	LOGIC	170
CHPT ENTRY INVALID. CHPT INCLUDED.....	DIOCS	24	NO FIELD SIZE. PAR. <i>xx</i> .....	ZERO	189
CHPT REQUIRES EOR1. EOR1 GENERATED.....	DIOCS	24		FILL	197
CODENAME NOT DEFINED BY A CODE.....	DECOD	161	NO FIELD SIZE. PARAMETER <i>xx</i> .....	MOVE	215
CODEVALUE NOT DEFINED UNDER A CODE.....	DECOD	162	NO FROM-FIELD IN MOVE MACRO LINE.....	MOVE	215
CODEVALUES AND BRANCHES NOT PAIRED.....	DECOD	162	NO TO-FIELD IN MOVE MACRO LINE.....	MOVE	215
COUNT GREATER THAN FIELD-SIZE.....	SHIFT	222	NOTHING TO TEST IN LOGIC STATEMENT.....	LOGIC	170
COUNT IS ZERO.....	SHIFT	222	NUMERIC FIELD GREATER THAN 20 DIGITS.....	COMP	146
COUNT NOT AN INDEXWORD.....	SHIFT	222	OPEN ENTRY INVALID. OPEN2 ASSUMED.....	DIOCS	25
EOR ENTRY INVALID. EOR1 ASSUMED.....	DIOCS	24	OPEN3 REQUIRES CHPT. OPEN2 GENERATED.....	DIOCS	25
EQUAL SIGN BEGINS INPUT — WILL IGNORE.....	LOGIC	169	OPERAND BLANK.....	OPEN	111
ERROR — IMPROPER OPERAND.....	CYCLE	156		GET	113
FIELD 1 BLANK.....	COMP	146		PUT	116
FIELD 2 BLANK.....	COMP	146		CLOSE	121
FIELD GREATER THAN 20 DIGITS.....	SHIFT	222	OPERAND HAS TWO PARAMETERS.....	GET	113
FIELD MISSING.....	ZSIGN	179		PUT	116
FIELD 1 NOT ACCEPTABLE.....	COMP	146	OPERAND OUTSIDE OF ALLOWABLE RANGE.....	DUF	31
FIELD 2 NOT ACCEPTABLE.....	COMP	146	OUTPUT SRBFORM4 BLANK, ASSUMED 10.....	PUT	116
FIELD UNACCEPTABLE.....	ZSIGN	179	PARAM 01 (03) FILE FORM INVALID.....	PUTX	119
	COUNT	222	PARAM 03 IS A FILE. PARAM 03 IGNORED.....	GET	113
FILLING INSTRUCTION. PAR. <i>xx</i> .....	FILL	197	PARAM011 NOT A FILE.....	OPEN	111
HIGH-ORDER DIGITS LOST OF PARAMETER <i>xx</i> .....	EDMOV	209		CLOSE	121
IGEN ENTRY INVALID. IGEN1 ASSUMED.....	DIOCS	25	PARAM 01 NOT A FILE.....	GET	113
ILLEGAL PUNCTUATION MARK USED.....	LOGIC	169	PARAM 01(03) NOT A FILE.....	PUT	116
ILLEGAL TERM ENDS LOGICAL EXPRESSION.....	LOGIC	169		PUTX	119
ILLEGAL TERM PRECEDES A BINARY OPERATOR...	LOGIC	169	PARAM 03 NOT DEFINED.....	GET	114
ILLEGAL TERM PRECEDES A LEFT PAREN.....	LOGIC	169	PARAM 01 NOT INPUT FILE.....	GET	113
ILLEGAL TERM PRECEDES A NOT.....	LOGIC	169	PARAM 01(03) NOT OUTPUT FILE.....	PUT	116
ILLEGAL TERM PRECEDES A RIGHT PAREN.....	LOGIC	170	PARAM NOT LABEL OF CYCLE MACRO <i>xx</i> .....	RECYC	156
ILLEGAL TERM PRECEDES PARAMETER <i>xx</i> .....	LOGIC	170	PARAM 03 — SRBFORM4 BLANK, ASSUMED 10....	PUTX	119
IMPROPER OPERAND.....	PUTX	119	PARAM 01 UNDEFINED.....	PUT	116
IMPROPER OPERAND. CHAN2, OPEN2, EOR1, CHPT,			PAREN. MISSING AROUND ARITH-REL.....	LOGIC	170
IGEN1 ASSUMED.....	DIOCS	25	PARENTHESIS NOT CLOSED.....	LOGIC	170
			PRINT SUPPRESSED IF ALL ZERO.....	EDMOV	209
			RECLENGTHS UNEQUAL.....	PUTX	119

<i>Message</i>	<i>Macro</i>	<i>Page</i>
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START GREATER THAN FIELD-SIZE .....	SHIFT	222
STMNT SHOULD OR SEEMS TO BE ENDED BUT CARDS REMAIN .....	DUF	31
SYMBOLIC ADDRESS NOT ALLOWED .....	DUF	31
TO-FIELD(S) SMALLER THAN FROM-FIELD(S) .....	MOVE	215
TO-FIELD NOT ALPHA. PARAMETER xx .....	EDMOV	209
TOO MANY PARAMETERS. WILL IGNORE .....	SNAP	226
TOO MANY RIGHT PARENTHESES .....	LOGIC	170
UNACCEPTABLE PARAMETER xx .....	EDMOV	209
	MOVE	215
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w — BOTH FIELDS NOT ALPHA — NOFORM .....	COMP	146
w — UNUSUAL BRANCH CONDITION .....	COMP	146

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x # # (SEE NOTE)		
ZEROING DC. PAR. xx .....	ZERO	189
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NOTE. For ARITH messages, which begin with an  
n, w or x followed by a two-digit number,  
see page 133.

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